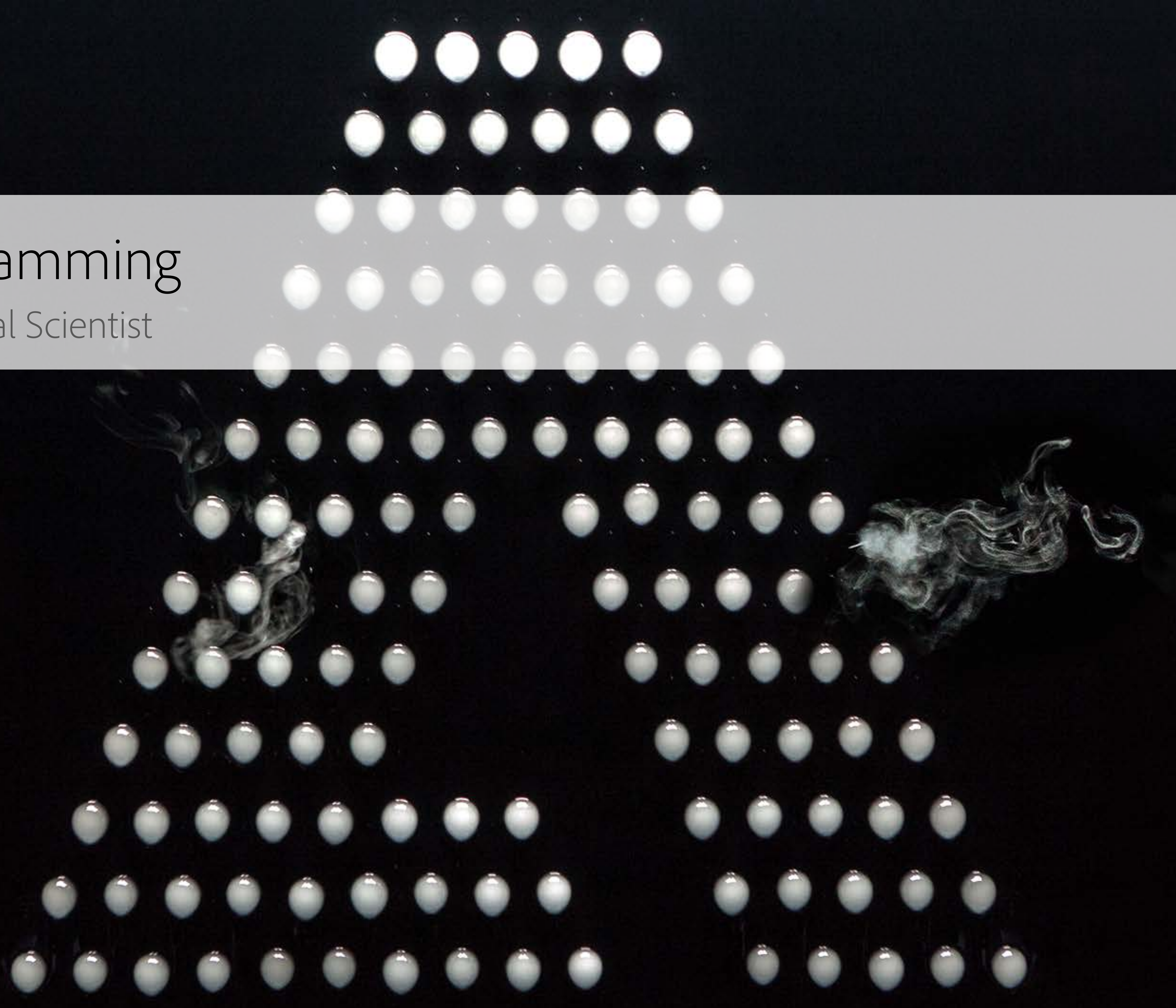
Generic Programming

Sean Parent | Principal Scientist



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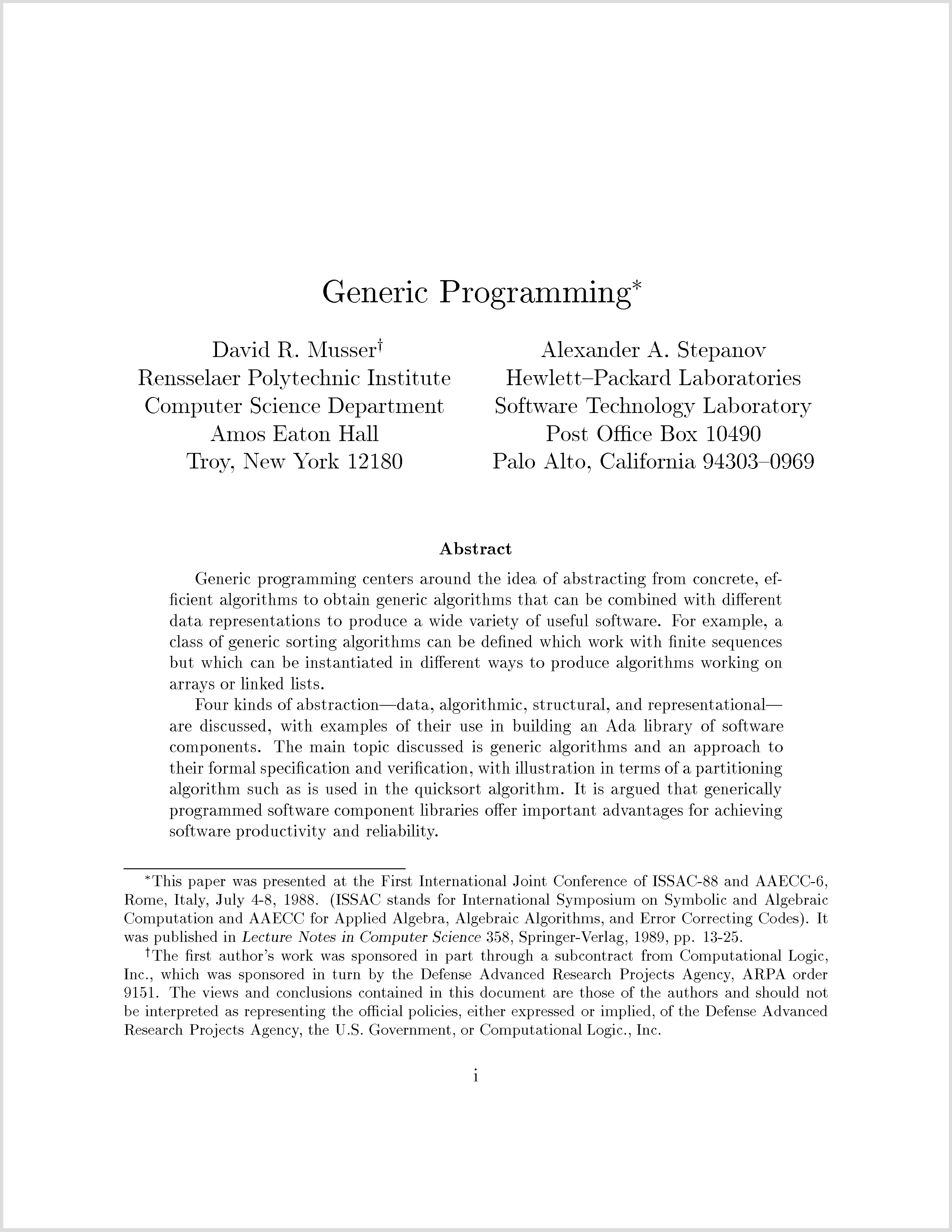
Thomas Wirtz

“You cannot fully grasp mathematics until you understand its historical context.” – Alex Stepanov

1988



Generic Programming*

David R. Musser[†]
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Computer Science Department
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Abstract

Generic programming centers around the idea of abstracting from concrete, efficient algorithms to obtain generic algorithms that can be combined with different data representations to produce a wide variety of useful software. For example, a class of generic sorting algorithms can be defined which work with finite sequences but which can be instantiated in different ways to produce algorithms working on arrays or linked lists.

Four kinds of abstraction—data, algorithmic, structural, and representational—are discussed, with examples of their use in building an Ada library of software components. The main topic discussed is generic algorithms and an approach to their formal specification and verification, with illustration in terms of a partitioning algorithm such as is used in the quicksort algorithm. It is argued that generically programmed software component libraries offer important advantages for achieving software productivity and reliability.

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[*] This paper was presented at the First International Joint Conference of ISSAC-88 and AAECC-6, Rome, Italy, July 4-8, 1988. (ISSAC stands for International Symposium on Symbolic and Algebraic Computation and AAECC for Applied Algebra, Algebraic Algorithms, and Error Correcting Codes). It was published in *Lecture Notes in Computer Science* 358, Springer-Verlag, 1989, pp. 13-25.

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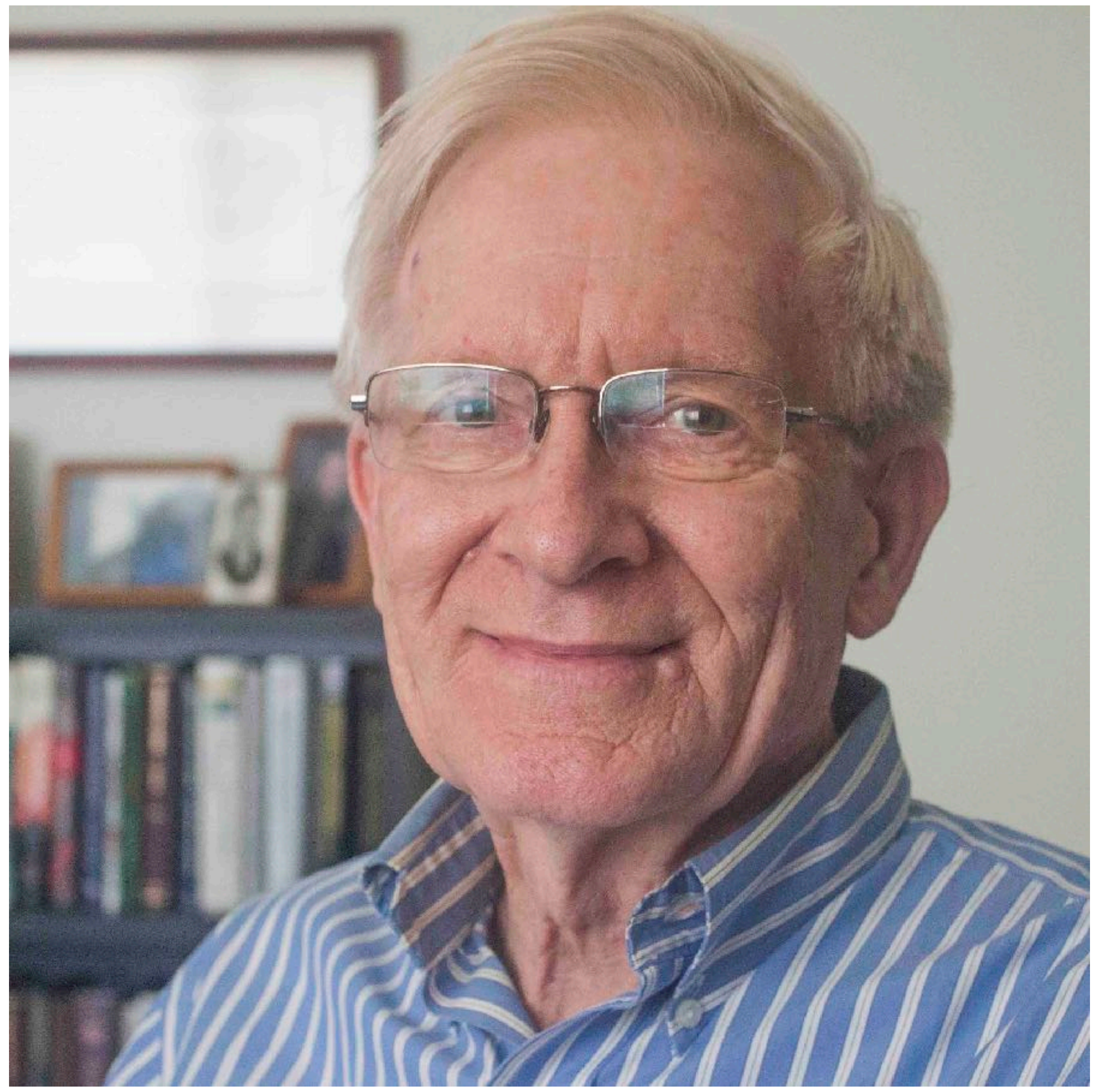

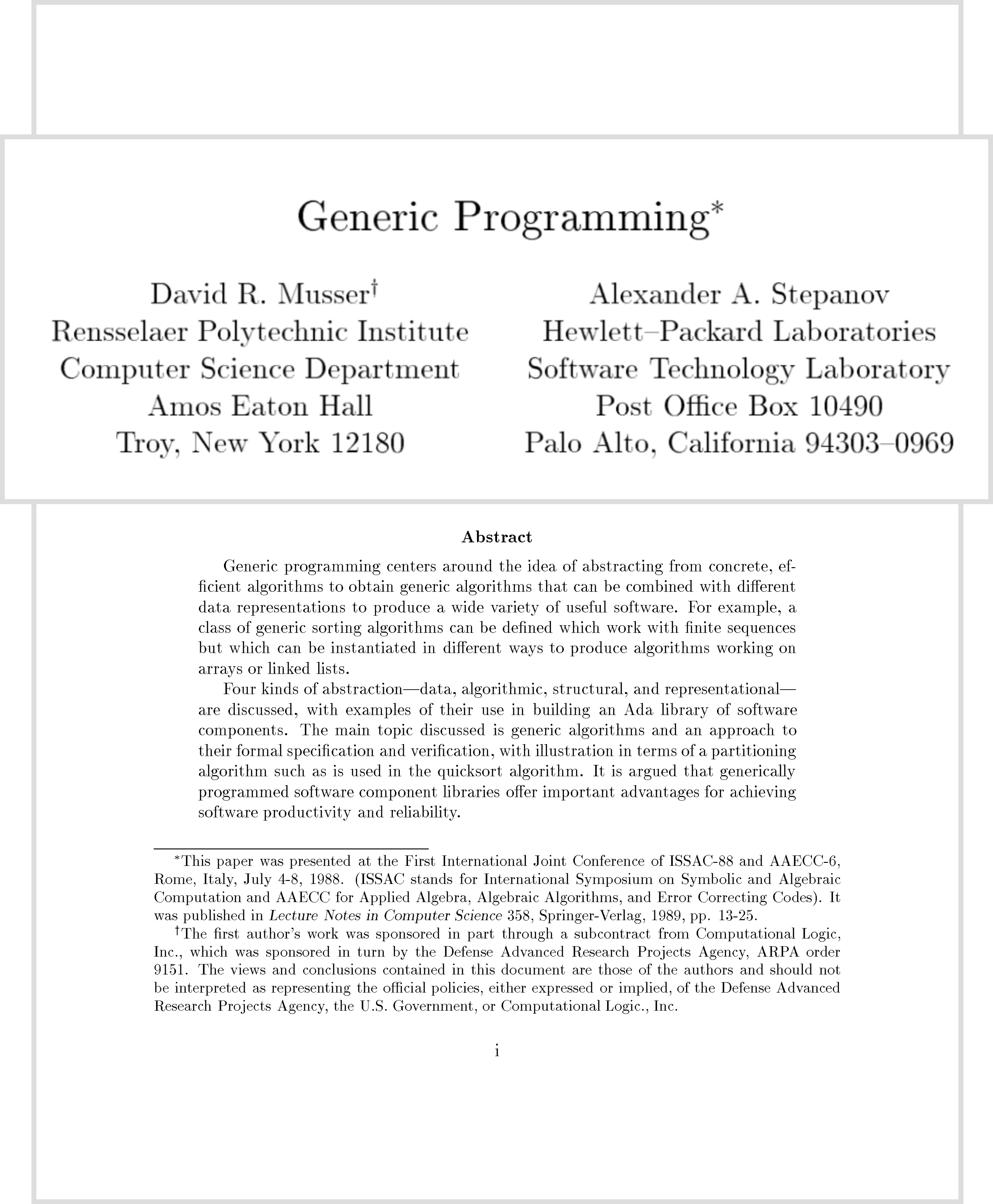

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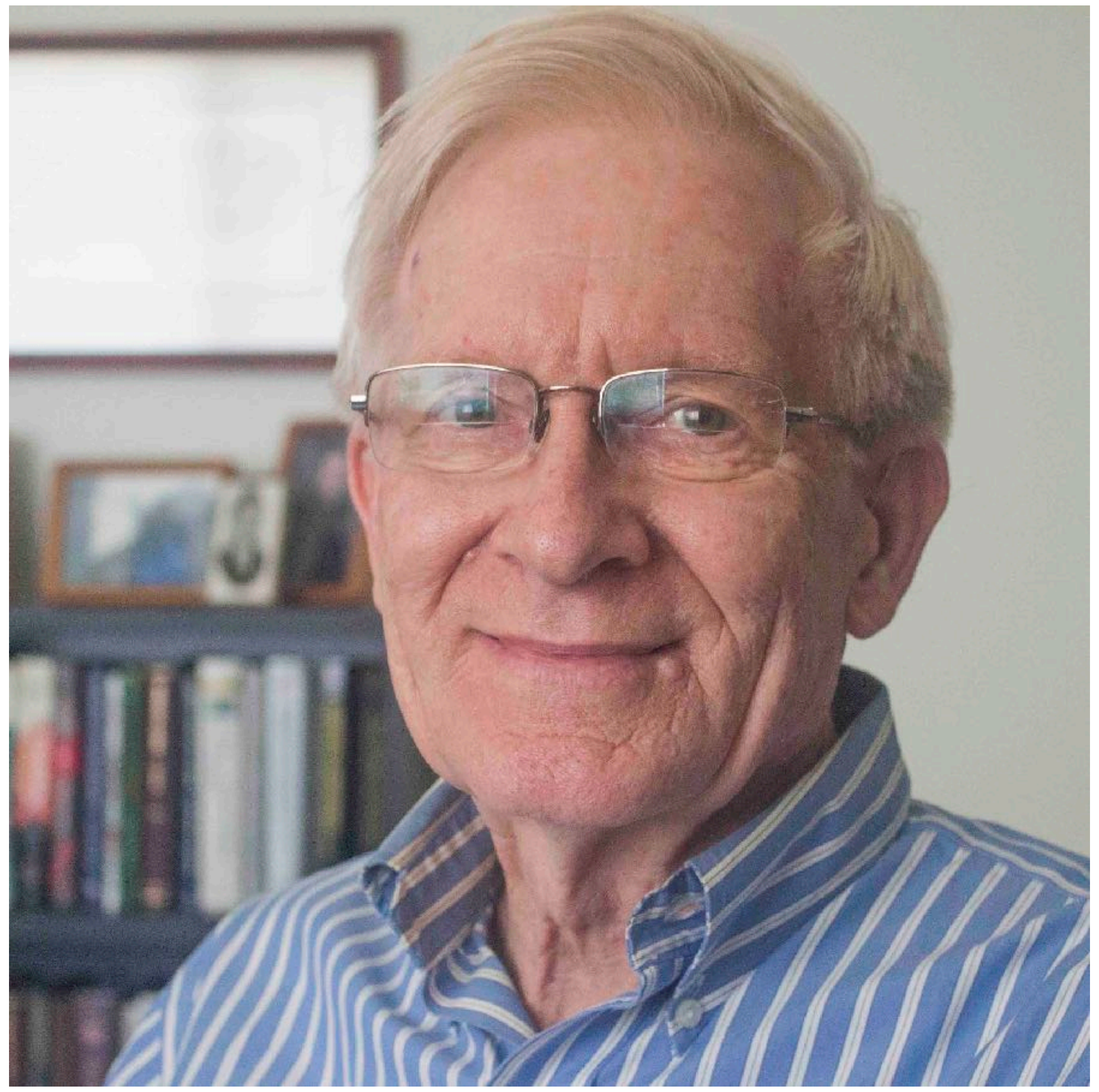

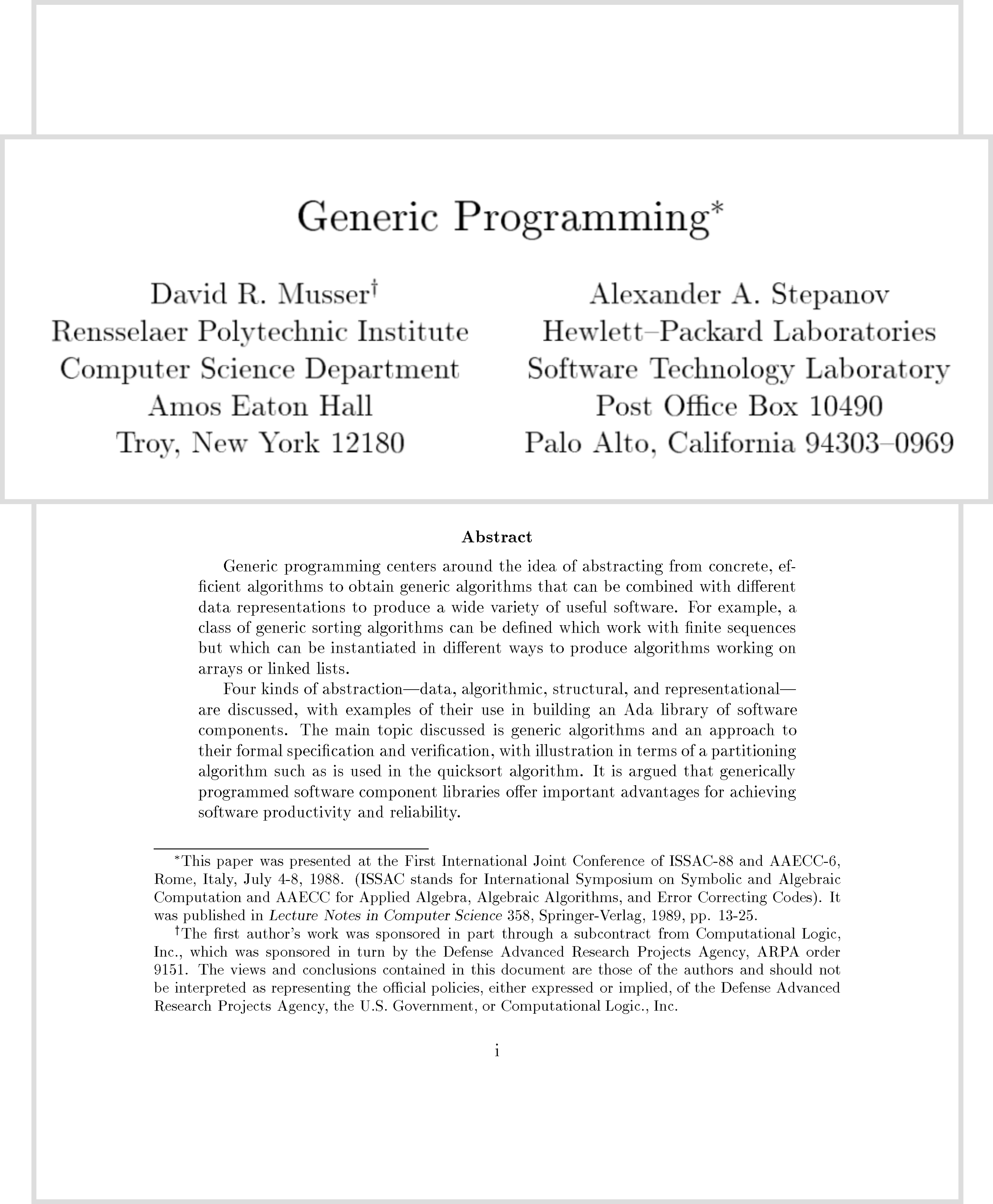

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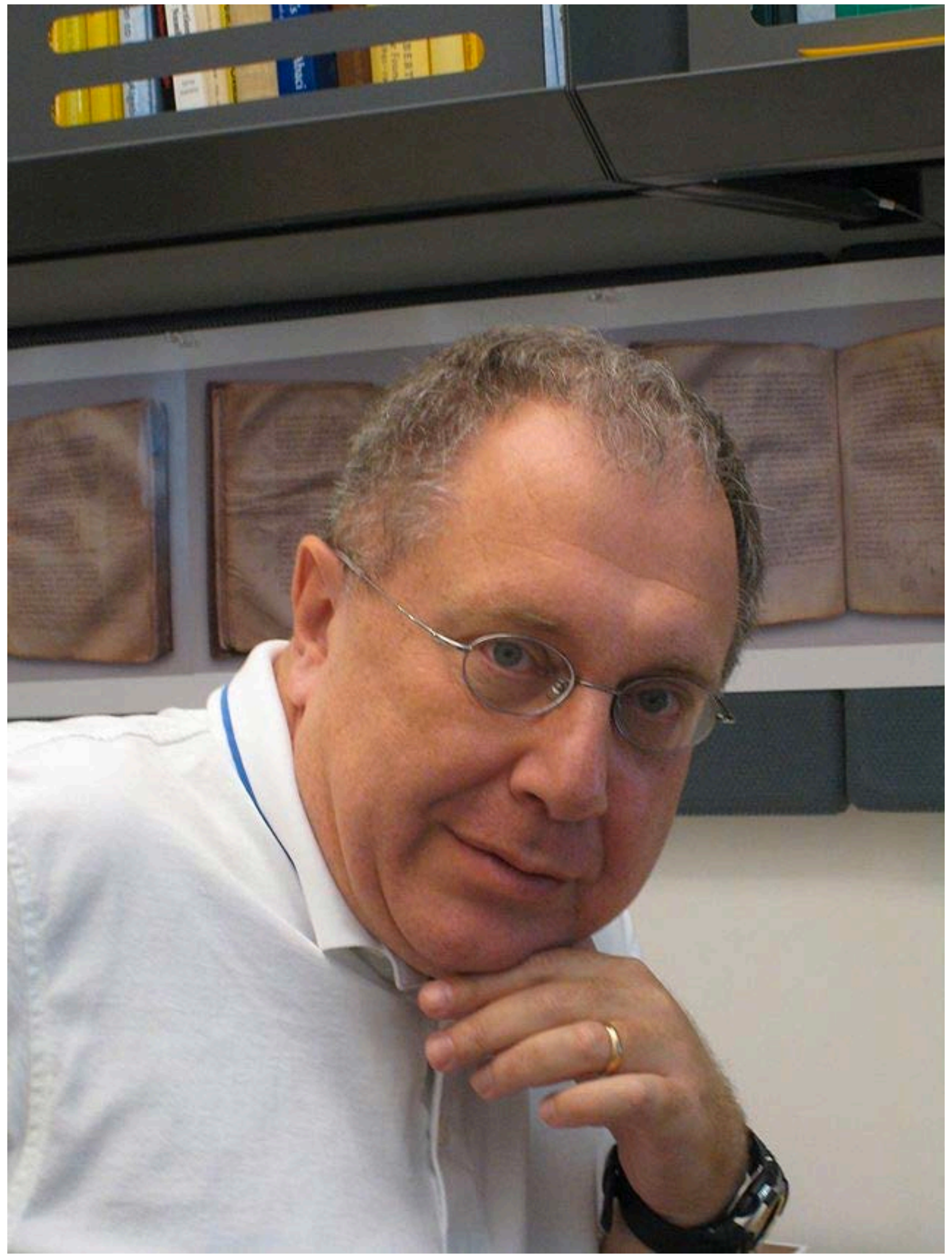

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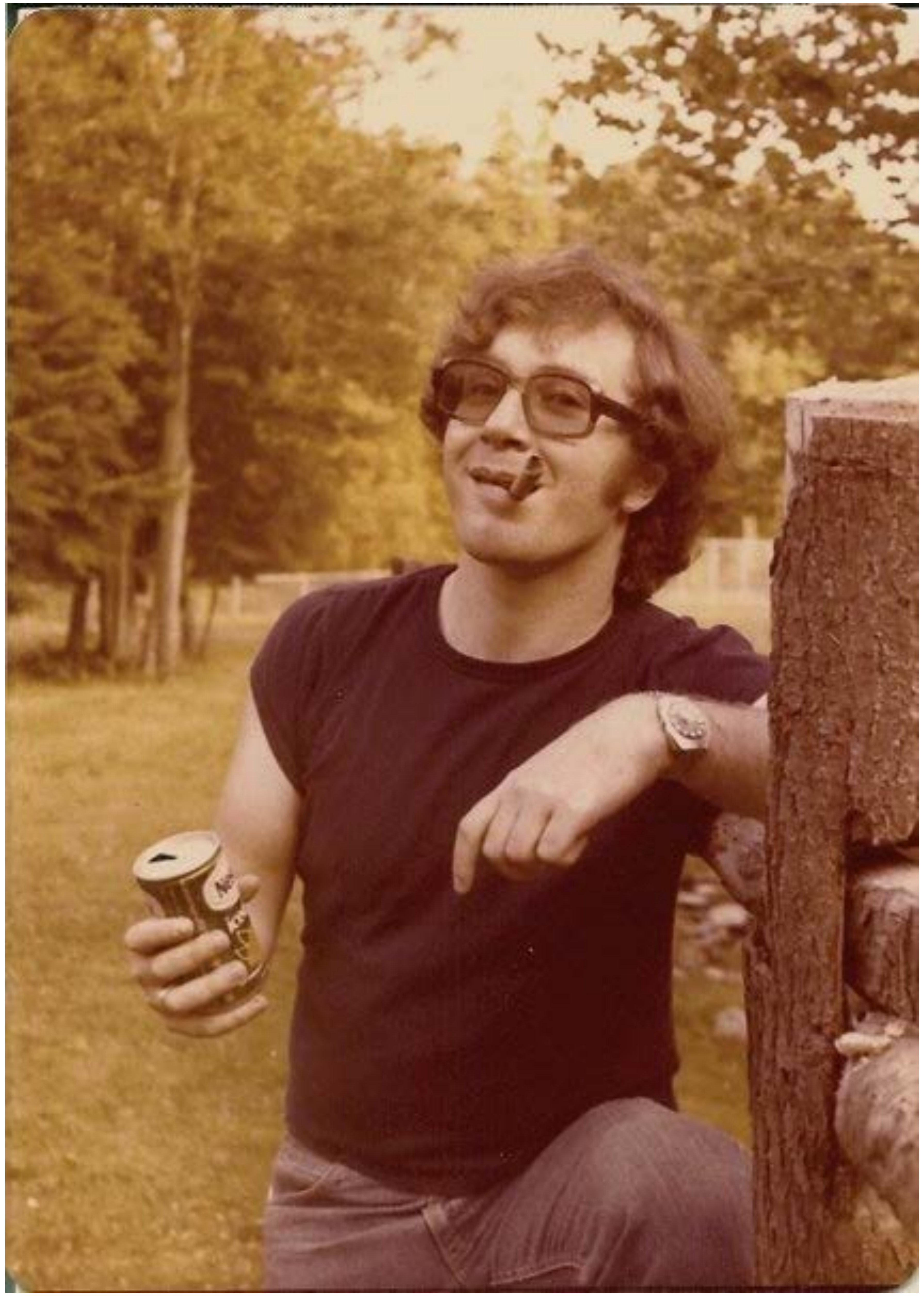
1976-1987



1976 Parallel Computation and Associative Property

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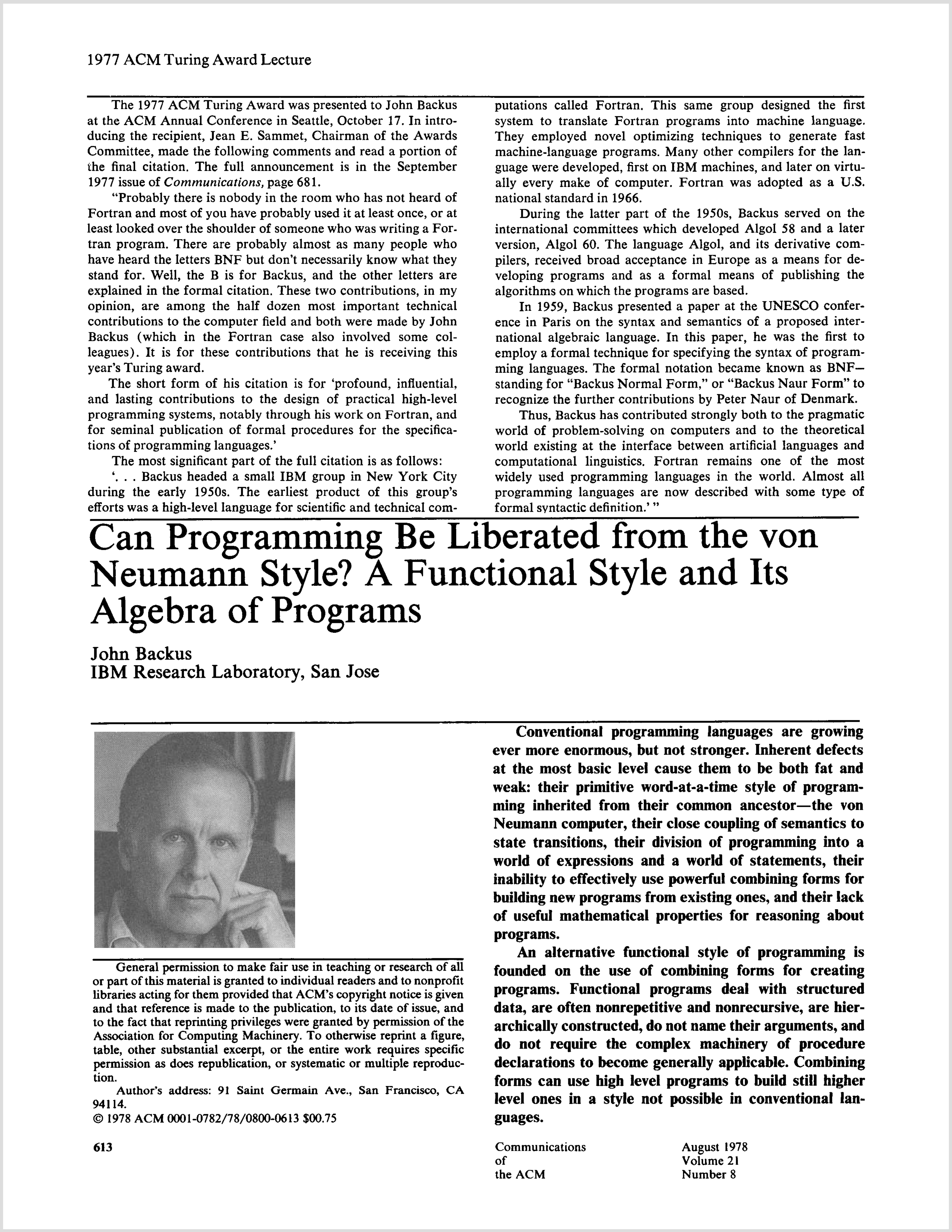
Parallel reduction is associated with monoids

Software is associated with Algebraic Structures

1977 John Backus



The 1977 ACM Turing Award was presented to John Backus at the ACM Annual Conference in Seattle, October 17. In introducing the recipient, Jean E. Sammet, Chairman of the Awards Committee, made the following comments and read a portion of the final citation. The full announcement is in the September 1977 issue of *Communications,* page 681.

"Probably there is nobody in the room who has not heard of Fortran and most of you have probably used it at least once, or at least looked over the shoulder of someone who was writing a Fortran program. There are probably almost as many people who have heard the letters BNF but don't necessarily know what they stand for. Well, the B is for Backus, and the other letters are explained in the formal citation. These two contributions, in my opinion, are among the half dozen most important technical contributions to the computer field and both were made by John Backus (which in the Fortran case also involved some colleagues). It is for these contributions that he is receiving this year's Turing award.

The short form of his citation is for 'profound, influential, and lasting contributions to the design of practical high-level programming systems, notably through his work on Fortran, and for seminal publication of formal procedures for the specifications of programming languages.'

The most significant part of the full citation is as follows:

'. . . Backus headed a small IBM group in New York City during the early 1950s. The earliest product of this group's efforts was a high-level language for scientific and technical computations called Fortran. This same group designed the first system to translate Fortran programs into machine language. They employed novel optimizing techniques to generate fast machine-language programs. Many other compilers for the language were developed, first on IBM machines, and later on virtually every make of computer. Fortran was adopted as a U.S. national standard in 1966.

During the latter part of the 1950s, Backus served on the international committees which developed Algol 58 and a later version, Algol 60. The language Algol, and its derivative compilers, received broad acceptance in Europe as a means for developing programs and as a formal means of publishing the algorithms on which the programs are based.

In 1959, Backus presented a paper at the UNESCO conference in Paris on the syntax and semantics of a proposed international algebraic language. In this paper, he was the first to employ a formal technique for specifying the syntax of programming languages. The formal notation became known as BNF—standing for "Backus Normal Form," or "Backus Naur Form" to recognize the further contributions by Peter Naur of Denmark.

Thus, Backus has contributed strongly both to the pragmatic world of problem-solving on computers and to the theoretical world existing at the interface between artificial languages and computational linguistics. Fortran remains one of the most widely used programming languages in the world. Almost all programming languages are now described with some type of formal syntactic definition.' "

Can Programming Be Liberated from the von Neumann Style? A Functional Style and Its Algebra of Programs

John Backus
IBM Research Laboratory, San Jose

Conventional programming languages are growing ever more enormous, but not stronger. Inherent defects at the most basic level cause them to be both fat and weak: their primitive word-at-a-time style of programming inherited from their common ancestor—the von Neumann computer, their close coupling of semantics to state transitions, their division of programming into a world of expressions and a world of statements, their inability to effectively use powerful combining forms for building new programs from existing ones, and their lack of useful mathematical properties for reasoning about programs.

An alternative functional style of programming is founded on the use of combining forms for creating programs. Functional programs deal with structured data, are often nonrepetitive and nonrecursive, are hierarchically constructed, do not name their arguments, and do not require the complex machinery of procedure declarations to become generally applicable. Combining forms can use high level programs to build still higher level ones in a style not possible in conventional languages.

Author's address: 91 Saint Germain Ave., San Francisco, CA 94114.

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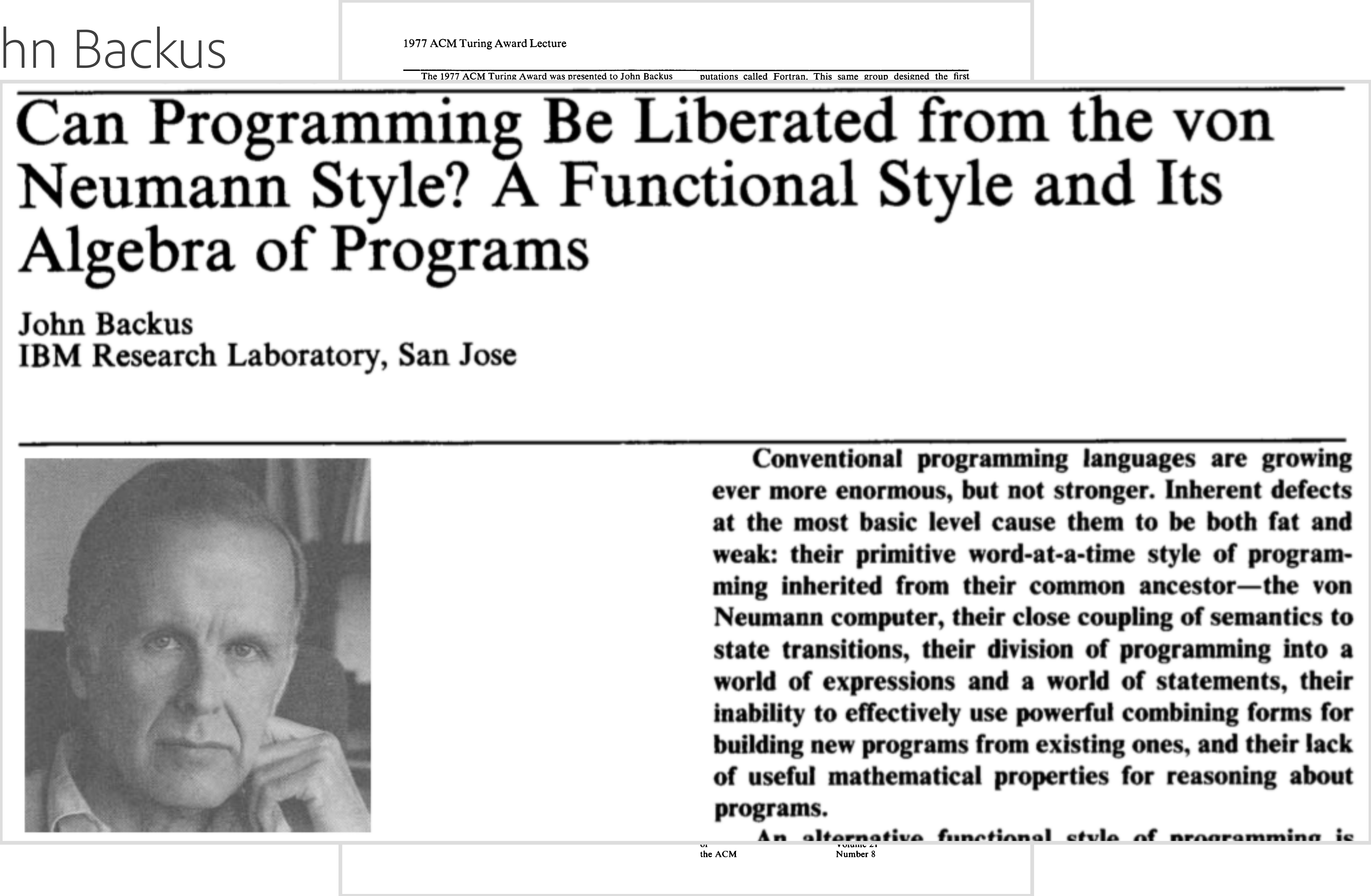

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Volume 21
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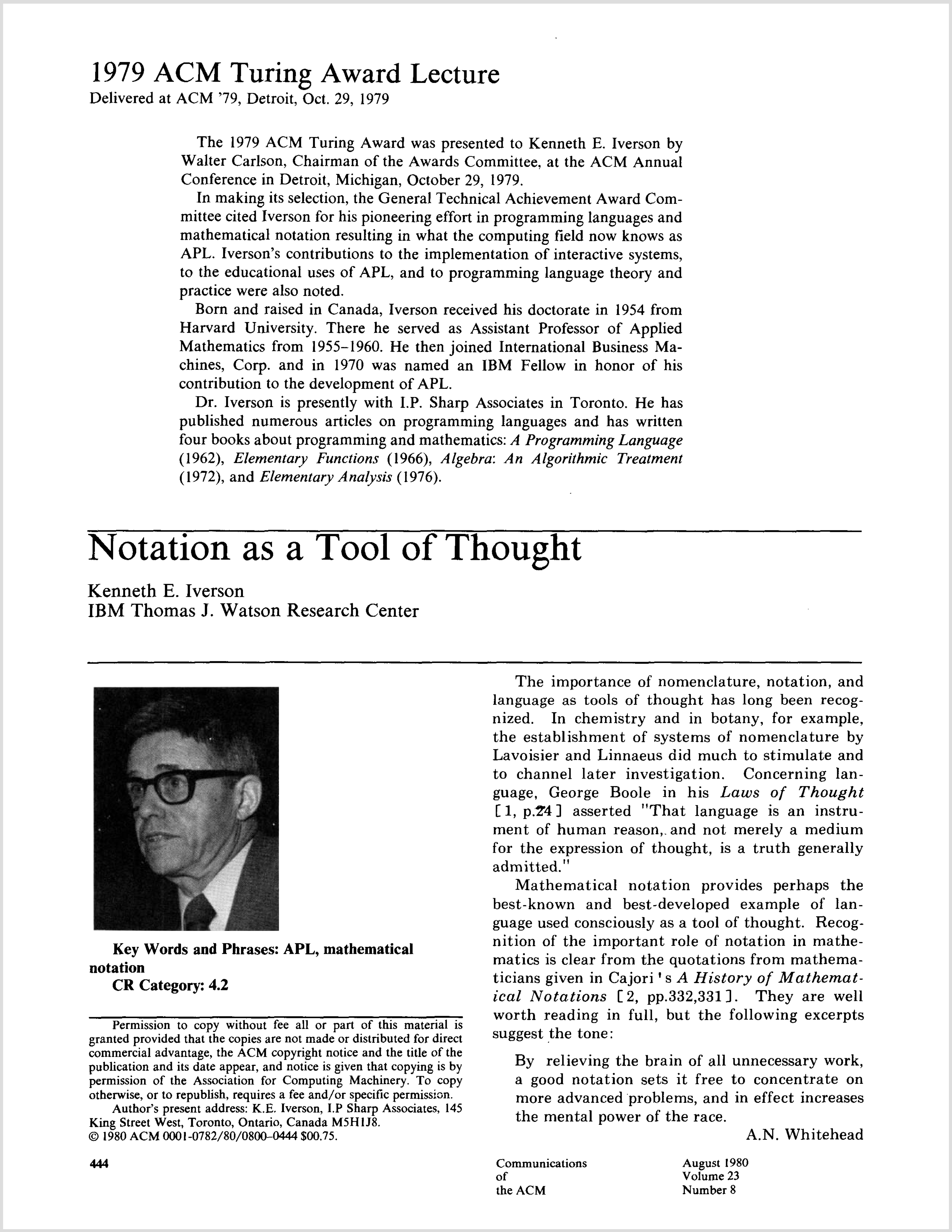
1979 Ken Iverson



1979 ACM Turing Award Lecture
Delivered at ACM '79, Detroit, Oct. 29, 1979

The 1979 ACM Turing Award was presented to Kenneth E. Iverson by Walter Carlson, Chairman of the Awards Committee, at the ACM Annual Conference in Detroit, Michigan, October 29, 1979.

In making its selection, the General Technical Achievement Award Committee cited Iverson for his pioneering effort in programming languages and mathematical notation resulting in what the computing field now knows as APL. Iverson's contributions to the implementation of interactive systems, to the educational uses of APL, and to programming language theory and practice were also noted.

Born and raised in Canada, Iverson received his doctorate in 1954 from Harvard University. There he served as Assistant Professor of Applied Mathematics from 1955–1960. He then joined International Business Machines, Corp. and in 1970 was named an IBM Fellow in honor of his contribution to the development of APL.

Dr. Iverson is presently with I.P. Sharp Associates in Toronto. He has published numerous articles on programming languages and has written four books about programming and mathematics: *A Programming Language* (1962), *Elementary Functions* (1966), *Algebra: An Algorithmic Treatment* (1972), and *Elementary Analysis* (1976).

Notation as a Tool of Thought

Kenneth E. Iverson
IBM Thomas J. Watson Research Center

Key Words and Phrases: APL, mathematical notation
CR Category: 4.2

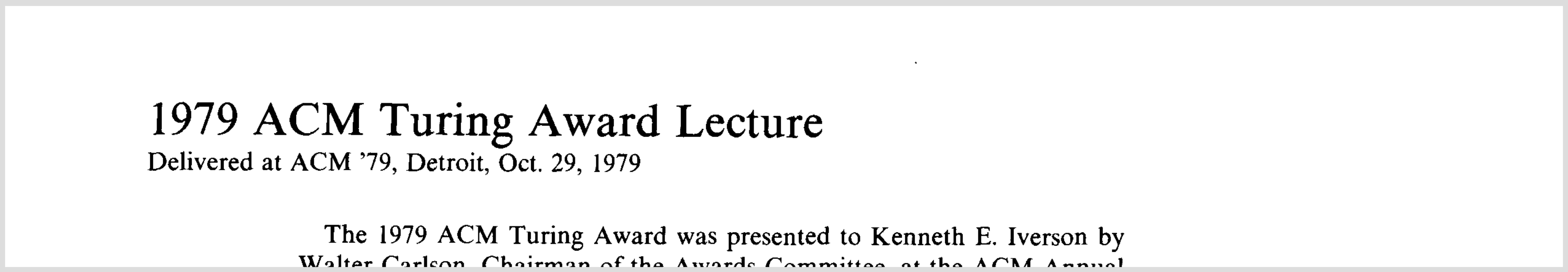
The importance of nomenclature, notation, and language as tools of thought has long been recognized. In chemistry and in botany, for example, the establishment of systems of nomenclature by Lavoisier and Linnaeus did much to stimulate and to channel later investigation. Concerning language, George Boole in his *Laws of Thought* [1, p.24] asserted "That language is an instrument of human reason, and not merely a medium for the expression of thought, is a truth generally admitted."

Mathematical notation provides perhaps the best-known and best-developed example of language used consciously as a tool of thought. Recognition of the important role of notation in mathematics is clear from the quotations from mathematicians given in Cajori's *A History of Mathematical Notations* [2, pp.332,331]. They are well worth reading in full, but the following excerpts suggest the tone:

> By relieving the brain of all unnecessary work, a good notation sets it free to concentrate on more advanced problems, and in effect increases the mental power of the race.
>
> A.N. Whitehead

Author's present address: K.E. Iverson, I.P Sharp Associates, 145 King Street West, Toronto, Ontario, Canada M5H1J8.

1979 Ken Iverson



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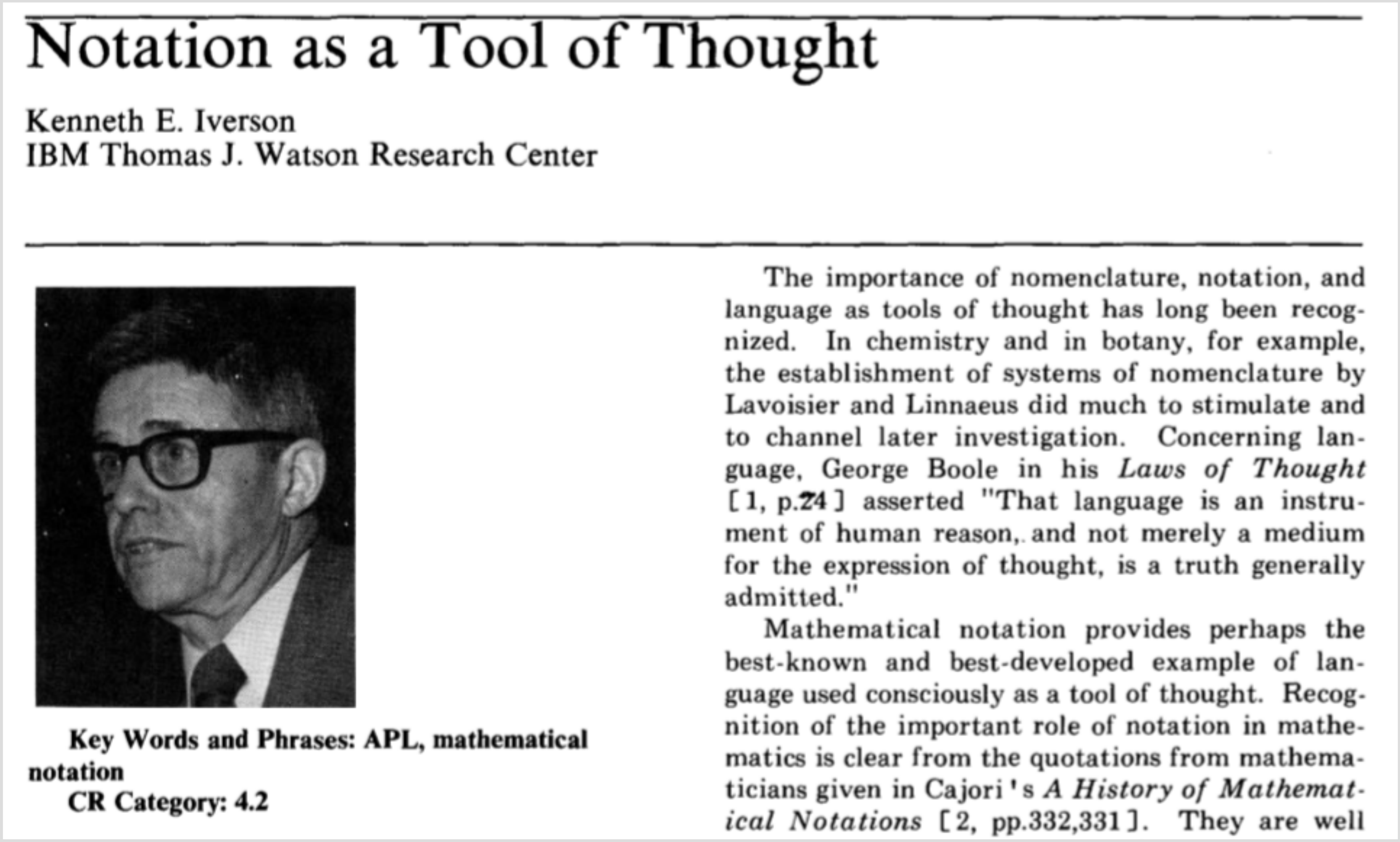

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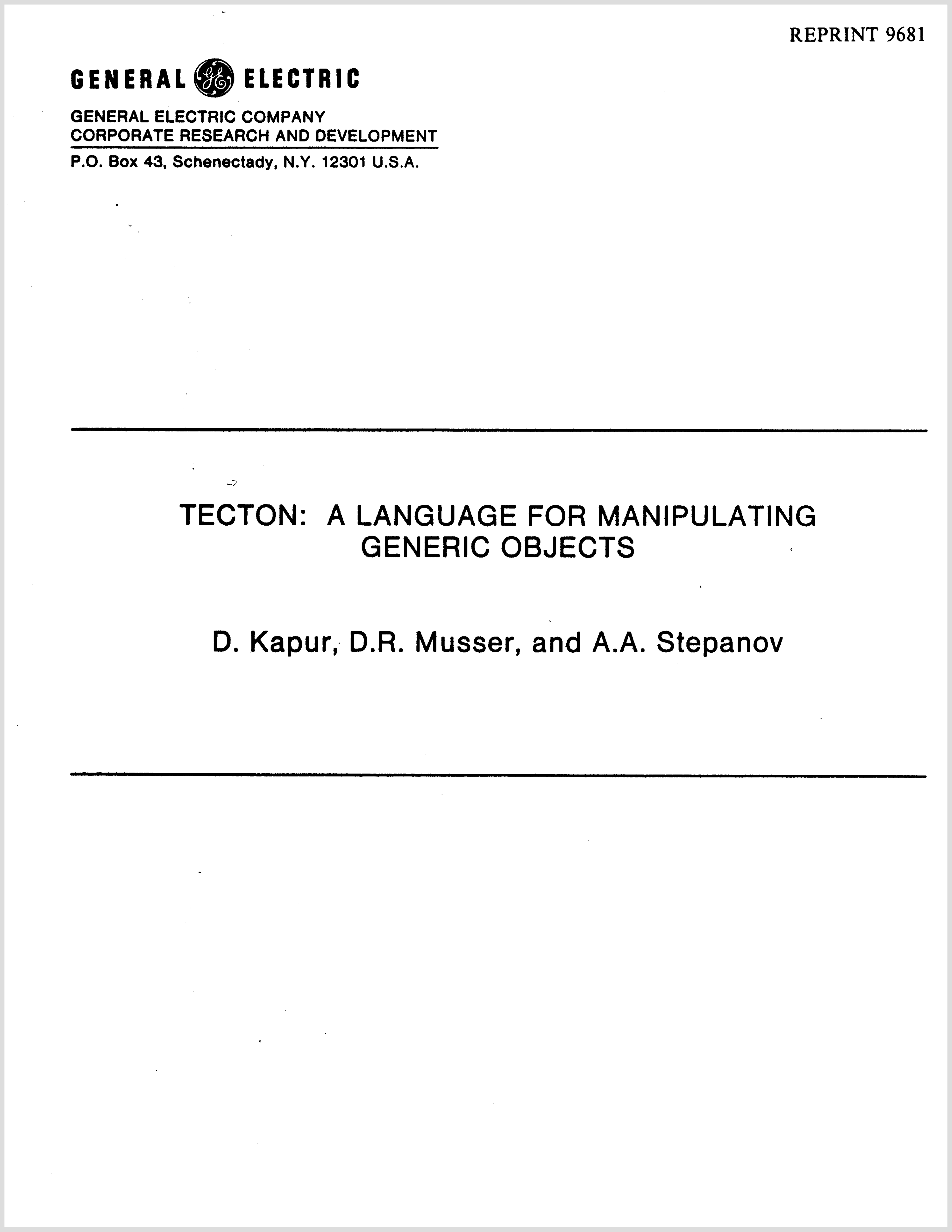
1981 Tecton

The Tecton language



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TECTON: A LANGUAGE FOR MANIPULATING GENERIC OBJECTS

D. Kapur, D.R. Musser, and A.A. Stepanov

Adobe

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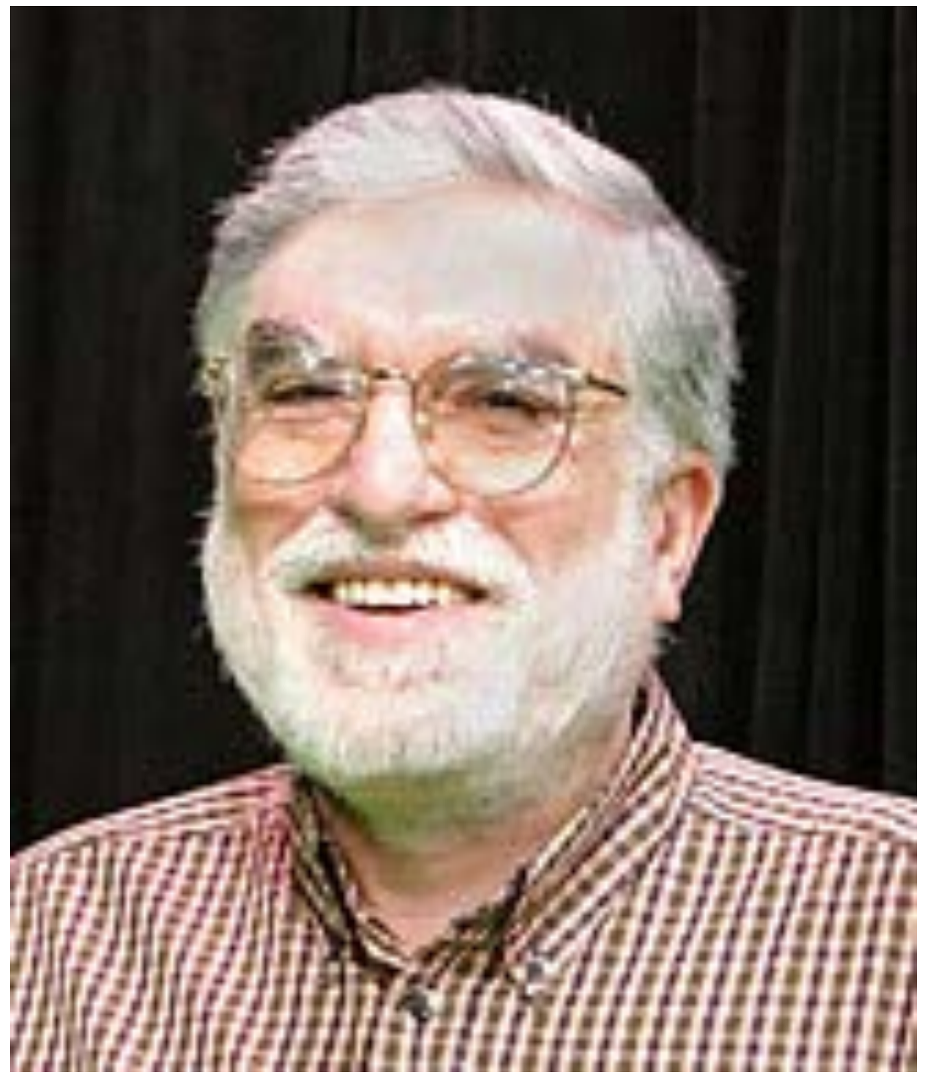
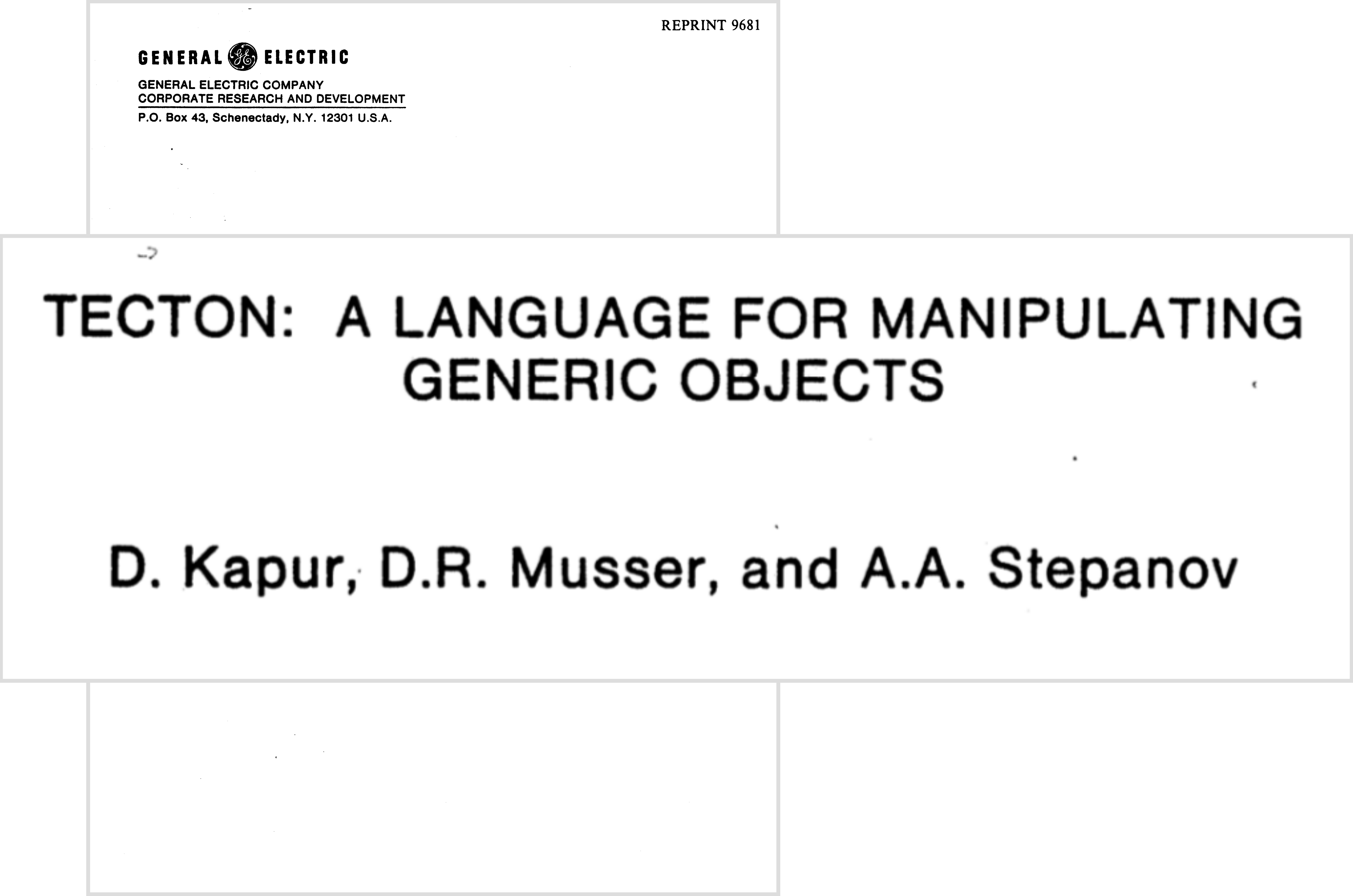
The Tecton language





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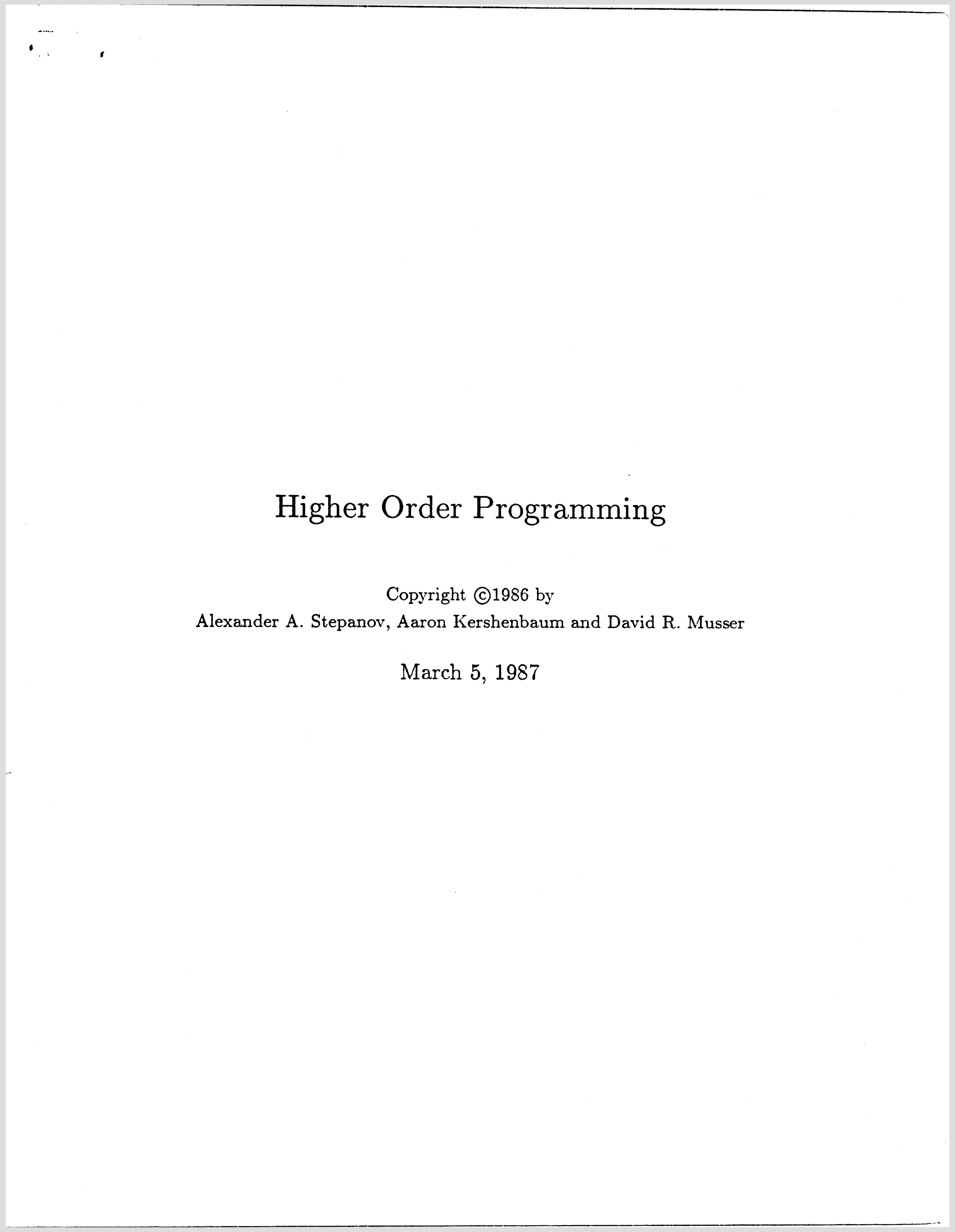
1986-87 Libraries

Higher Order Programming



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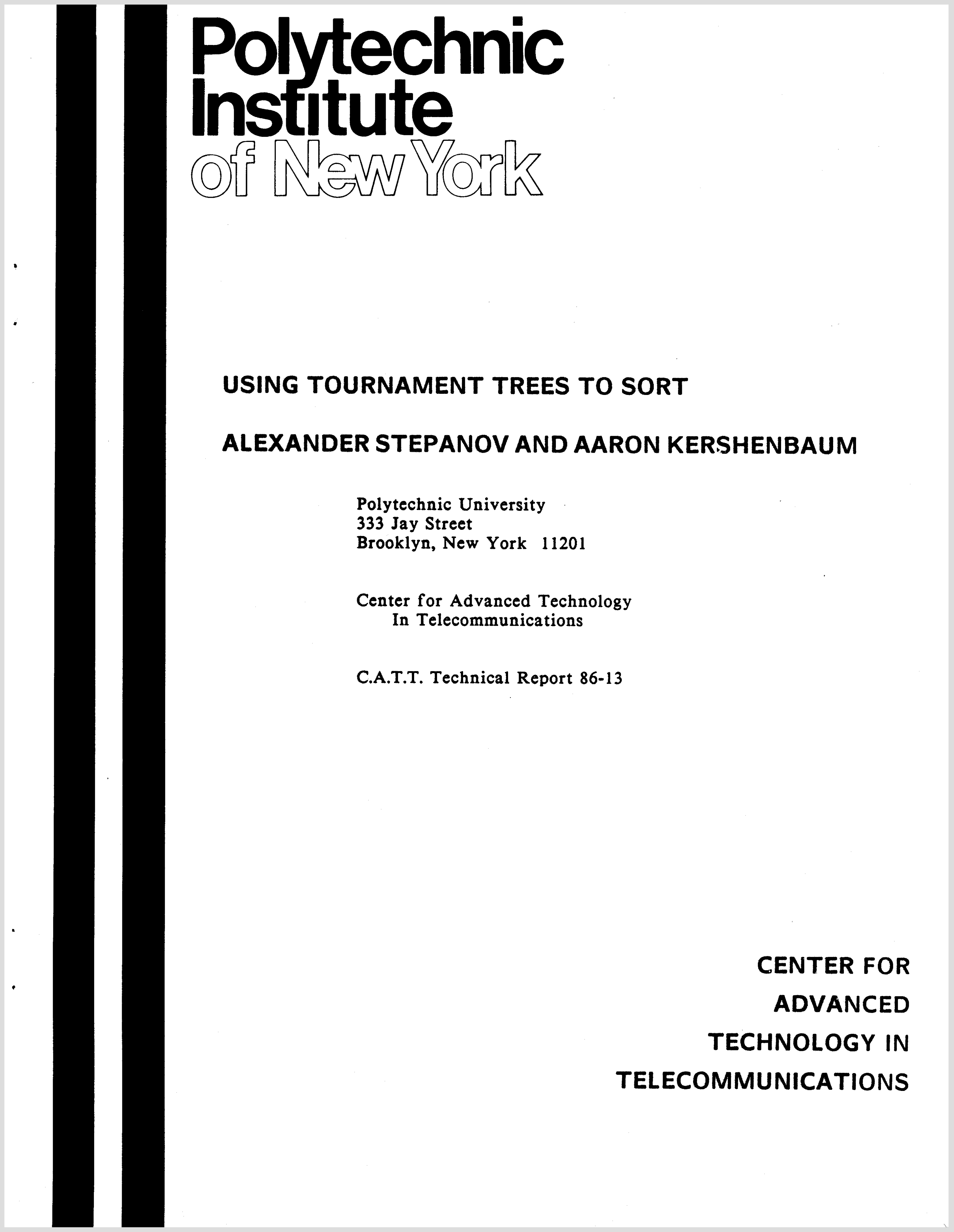
March 5, 1987



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ALEXANDER STEPANOV AND AARON KERSHENBAUM

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333 Jay Street
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In Telecommunications

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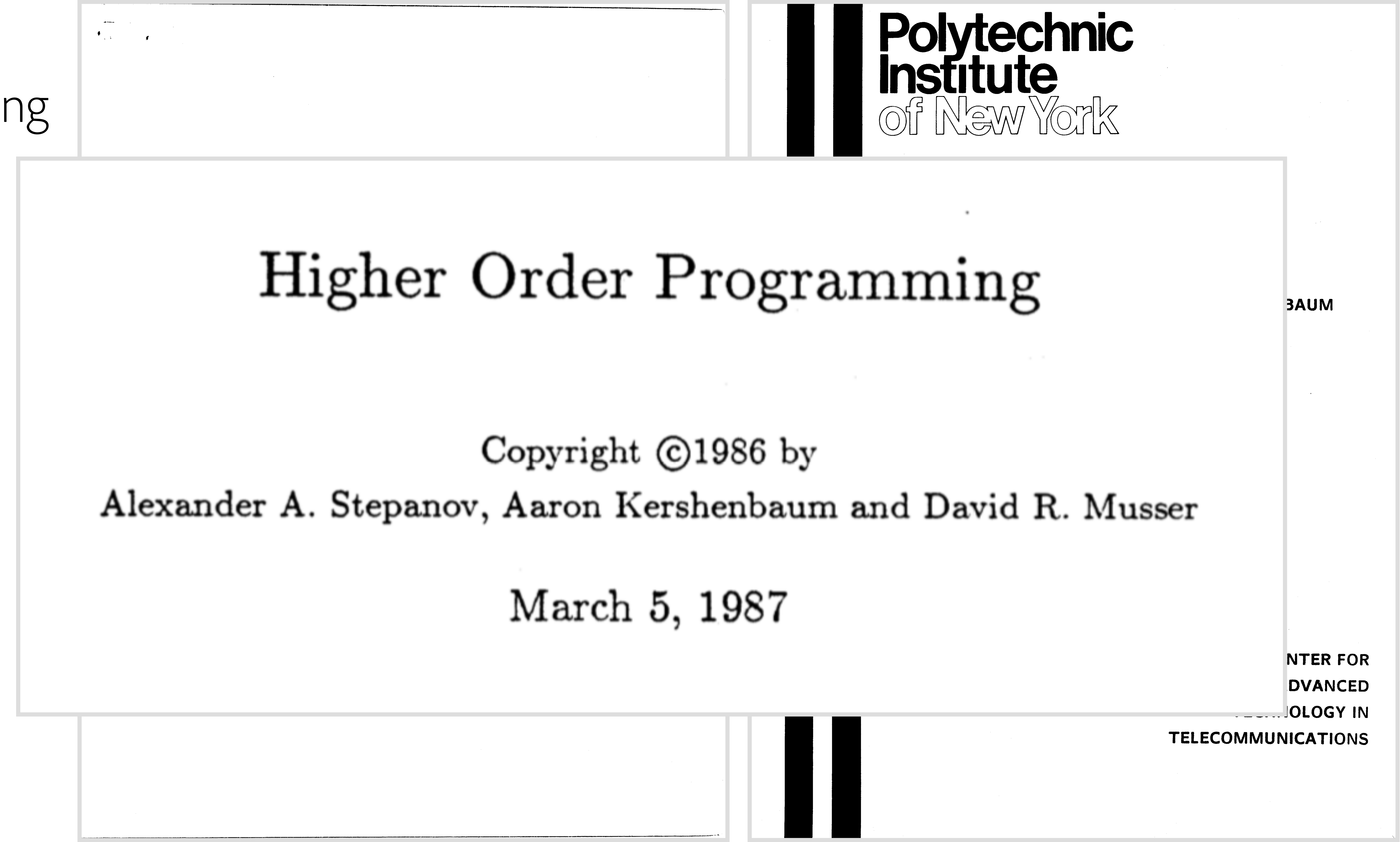

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March 5, 1987

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Alex works briefly at Bell Labs

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“De-Bourbakized”

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    Nicolas Bourbaki

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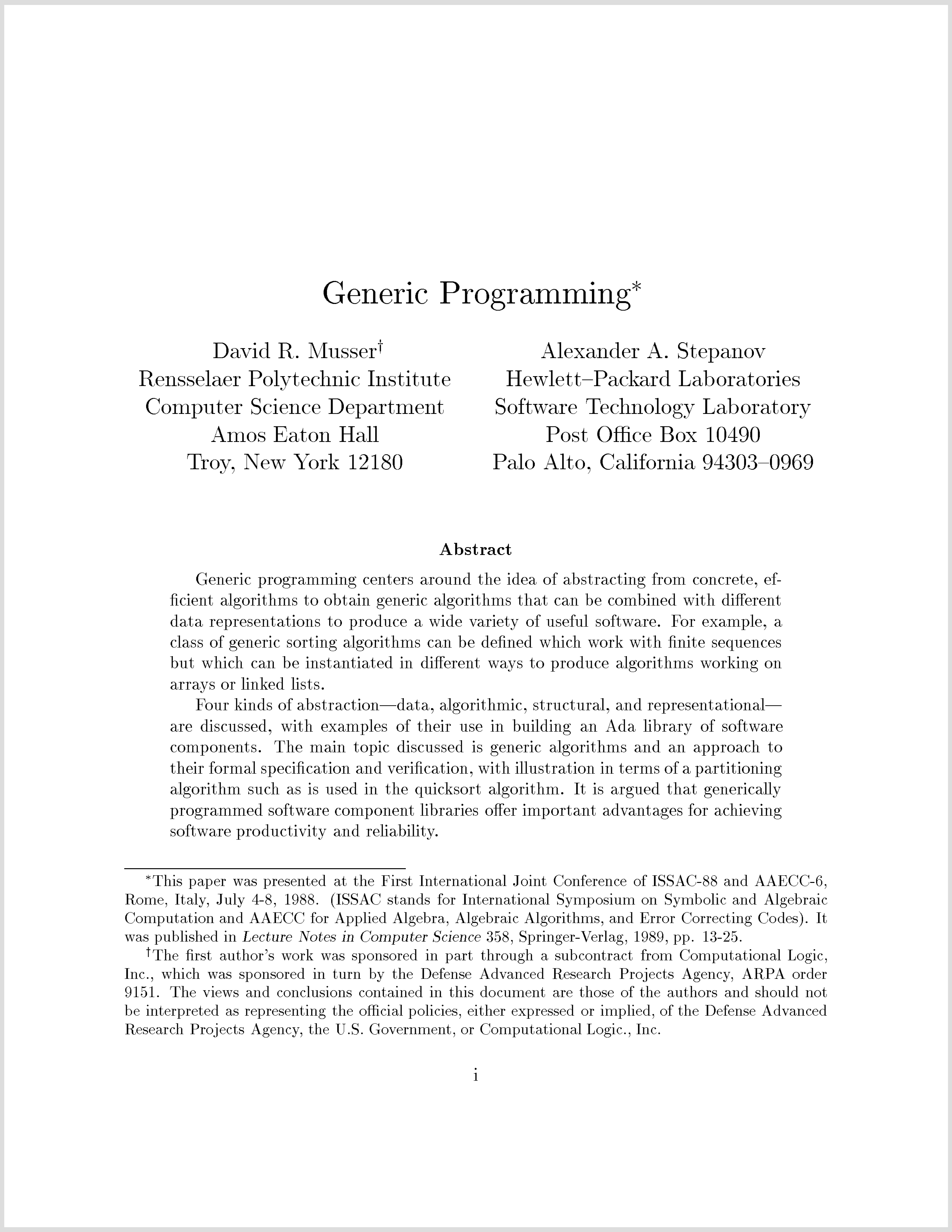
Mathematics is discovery, not invention

Software is defined on Algebraic Structures

1988



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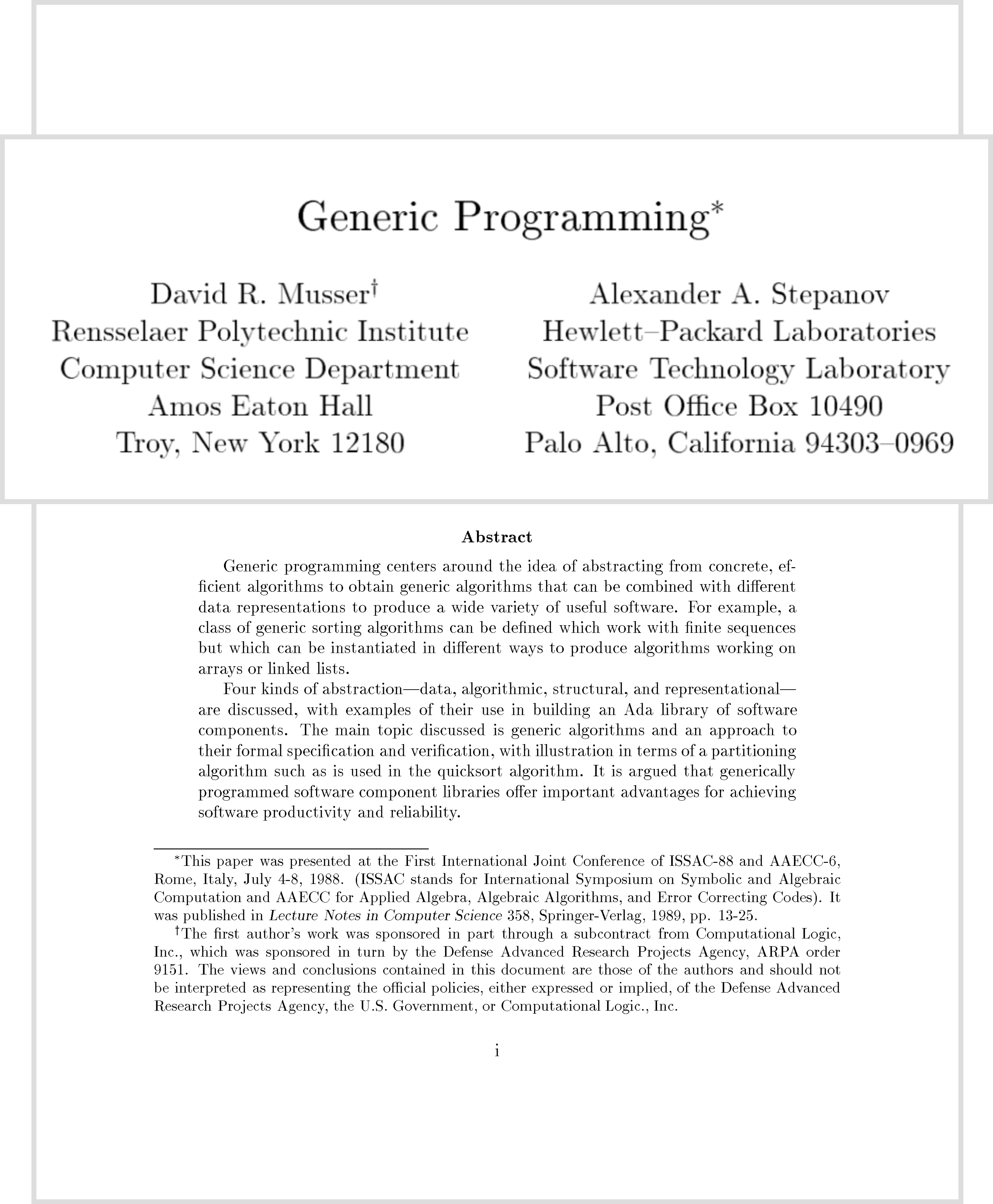

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```
procedure Partition(S         : in out Sequence;
                    F, L      : in Coordinate;
                    Middle    : out Coordinate;
                    Middle_OK : out Boolean) is
  First : Coordinate := F;
  Last  : Coordinate := L;
begin
  loop
    loop
      if First = Last then
        Middle := First;
        Middle_OK := Test(S, First);
        return;
      end if;
      exit when not Test(S, First);
      First := Next(First);
    end loop;
    loop
      exit when Test(S, Last);
      Last := Prev(Last);
      if First = Last then
        Middle := First;
        Middle_OK := False;
        return;
      end if;
    end loop;
    Swap(S, First, Last);
    First := Next(First);
    if First = Last then
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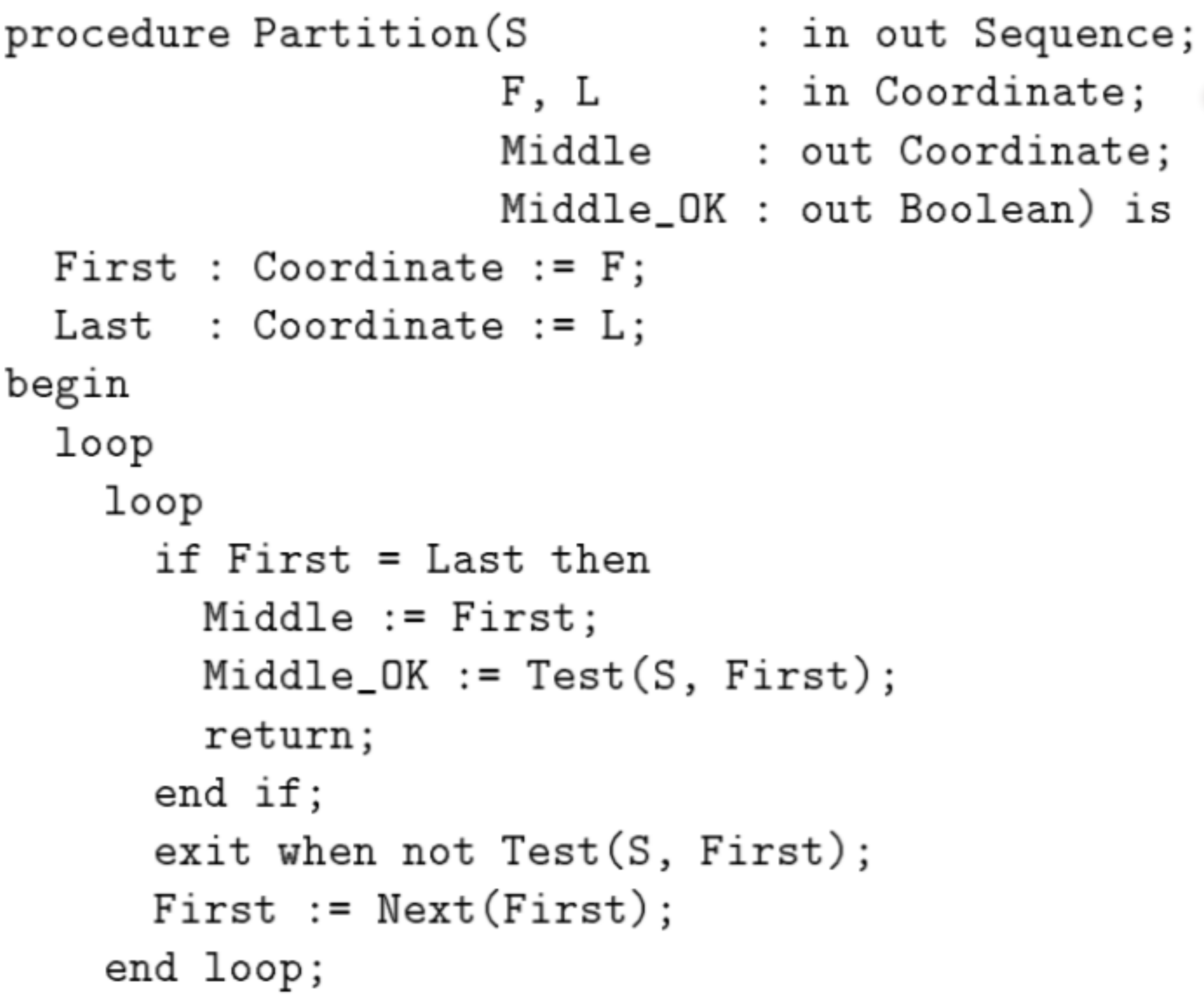
      Middle_OK := False;
      return;
    end if;
    Last := Prev(Last);
  end loop;
end Partition;
```

Figure 1: Body of Partition Algorithm



```
procedure Partition(S          : in out Sequence;
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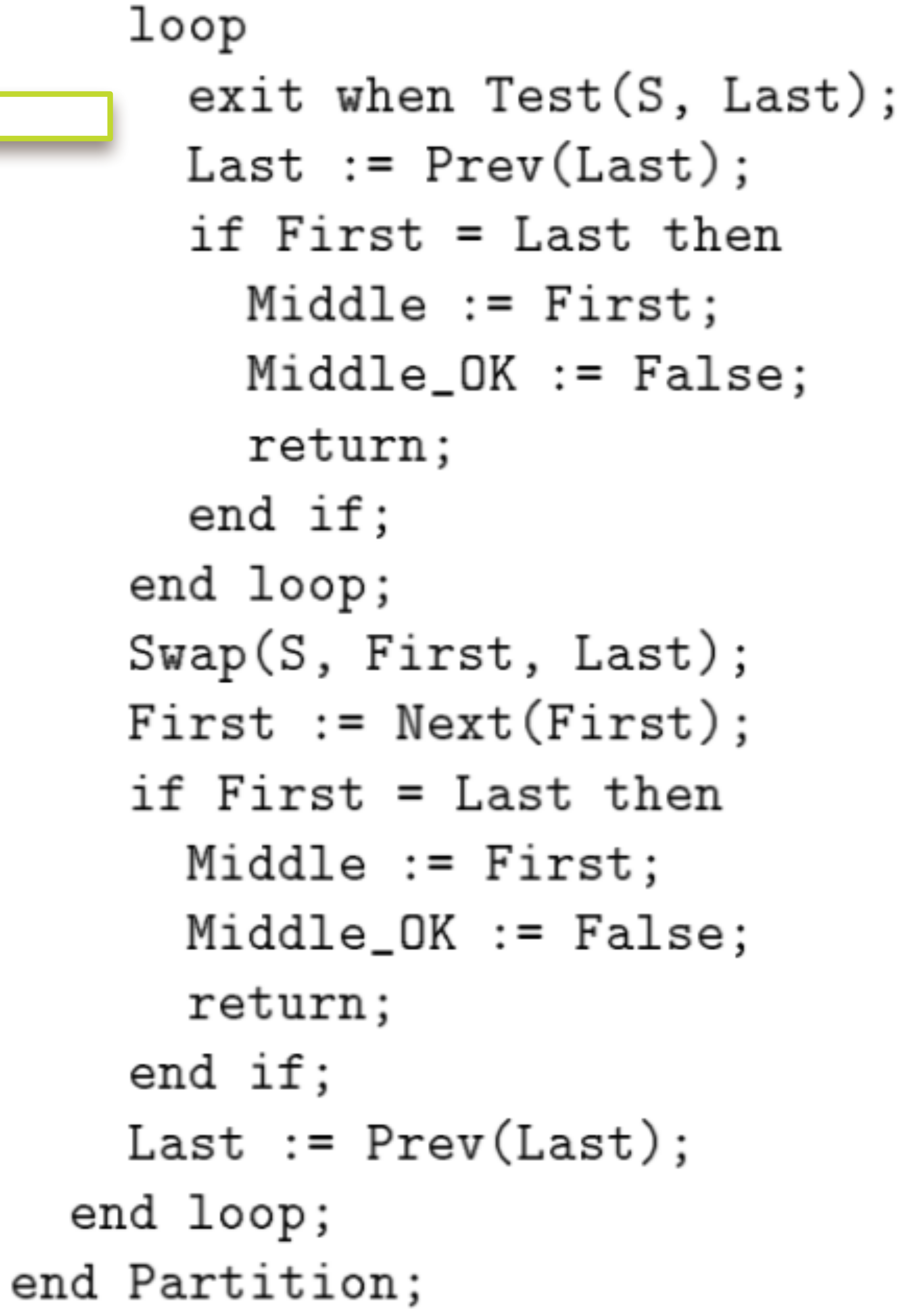
                    Middle     : out Coordinate;
                    Middle_OK  : out Boolean) is
  First : Coordinate := F;
  Last  : Coordinate := L;
begin
  loop
    loop
      if First = Last then
        Middle := First;
        Middle_OK := Test(S, First);
        return;
      end if;
      exit when not Test(S, First);
      First := Next(First);
    end loop;
```



```
    loop
      exit when Test(S, Last);
      Last := Prev(Last);
      if First = Last then
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        Middle_OK := False;
        return;
      end if;
    end loop;
    Swap(S, First, Last);
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    loop
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      Last := Prev(Last);
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      end if;
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      end if;
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      exit when Test(S, Last);
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      end if;
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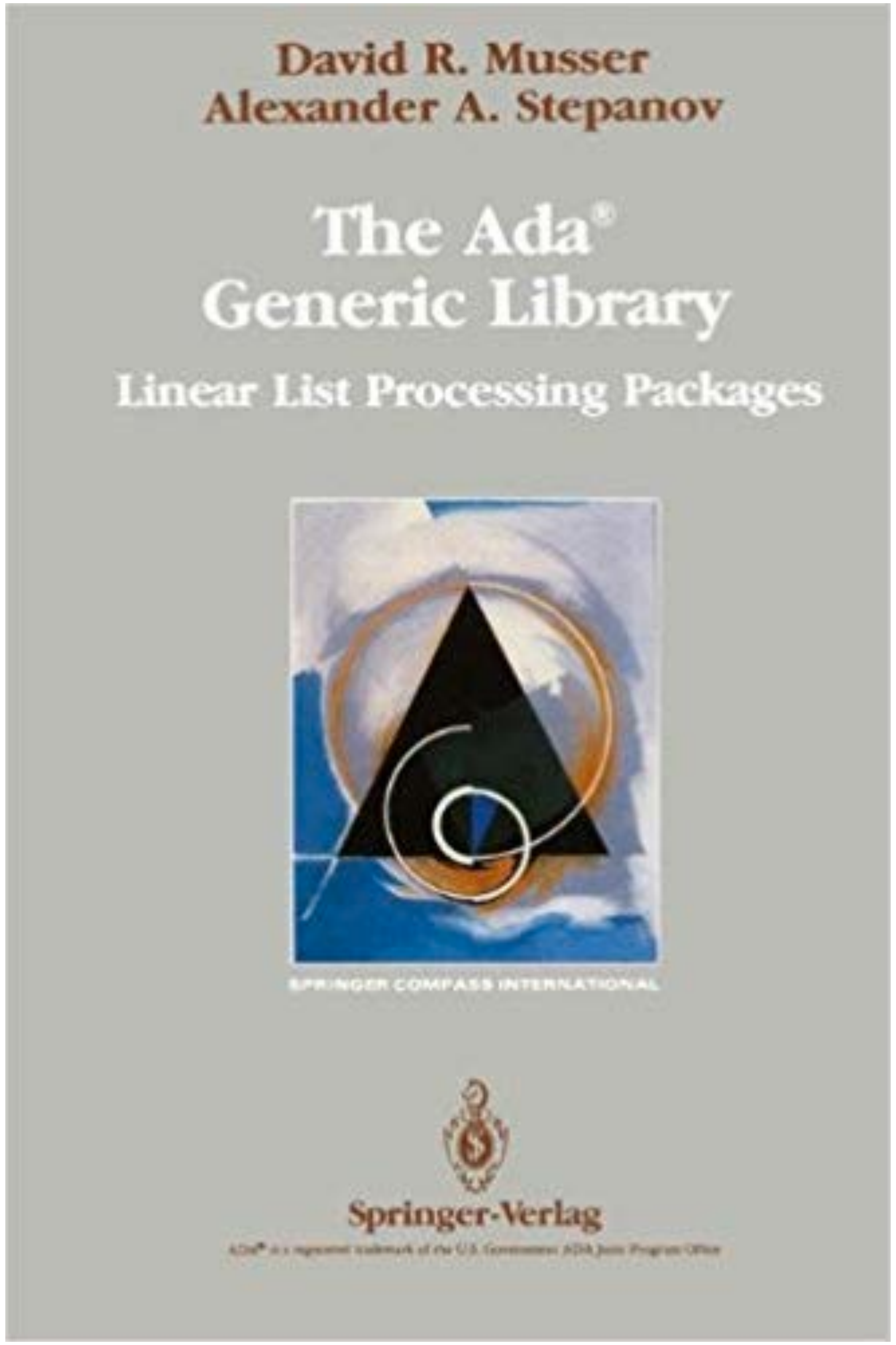

David R. Musser
Alexander A. Stepanov

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Generic Library
Linear List Processing Packages

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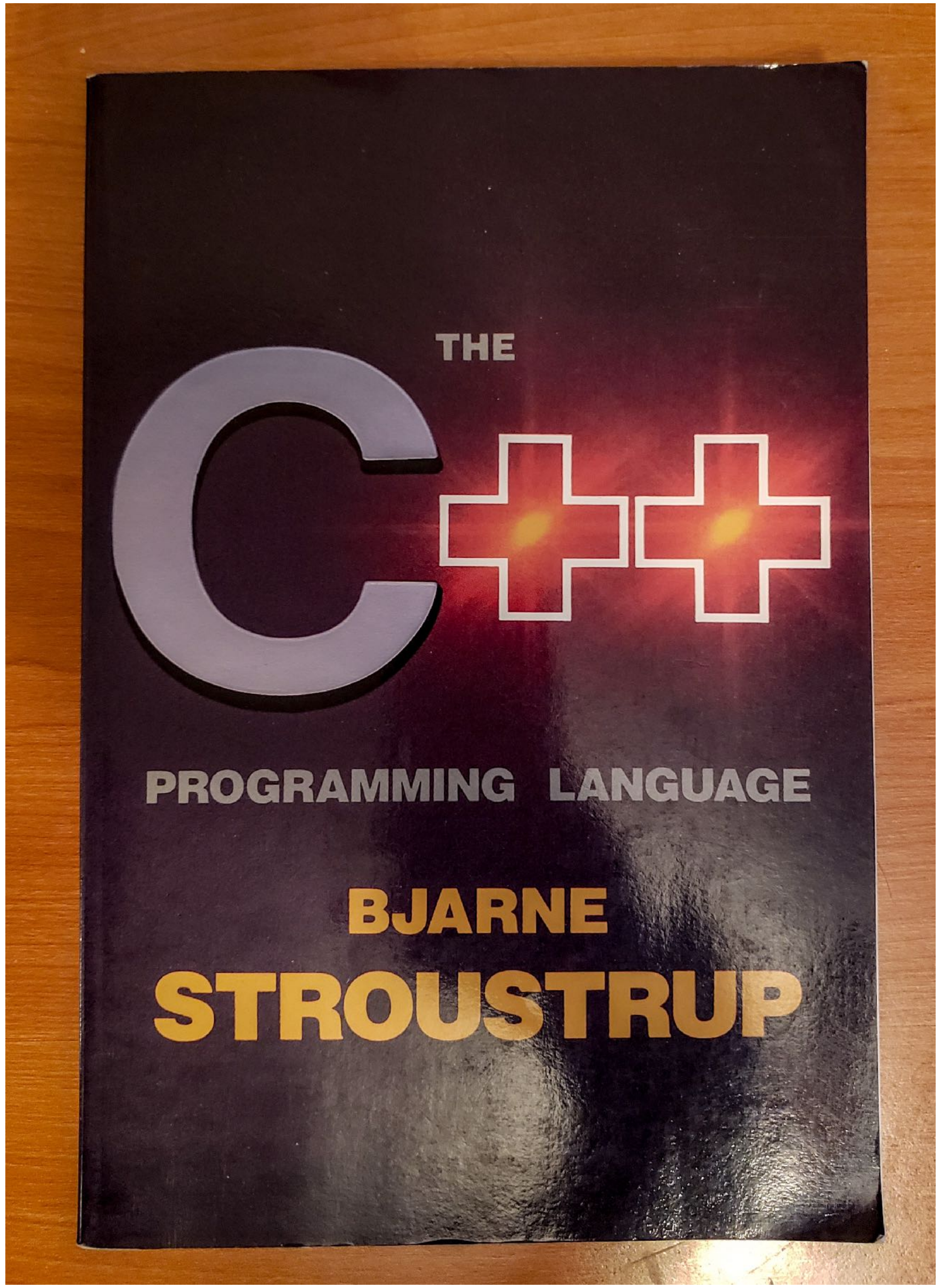
THE
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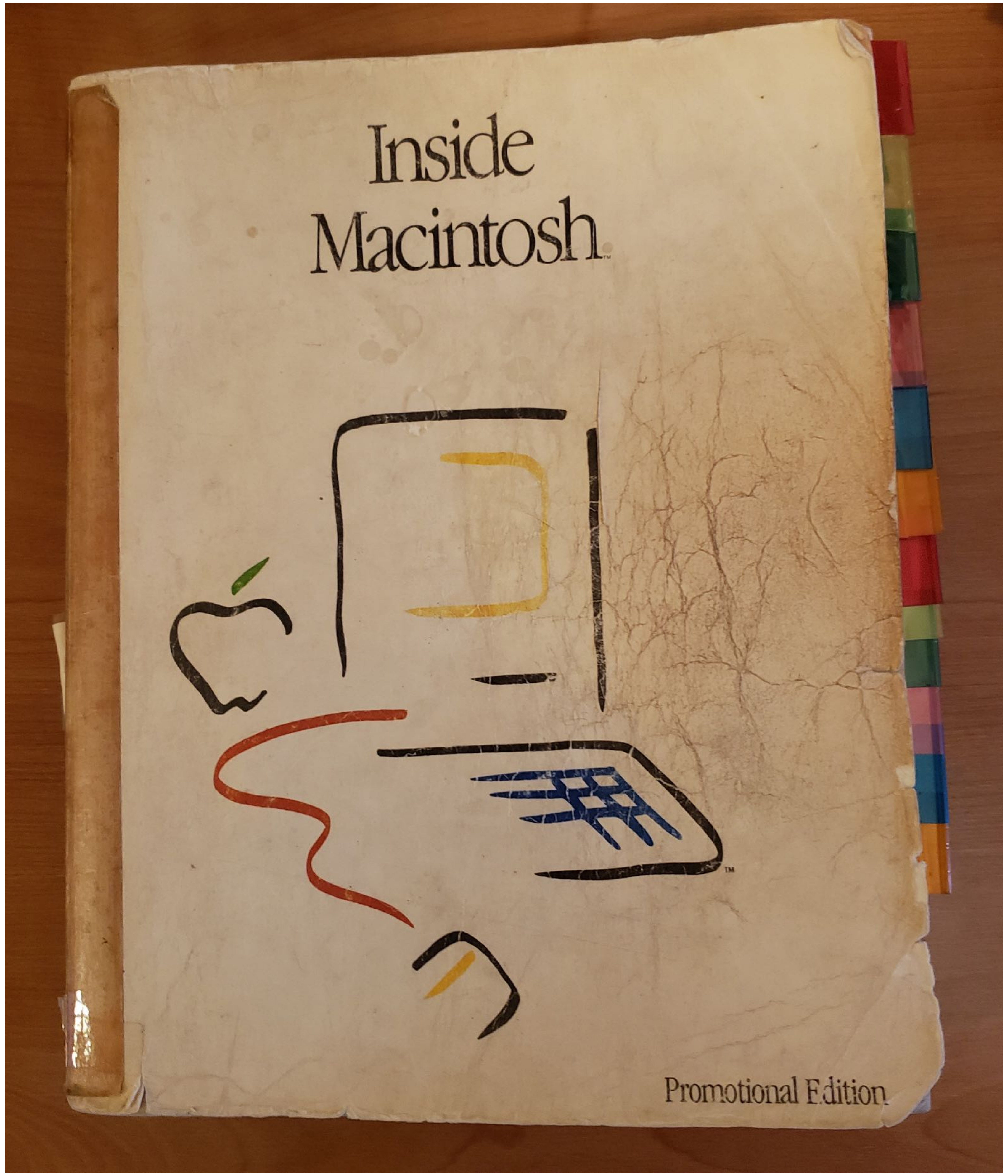
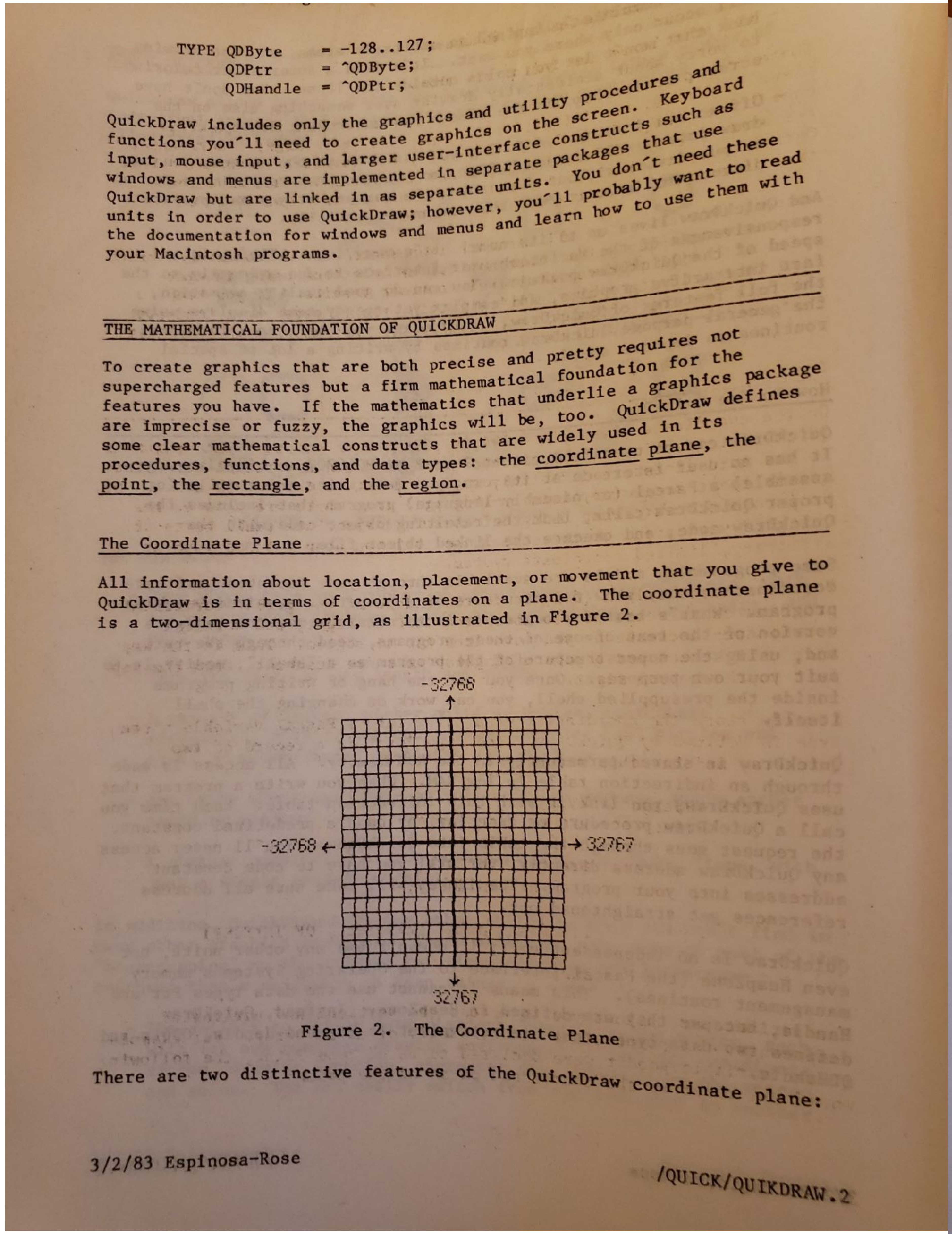
PROGRAMMING LANGUAGE
BJARNE
STROUSTRUP


Inside
Macintosh
Promotional Edition



```
TYPE QDByte    = -128..127;
     QDPtr     = ^QDByte;
     QDHandle  = ^QDPtr;
```

QuickDraw includes only the graphics and utility procedures and functions you´ll need to create graphics on the screen. Keyboard input, mouse input, and larger user-interface constructs such as windows and menus are implemented in separate packages that use QuickDraw but are linked in as separate units. You don´t need these units in order to use QuickDraw; however, you´ll probably want to read the documentation for windows and menus and learn how to use them with your Macintosh programs.

THE MATHEMATICAL FOUNDATION OF QUICKDRAW

To create graphics that are both precise and pretty requires not supercharged features but a firm mathematical foundation for the features you have. If the mathematics that underlie a graphics package are imprecise or fuzzy, the graphics will be, too. QuickDraw defines some clear mathematical constructs that are widely used in its procedures, functions, and data types: the coordinate plane, the point, the rectangle, and the region.

The Coordinate Plane

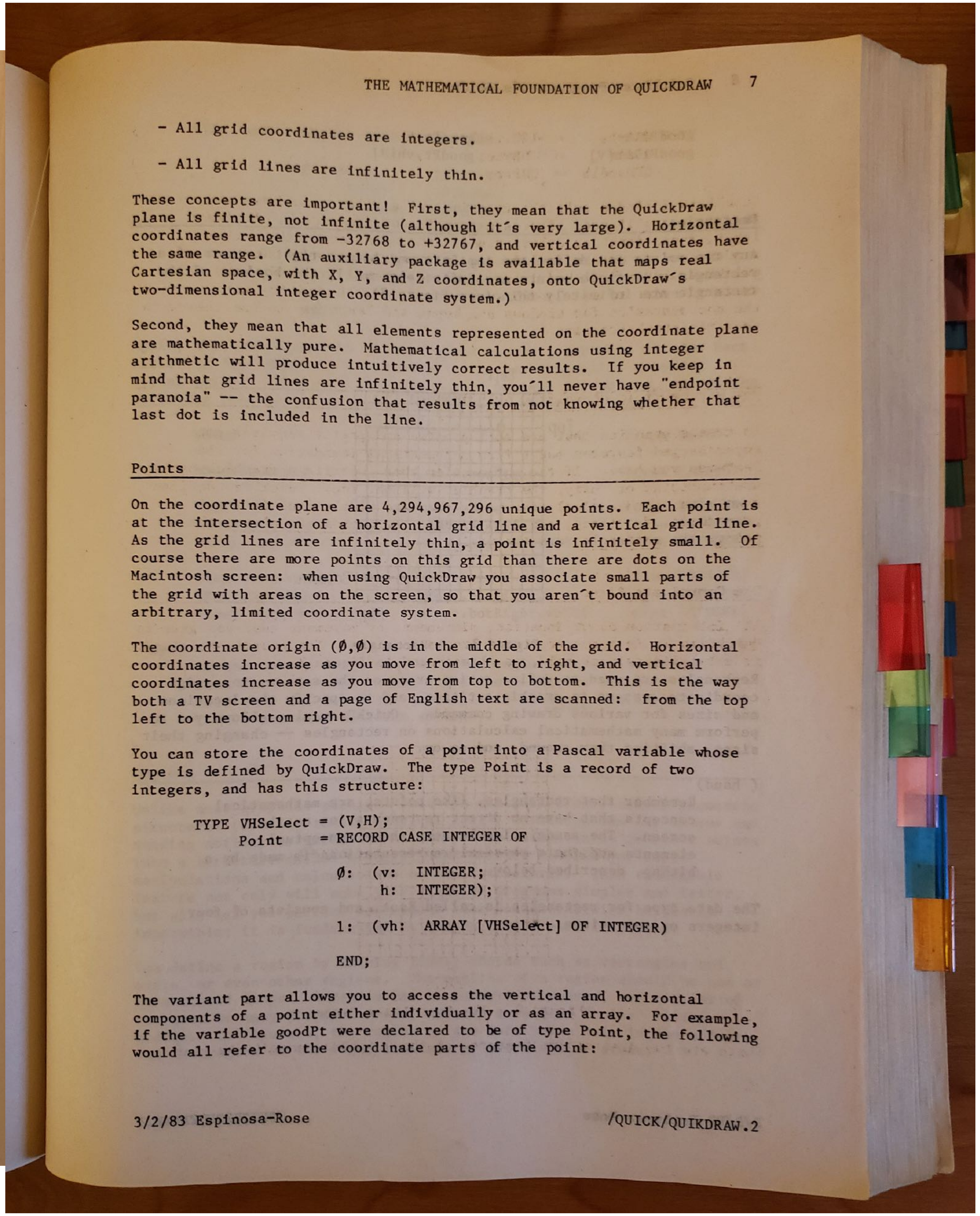
All information about location, placement, or movement that you give to QuickDraw is in terms of coordinates on a plane. The coordinate plane is a two-dimensional grid, as illustrated in Figure 2.

Figure 2. The Coordinate Plane

There are two distinctive features of the QuickDraw coordinate plane:

3/2/83 Espinosa-Rose

/QUICK/QUIKDRAW.2



THE MATHEMATICAL FOUNDATION OF QUICKDRAW 7

- All grid coordinates are integers.
- All grid lines are infinitely thin.

These concepts are important! First, they mean that the QuickDraw plane is finite, not infinite (although it´s very large). Horizontal coordinates range from -32768 to +32767, and vertical coordinates have the same range. (An auxiliary package is available that maps real Cartesian space, with X, Y, and Z coordinates, onto QuickDraw´s two-dimensional integer coordinate system.)

Second, they mean that all elements represented on the coordinate plane are mathematically pure. Mathematical calculations using integer arithmetic will produce intuitively correct results. If you keep in mind that grid lines are infinitely thin, you´ll never have "endpoint paranoia" -- the confusion that results from not knowing whether that last dot is included in the line.

Points

On the coordinate plane are 4,294,967,296 unique points. Each point is at the intersection of a horizontal grid line and a vertical grid line. As the grid lines are infinitely thin, a point is infinitely small. Of course there are more points on this grid than there are dots on the Macintosh screen: when using QuickDraw you associate small parts of the grid with areas on the screen, so that you aren´t bound into an arbitrary, limited coordinate system.

The coordinate origin (0,0) is in the middle of the grid. Horizontal coordinates increase as you move from left to right, and vertical coordinates increase as you move from top to bottom. This is the way both a TV screen and a page of English text are scanned: from the top left to the bottom right.

You can store the coordinates of a point into a Pascal variable whose type is defined by QuickDraw. The type Point is a record of two integers, and has this structure:

```
TYPE VHSelect = (V,H);
     Point    = RECORD CASE INTEGER OF

                  0:  (v:  INTEGER;
                       h:  INTEGER);

                  1:  (vh:  ARRAY [VHSelect] OF INTEGER)

                  END;
```

The variant part allows you to access the vertical and horizontal components of a point either individually or as an array. For example, if the variable goodPt were declared to be of type Point, the following would all refer to the coordinate parts of the point:

3/2/83 Espinosa-Rose

/QUICK/QUIKDRAW.2

```
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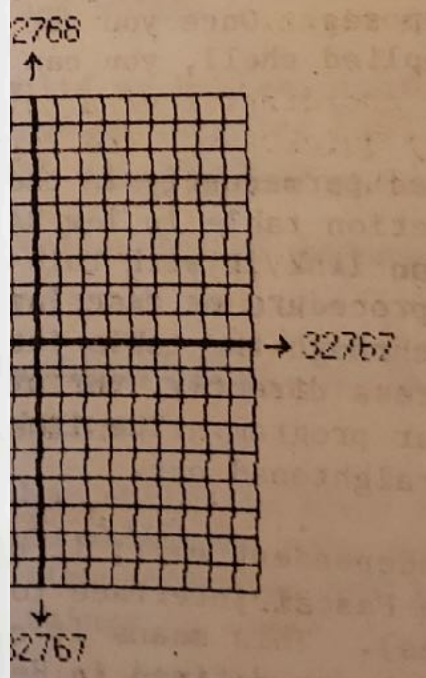
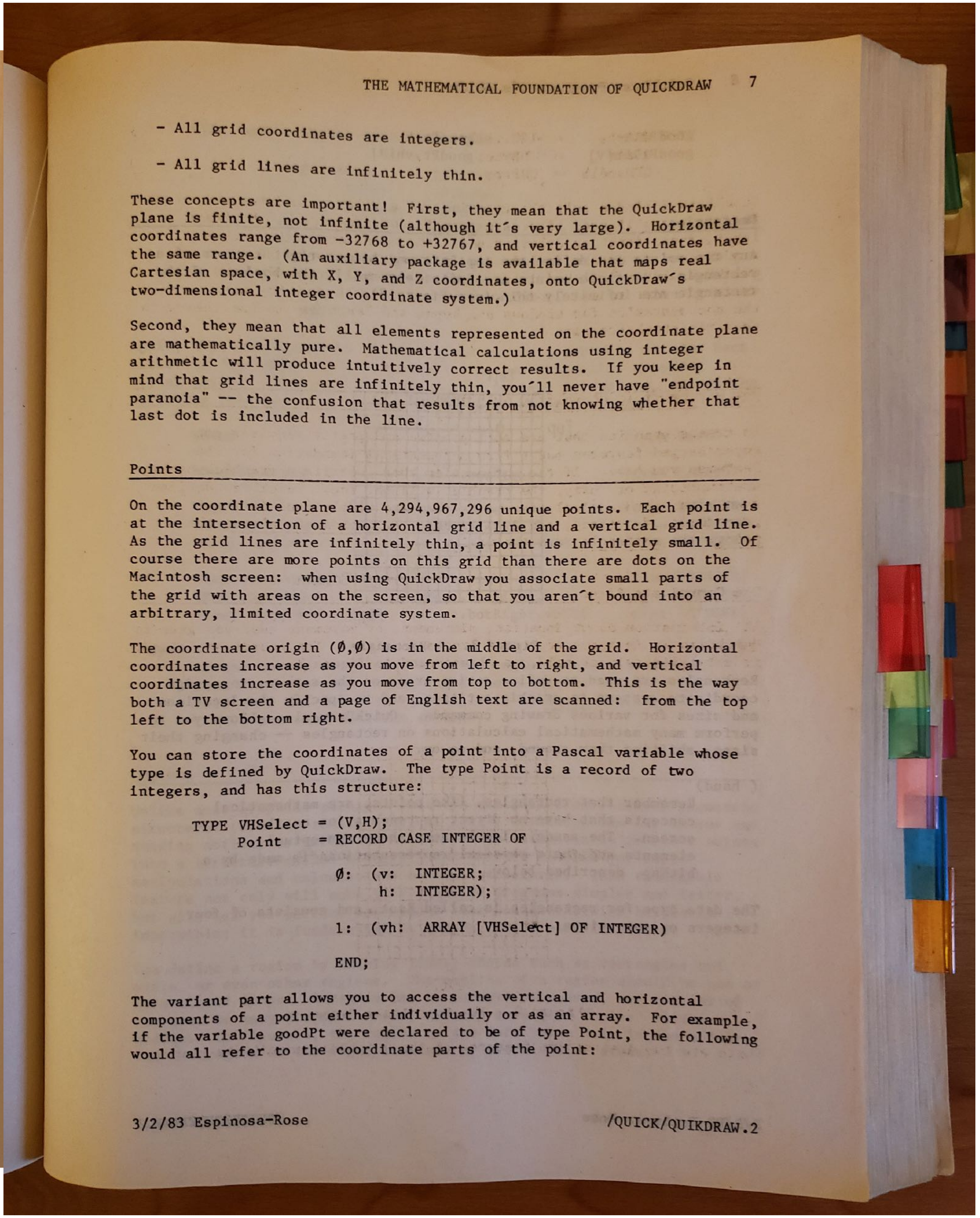


e Coordinate Plane

s of the QuickDraw coordinate plane:



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                     h:  INTEGER);

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                END;
```

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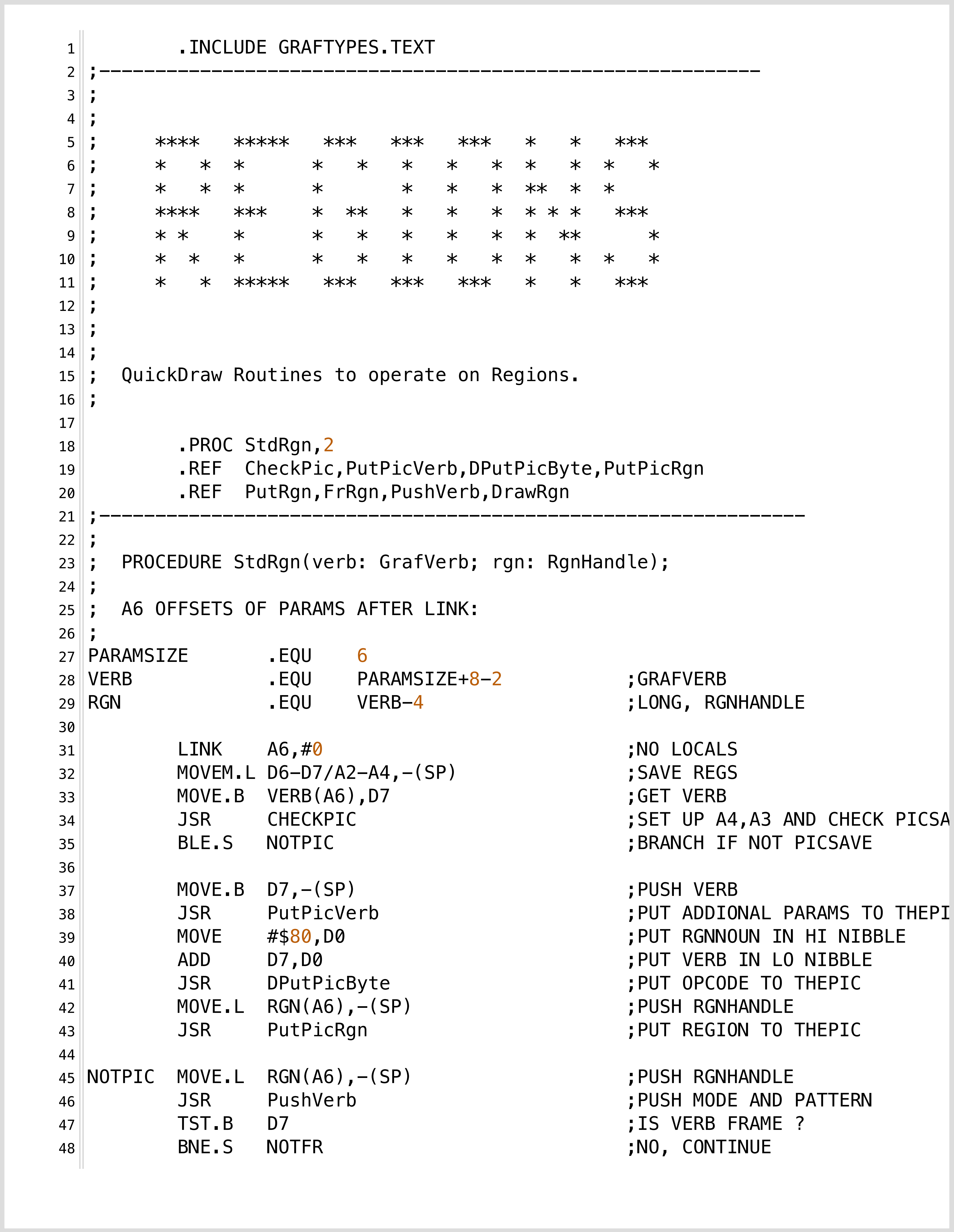


```
        .INCLUDE GRAFTYPES.TEXT
;-----------------------------------------------------------
;
;
;     ****   *****   ***   ***   ***   *   *   ***
;     *   *  *      *   *   *   *   *  *   *  *   *
;     *   *  *      *       *   *   *  **  *  *
;     ****   ***    *  **   *   *   *  * * *   ***
;     * *    *      *   *   *   *   *  *  **      *
;     *  *   *      *   *   *   *   *  *   *  *   *
;     *   *  *****   ***   ***   ***   *   *   ***
;
;
;
;  QuickDraw Routines to operate on Regions.
;

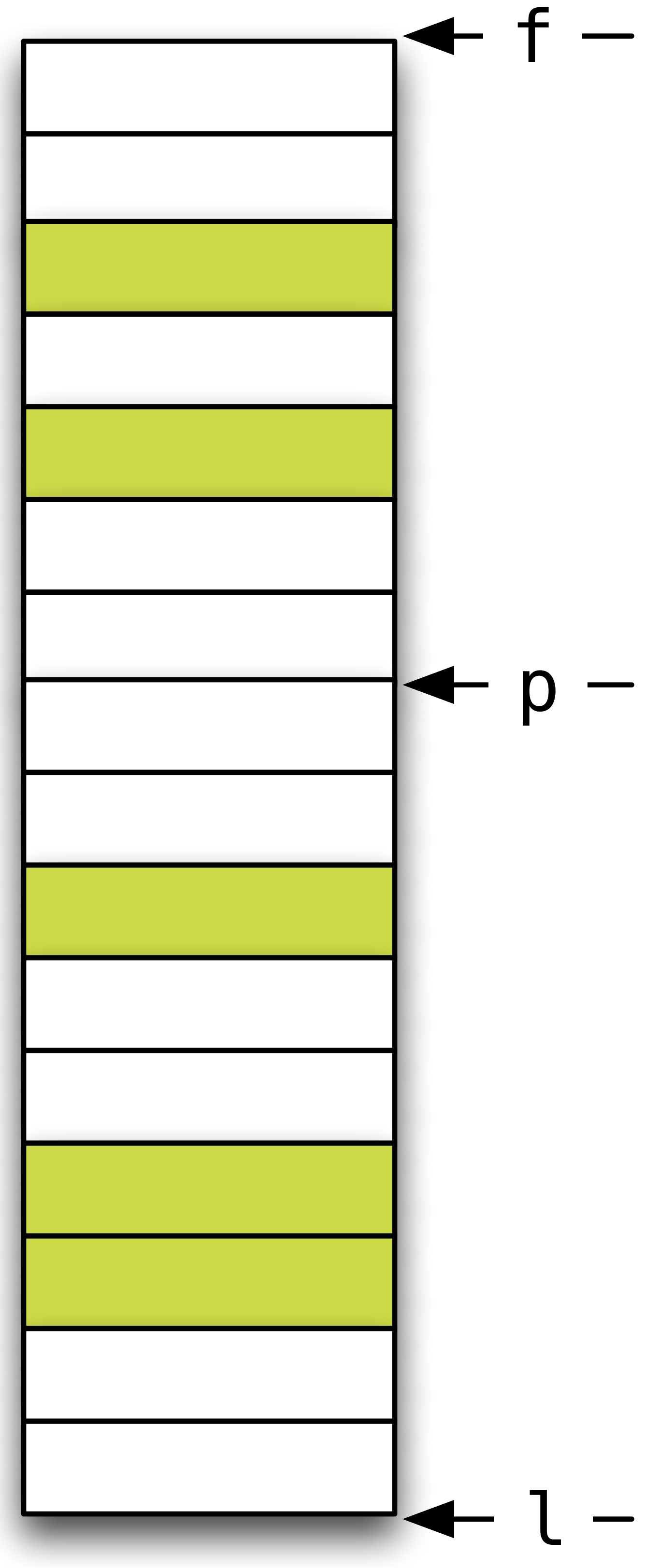
        .PROC StdRgn,2
        .REF  CheckPic,PutPicVerb,DPutPicByte,PutPicRgn
        .REF  PutRgn,FrRgn,PushVerb,DrawRgn
;---------------------------------------------------------------
;
;  PROCEDURE StdRgn(verb: GrafVerb; rgn: RgnHandle);
;
;  A6 OFFSETS OF PARAMS AFTER LINK:
;
PARAMSIZE       .EQU    6
VERB            .EQU    PARAMSIZE+8-2           ;GRAFVERB
RGN             .EQU    VERB-4                  ;LONG, RGNHANDLE

        LINK    A6,#0                           ;NO LOCALS
        MOVEM.L D6-D7/A2-A4,-(SP)               ;SAVE REGS
        MOVE.B  VERB(A6),D7                     ;GET VERB
        JSR     CHECKPIC                        ;SET UP A4,A3 AND CHECK PICSA
        BLE.S   NOTPIC                          ;BRANCH IF NOT PICSAVE

        MOVE.B  D7,-(SP)                        ;PUSH VERB
        JSR     PutPicVerb                      ;PUT ADDIONAL PARAMS TO THEPI
        MOVE    #$80,D0                         ;PUT RGNNOUN IN HI NIBBLE
        ADD     D7,D0                           ;PUT VERB IN LO NIBBLE
        JSR     DPutPicByte                     ;PUT OPCODE TO THEPIC
        MOVE.L  RGN(A6),-(SP)                   ;PUSH RGNHANDLE
        JSR     PutPicRgn                       ;PUT REGION TO THEPIC

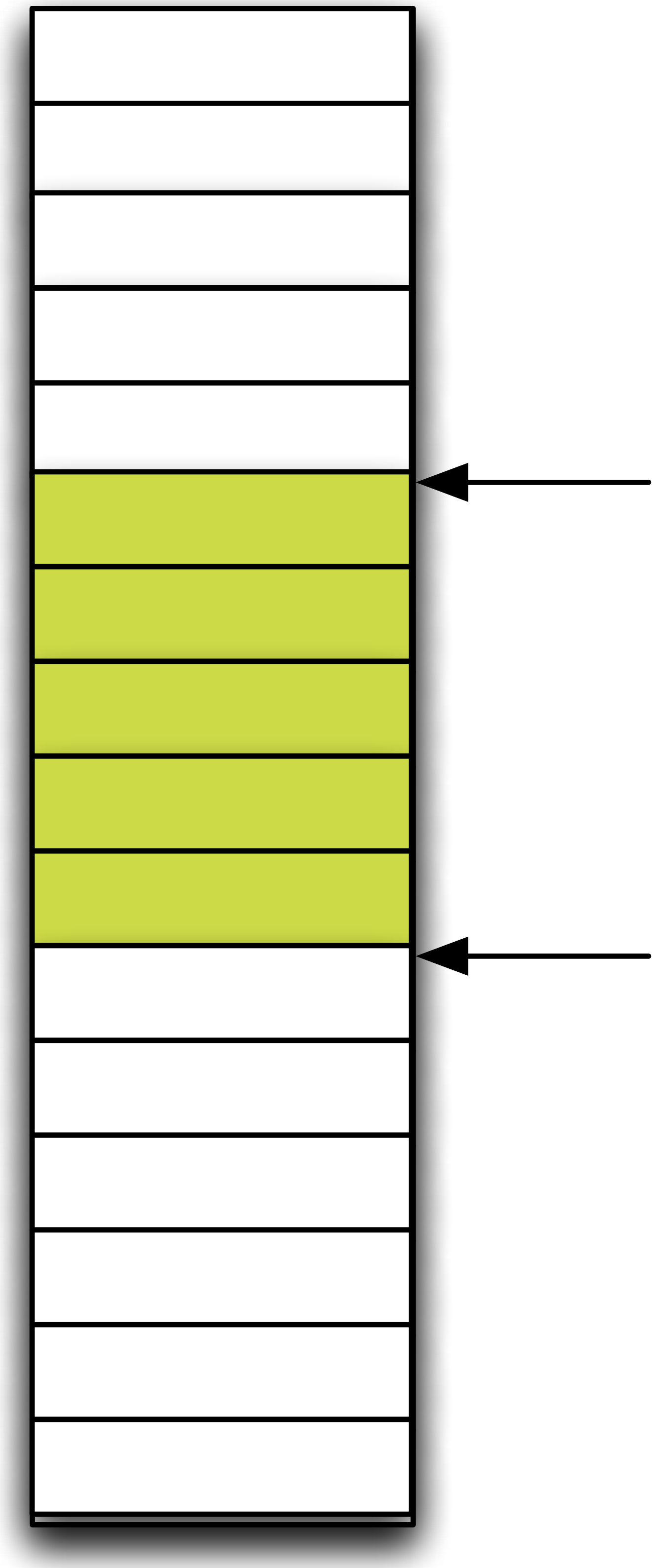
NOTPIC  MOVE.L  RGN(A6),-(SP)                   ;PUSH RGNHANDLE
        JSR     PushVerb                        ;PUSH MODE AND PATTERN
        TST.B   D7                              ;IS VERB FRAME ?
        BNE.S   NOTFR                           ;NO, CONTINUE
```

Gather



```
template <typename I, // I models BidirectionalIterator
          typename S> // S models UnaryPredicate
auto gather(I f, I l, I p, S s) -> pair<I, I>
{
    return { stable_partition(f, p, not1(s)),
             stable_partition(p, l, s) };
}
```

Gather



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             stable_partition(p, l, s) };
}
```

For a sequence of n elements there are $n + 1$ positions

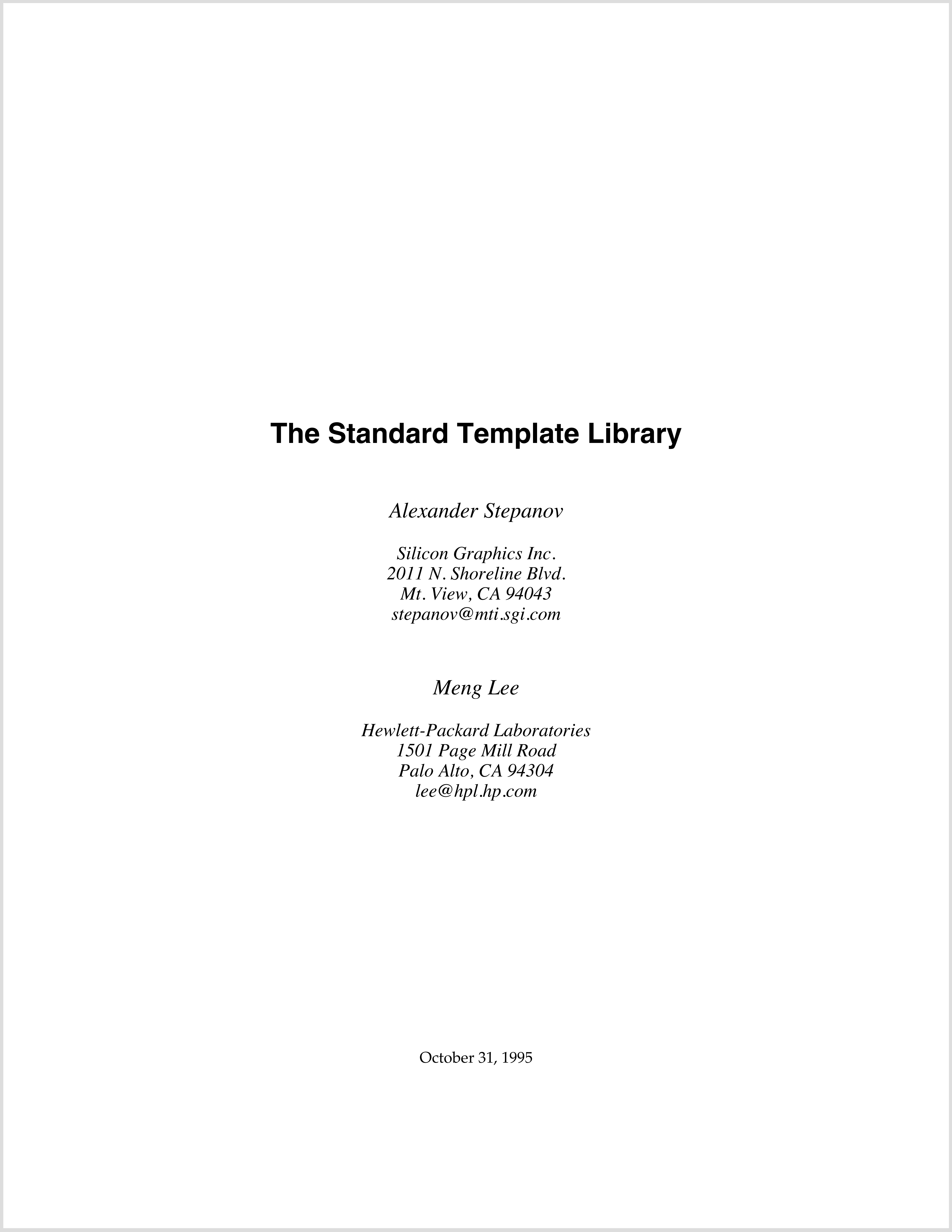
1993

1993

Alex resumes work on Generic Programming

Andrew Koenig suggests writing a standard library proposal

1994



The Standard Template Library

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October 31, 1995

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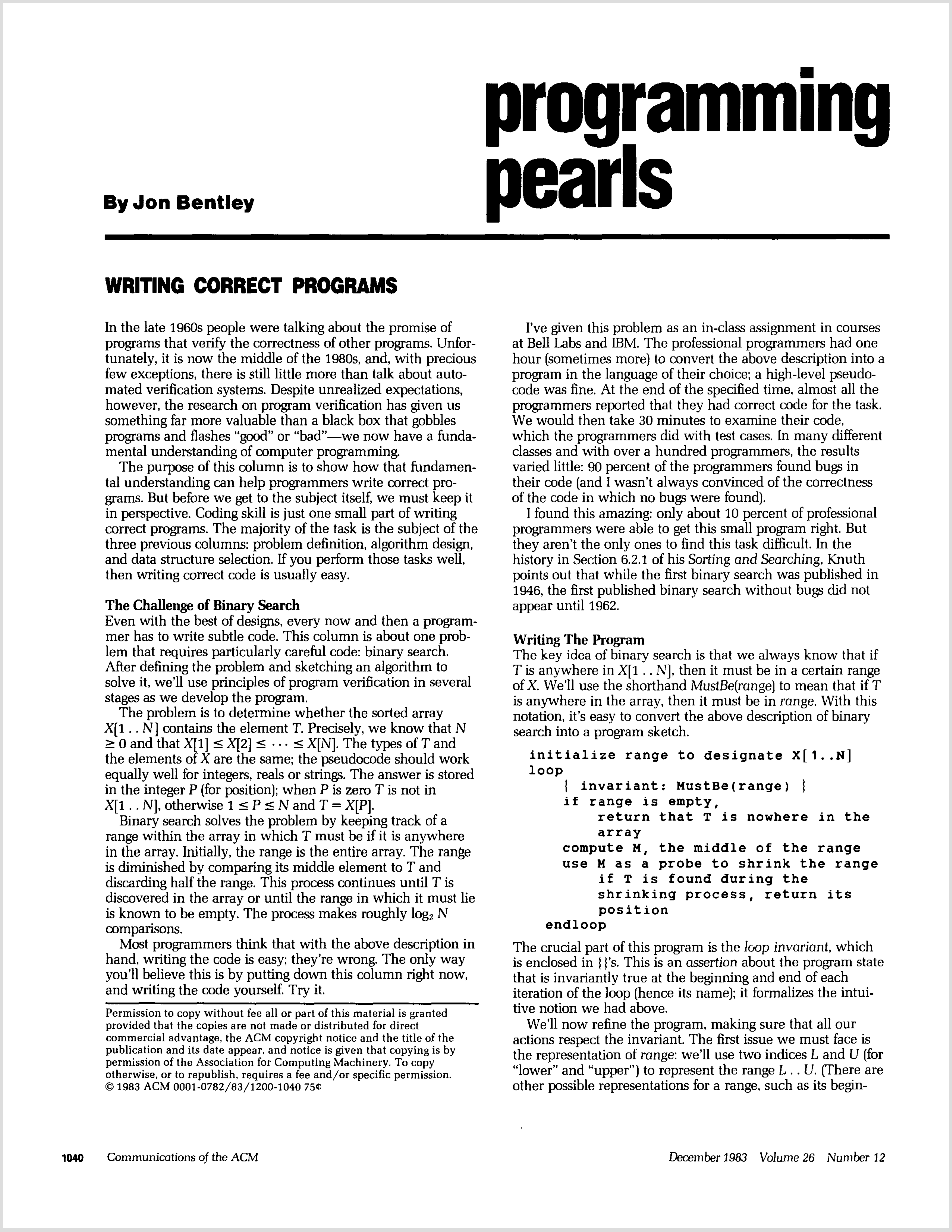
1983



programming pearls

By Jon Bentley

WRITING CORRECT PROGRAMS

In the late 1960s people were talking about the promise of programs that verify the correctness of other programs. Unfortunately, it is now the middle of the 1980s, and, with precious few exceptions, there is still little more than talk about automated verification systems. Despite unrealized expectations, however, the research on program verification has given us something far more valuable than a black box that gobbles programs and flashes "good" or "bad"—we now have a fundamental understanding of computer programming.

The purpose of this column is to show how that fundamental understanding can help programmers write correct programs. But before we get to the subject itself, we must keep it in perspective. Coding skill is just one small part of writing correct programs. The majority of the task is the subject of the three previous columns: problem definition, algorithm design, and data structure selection. If you perform those tasks well, then writing correct code is usually easy.

The Challenge of Binary Search

Even with the best of designs, every now and then a programmer has to write subtle code. This column is about one problem that requires particularly careful code: binary search. After defining the problem and sketching an algorithm to solve it, we'll use principles of program verification in several stages as we develop the program.

The problem is to determine whether the sorted array $X[1\,..\,N]$ contains the element T. Precisely, we know that $N \geq 0$ and that $X[1] \leq X[2] \leq \cdots \leq X[N]$. The types of T and the elements of X are the same; the pseudocode should work equally well for integers, reals or strings. The answer is stored in the integer P (for position); when P is zero T is not in $X[1\,..\,N]$, otherwise $1 \leq P \leq N$ and $T = X[P]$.

Binary search solves the problem by keeping track of a range within the array in which T must be if it is anywhere in the array. Initially, the range is the entire array. The range is diminished by comparing its middle element to T and discarding half the range. This process continues until T is discovered in the array or until the range in which it must lie is known to be empty. The process makes roughly $\log_2 N$ comparisons.

Most programmers think that with the above description in hand, writing the code is easy; they're wrong. The only way you'll believe this is by putting down this column right now, and writing the code yourself. Try it.

I've given this problem as an in-class assignment in courses at Bell Labs and IBM. The professional programmers had one hour (sometimes more) to convert the above description into a program in the language of their choice; a high-level pseudocode was fine. At the end of the specified time, almost all the programmers reported that they had correct code for the task. We would then take 30 minutes to examine their code, which the programmers did with test cases. In many different classes and with over a hundred programmers, the results varied little: 90 percent of the programmers found bugs in their code (and I wasn't always convinced of the correctness of the code in which no bugs were found).

I found this amazing: only about 10 percent of professional programmers were able to get this small program right. But they aren't the only ones to find this task difficult. In the history in Section 6.2.1 of his *Sorting and Searching*, Knuth points out that while the first binary search was published in 1946, the first published binary search without bugs did not appear until 1962.

Writing The Program

The key idea of binary search is that we always know that if T is anywhere in $X[1\,..\,N]$, then it must be in a certain range of X. We'll use the shorthand *MustBe*(*range*) to mean that if T is anywhere in the array, then it must be in *range*. With this notation, it's easy to convert the above description of binary search into a program sketch.

```
initialize range to designate X[1..N]
loop
     { invariant: MustBe(range) }
     if range is empty,
          return that T is nowhere in the
          array
     compute M, the middle of the range
     use M as a probe to shrink the range
          if T is found during the
          shrinking process, return its
          position
  endloop
```

The crucial part of this program is the *loop invariant*, which is enclosed in {}'s. This is an *assertion* about the program state that is invariantly true at the beginning and end of each iteration of the loop (hence its name); it formalizes the intuitive notion we had above.

We'll now refine the program, making sure that all our actions respect the invariant. The first issue we must face is the representation of *range*: we'll use two indices L and U (for "lower" and "upper") to represent the range $L\,..\,U$. (There are other possible representations for a range, such as its begin-

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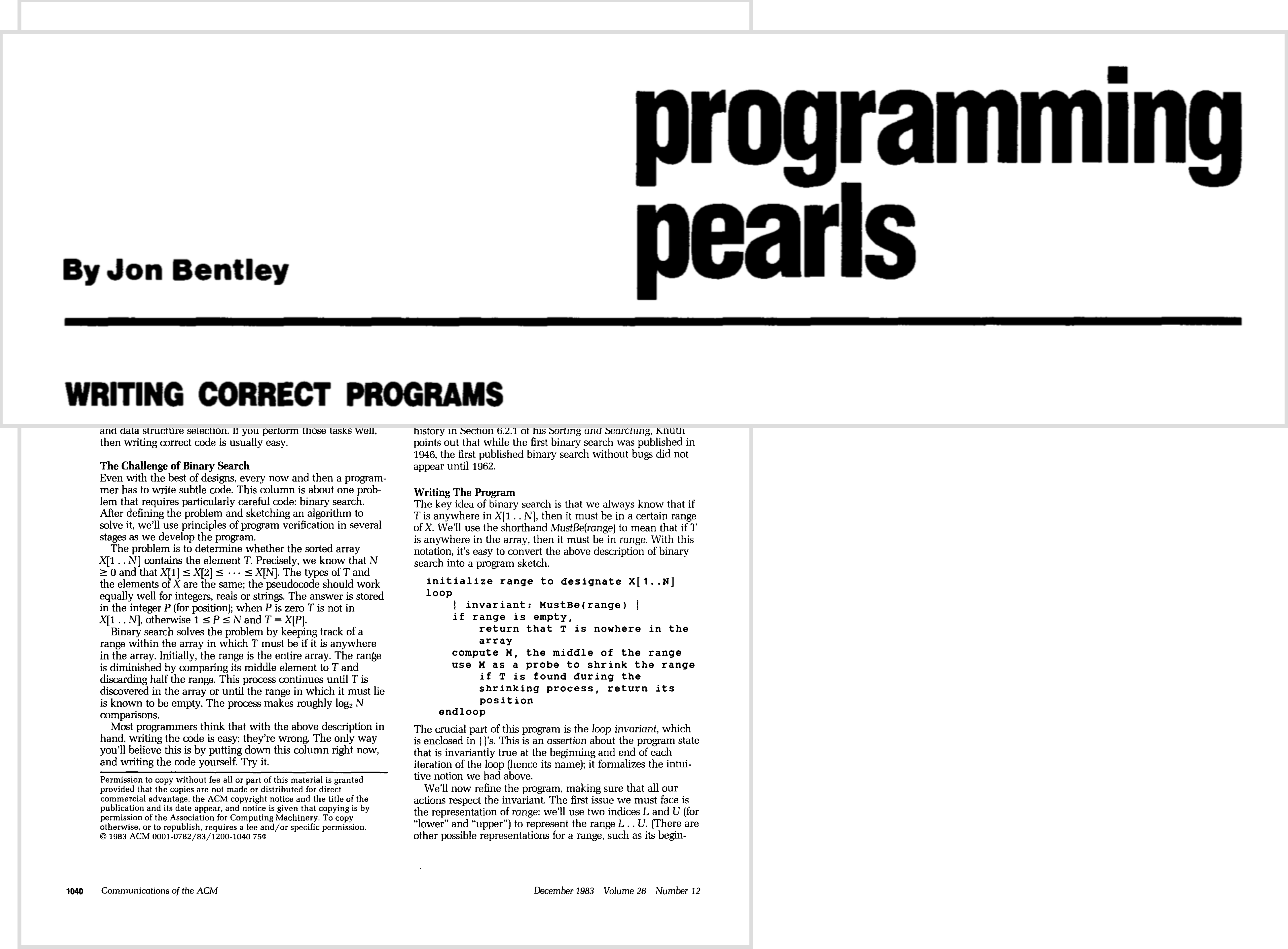

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>
> – Jon Bentley, Programming Pearls

“I want to hire the other ten percent.”
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Jon Bentley’s Solution (translated to C++)

```
int binary_search(int x[], int n, int v) {
    int l = 0;
    int u = n - 1;

    while (true) {
        if (l > u) return -1;

        int m = (l + u) / 2;

        if (x[m] < v) l = m + 1;
        else if (x[m] == v) return m;
        else /* (x[m] > v) */ u = m - 1;
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STL implementation

```
template <class I, // I models ForwardIterator
          class T> // T is value_type(I)
I lower_bound(I f, I l, const T& v) {
    while (f != l) {
        auto m = next(f, distance(f, l) / 2);

        if (*m < v) f = next(m);
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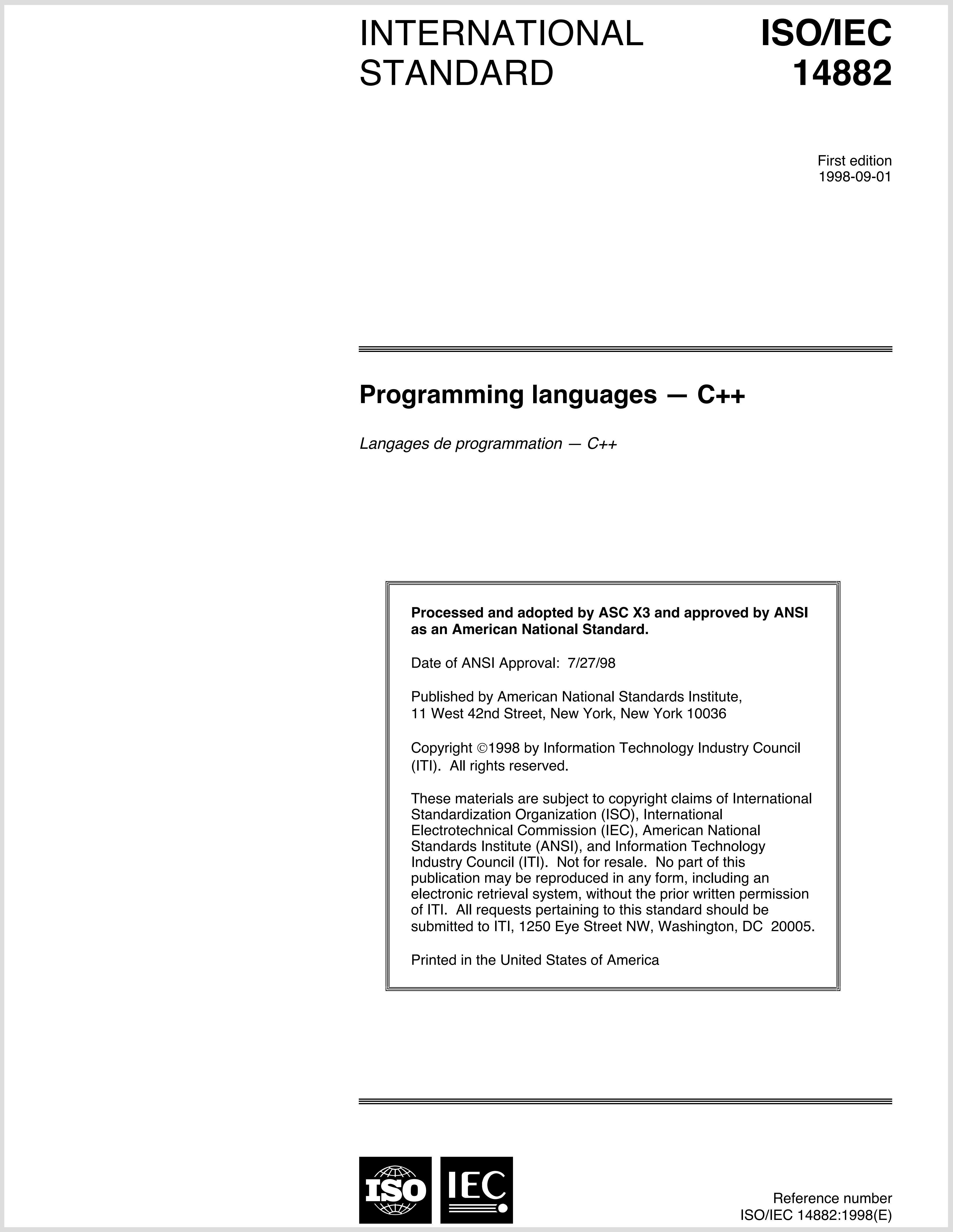

INTERNATIONAL STANDARD

ISO/IEC 14882

First edition
1998-09-01

Programming languages — C++

Langages de programmation — C++

Processed and adopted by ASC X3 and approved by ANSI as an American National Standard.

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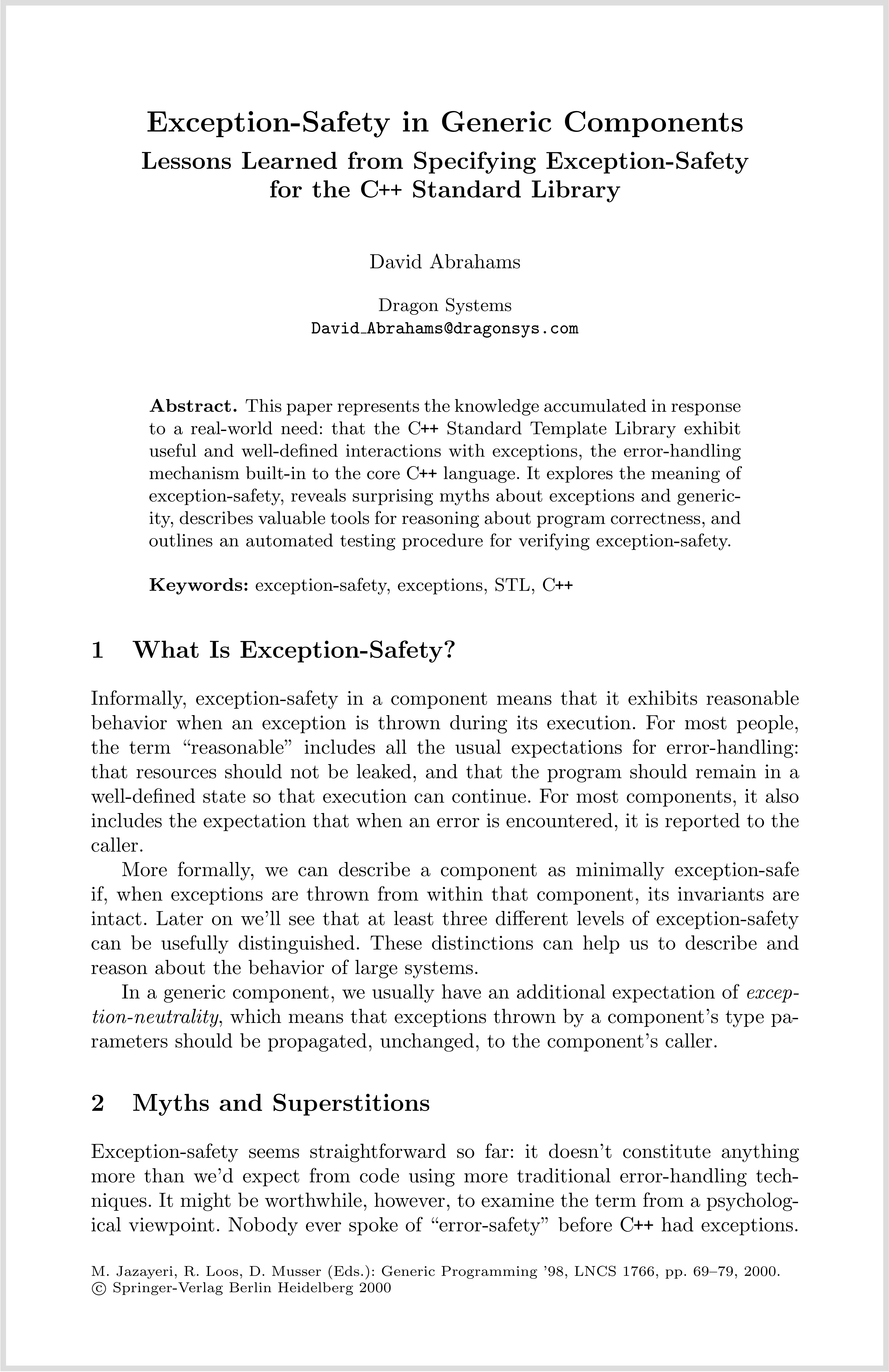


Exception-Safety in Generic Components

Lessons Learned from Specifying Exception-Safety for the C++ Standard Library

David Abrahams

Dragon Systems
David_Abrahams@dragonsys.com

Abstract. This paper represents the knowledge accumulated in response to a real-world need: that the C++ Standard Template Library exhibit useful and well-defined interactions with exceptions, the error-handling mechanism built-in to the core C++ language. It explores the meaning of exception-safety, reveals surprising myths about exceptions and genericity, describes valuable tools for reasoning about program correctness, and outlines an automated testing procedure for verifying exception-safety.

Keywords: exception-safety, exceptions, STL, C++

1 What Is Exception-Safety?

Informally, exception-safety in a component means that it exhibits reasonable behavior when an exception is thrown during its execution. For most people, the term “reasonable” includes all the usual expectations for error-handling: that resources should not be leaked, and that the program should remain in a well-defined state so that execution can continue. For most components, it also includes the expectation that when an error is encountered, it is reported to the caller.

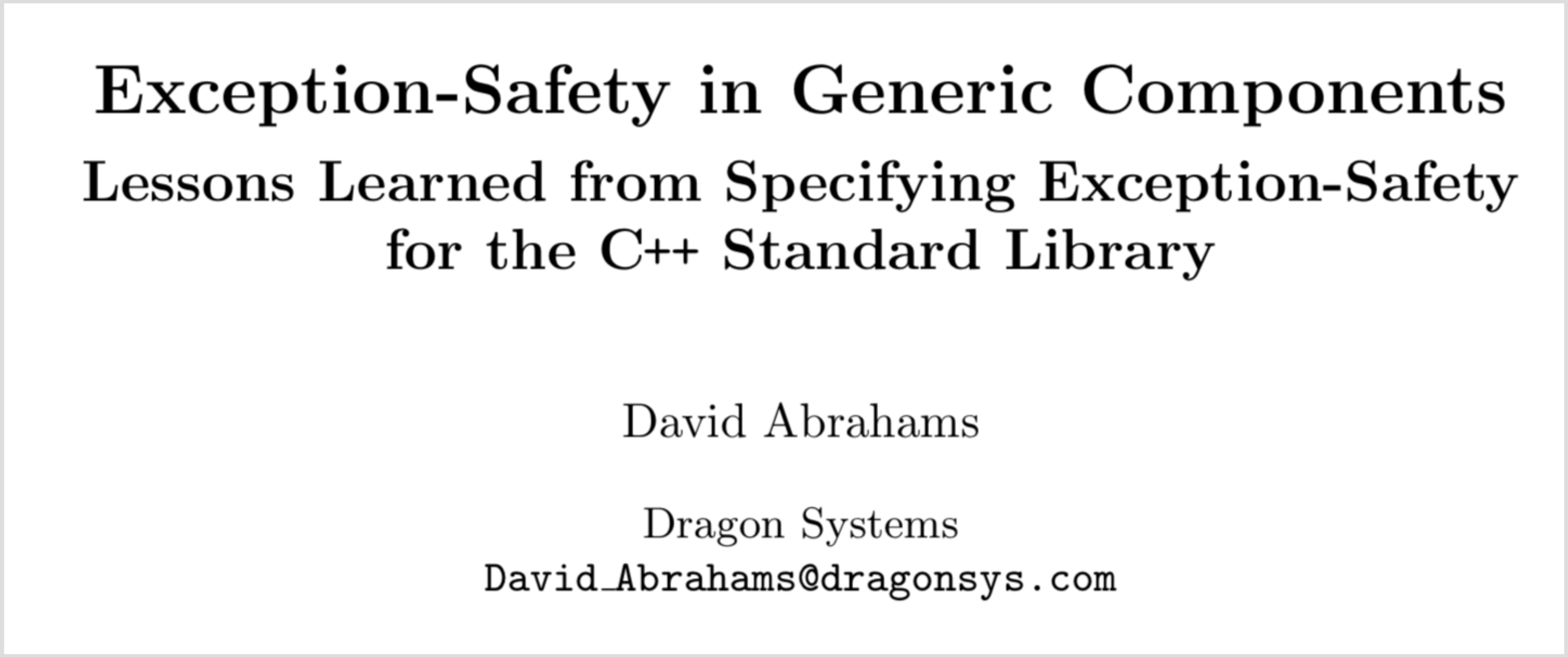
More formally, we can describe a component as minimally exception-safe if, when exceptions are thrown from within that component, its invariants are intact. Later on we’ll see that at least three different levels of exception-safety can be usefully distinguished. These distinctions can help us to describe and reason about the behavior of large systems.

In a generic component, we usually have an additional expectation of *exception-neutrality*, which means that exceptions thrown by a component’s type parameters should be propagated, unchanged, to the component’s caller.

2 Myths and Superstitions

Exception-safety seems straightforward so far: it doesn’t constitute anything more than we’d expect from code using more traditional error-handling techniques. It might be worthwhile, however, to examine the term from a psychological viewpoint. Nobody ever spoke of “error-safety” before C++ had exceptions.

M. Jazayeri, R. Loos, D. Musser (Eds.): Generic Programming ’98, LNCS 1766, pp. 69–79, 2000.



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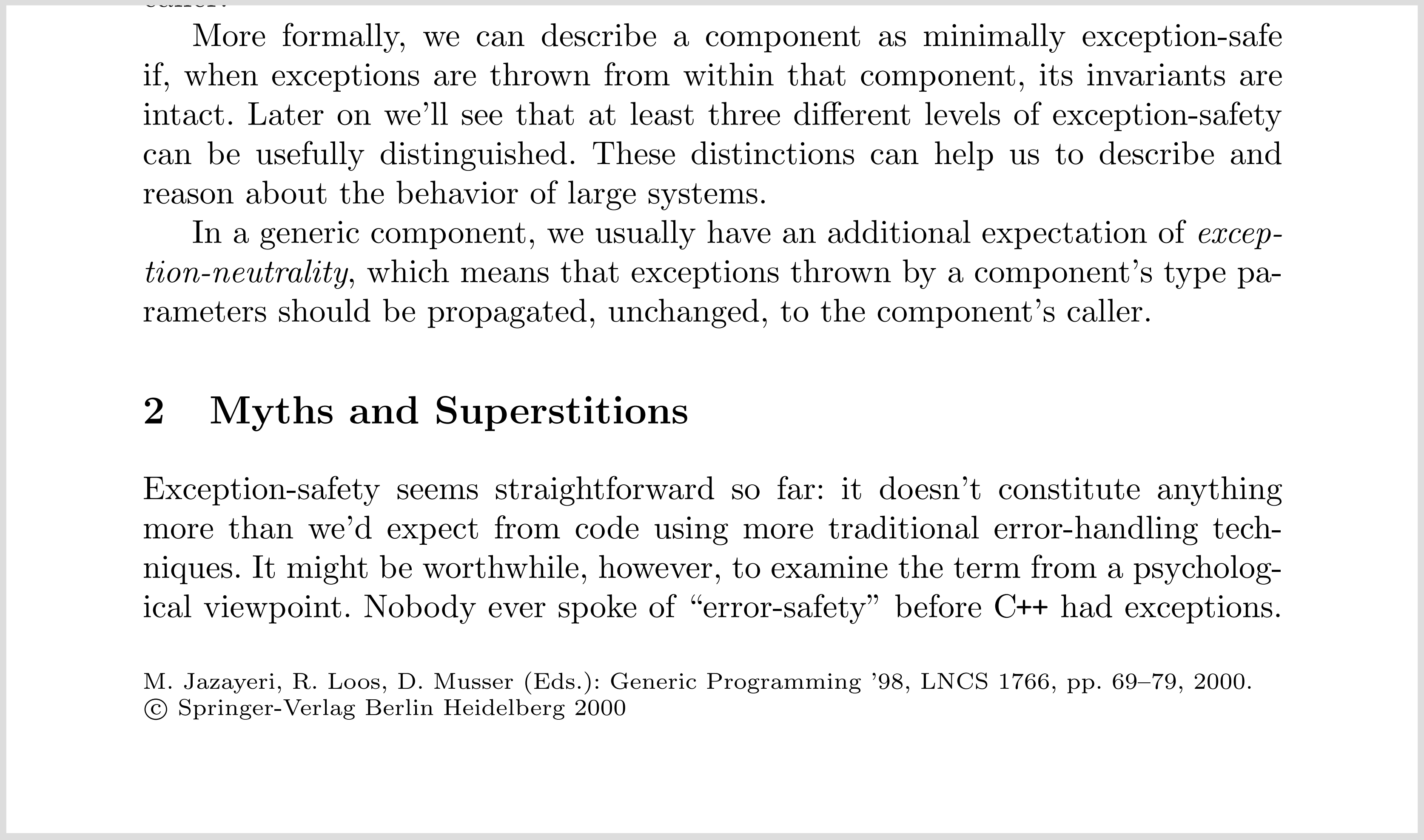

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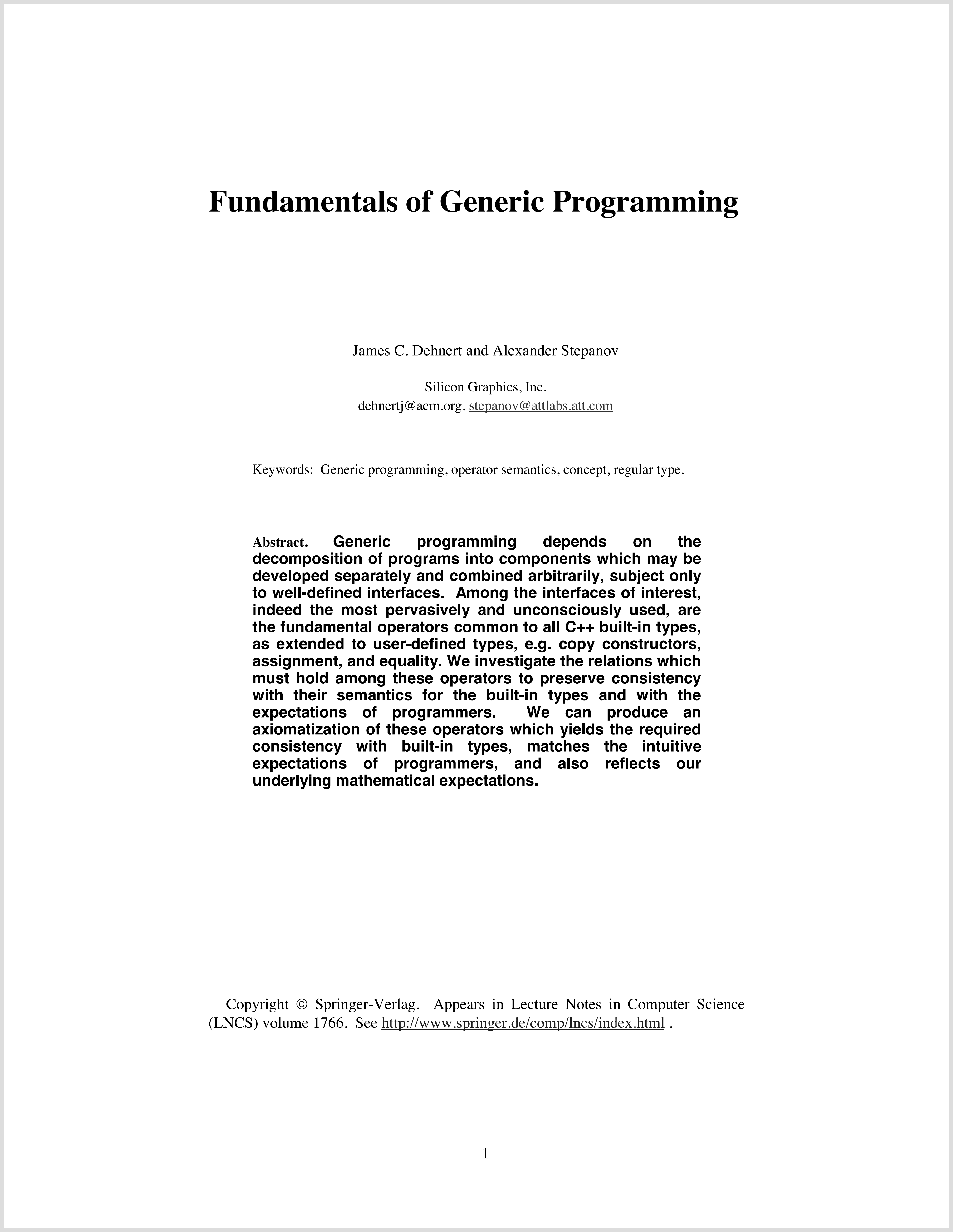


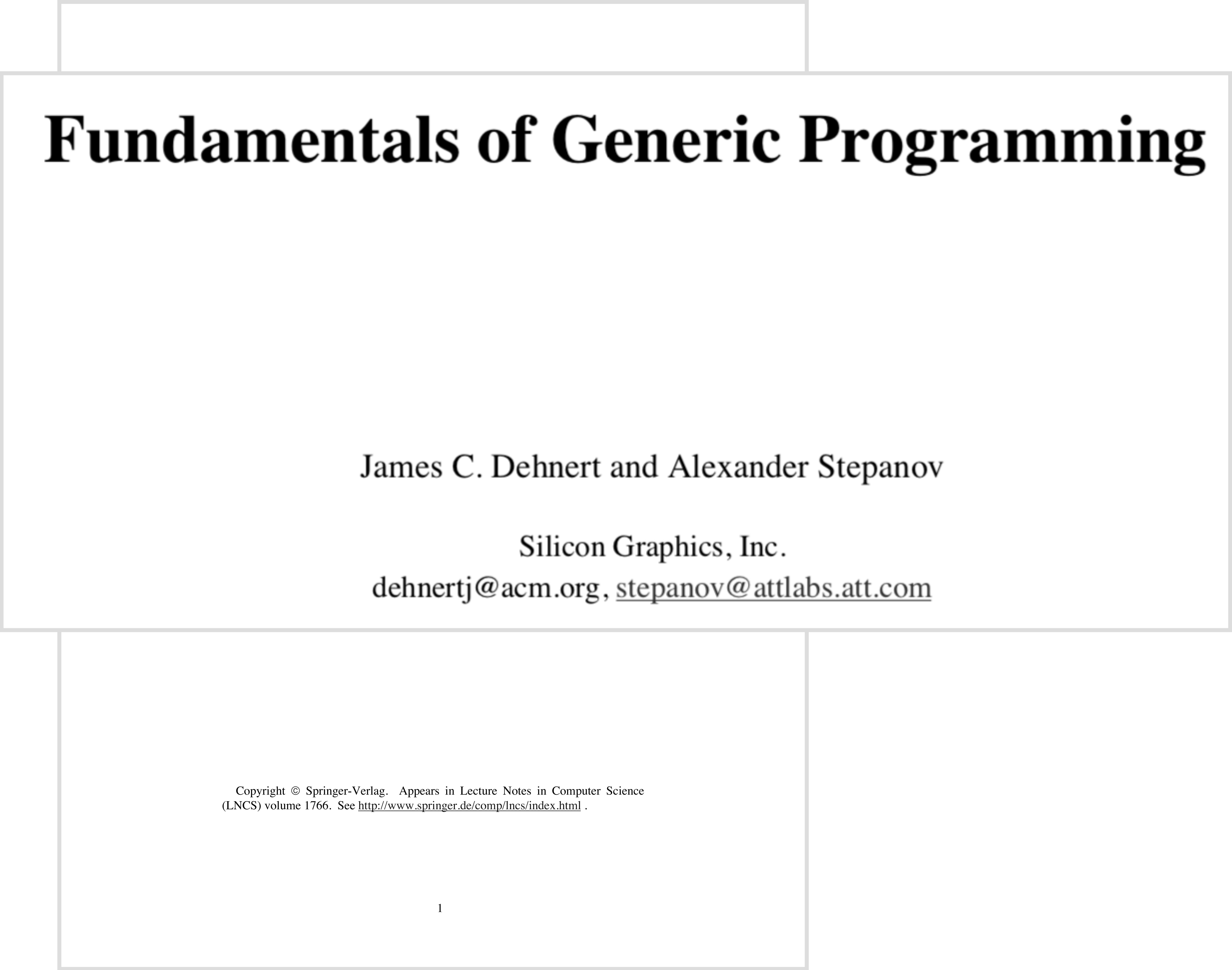
Fundamentals of Generic Programming

James C. Dehnert and Alexander Stepanov

Silicon Graphics, Inc.
dehnertj@acm.org, stepanov@attlabs.att.com

Keywords: Generic programming, operator semantics, concept, regular type.

Abstract. Generic programming depends on the decomposition of programs into components which may be developed separately and combined arbitrarily, subject only to well-defined interfaces. Among the interfaces of interest, indeed the most pervasively and unconsciously used, are the fundamental operators common to all C++ built-in types, as extended to user-defined types, e.g. copy constructors, assignment, and equality. We investigate the relations which must hold among these operators to preserve consistency with their semantics for the built-in types and with the expectations of programmers. We can produce an axiomatization of these operators which yields the required consistency with built-in types, matches the intuitive expectations of programmers, and also reflects our underlying mathematical expectations.



Fundamentals of Generic Programming

James C. Dehnert and Alexander Stepanov

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1

“We call the set of axioms satisfied by a data type and a set of operations on it a *concept*.”

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"Since we wish to extend semantics as well as syntax from built-in types to user types, we introduce the idea of a *regular type*, which matches the built-in type semantics, thereby making our user-defined types behave like built-in types as well."

"Since we wish to extend semantics as well as syntax from built-in types to user types, we introduce the idea of a ***regular type***, which matches the built-in type semantics, thereby making our user-defined types behave like built-in types as well."

2002

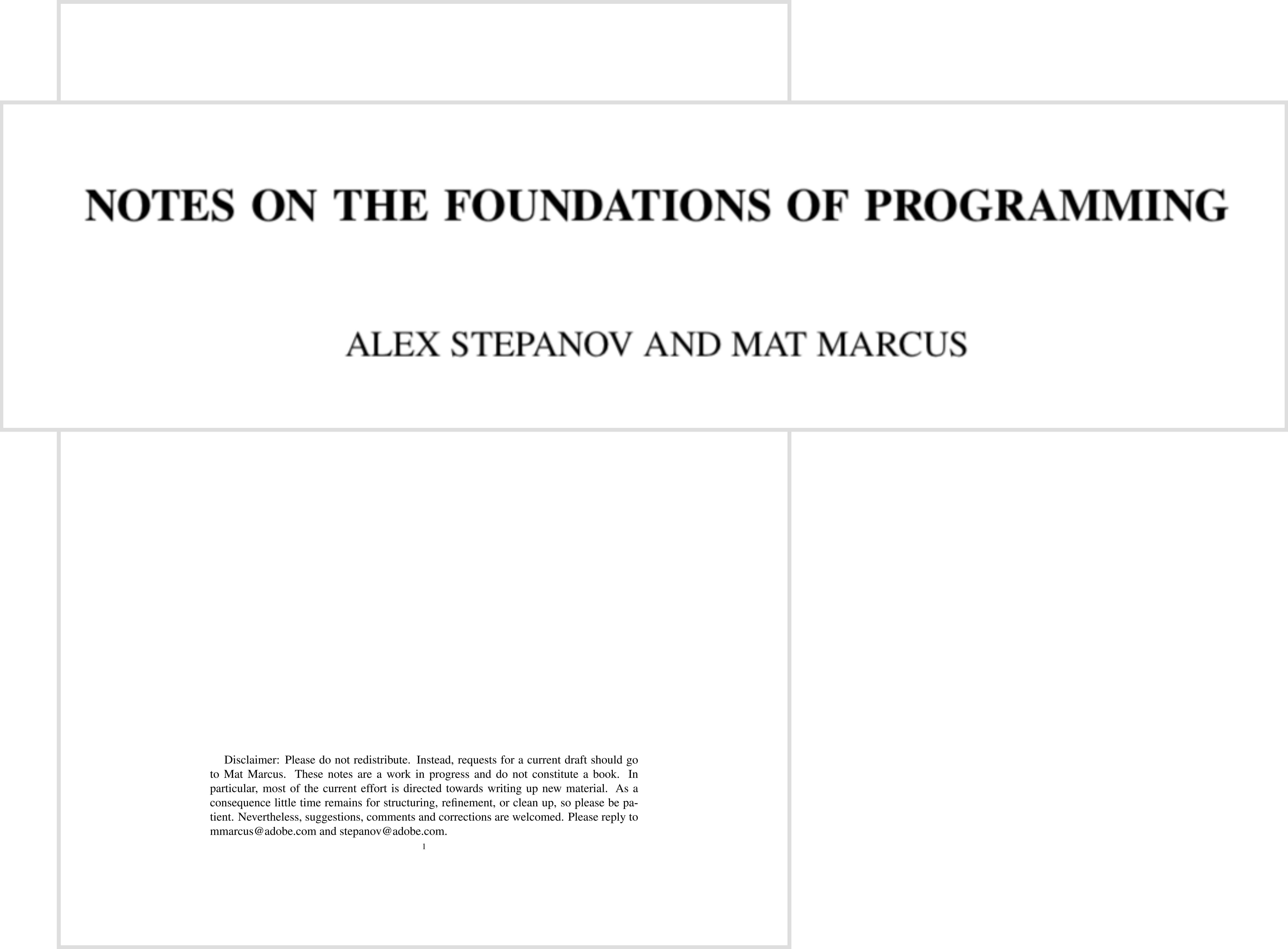
Software Technology Adobe Lab

Adobe Software Technology Lab

NOTES ON THE FOUNDATIONS OF PROGRAMMING

ALEX STEPANOV AND MAT MARCUS

Disclaimer: Please do not redistribute. Instead, requests for a current draft should go to Mat Marcus. These notes are a work in progress and do not constitute a book. In particular, most of the current effort is directed towards writing up new material. As a consequence little time remains for structuring, refinement, or clean up, so please be patient. Nevertheless, suggestions, comments and corrections are welcomed. Please reply to mmarcus@adobe.com and stepanov@adobe.com.



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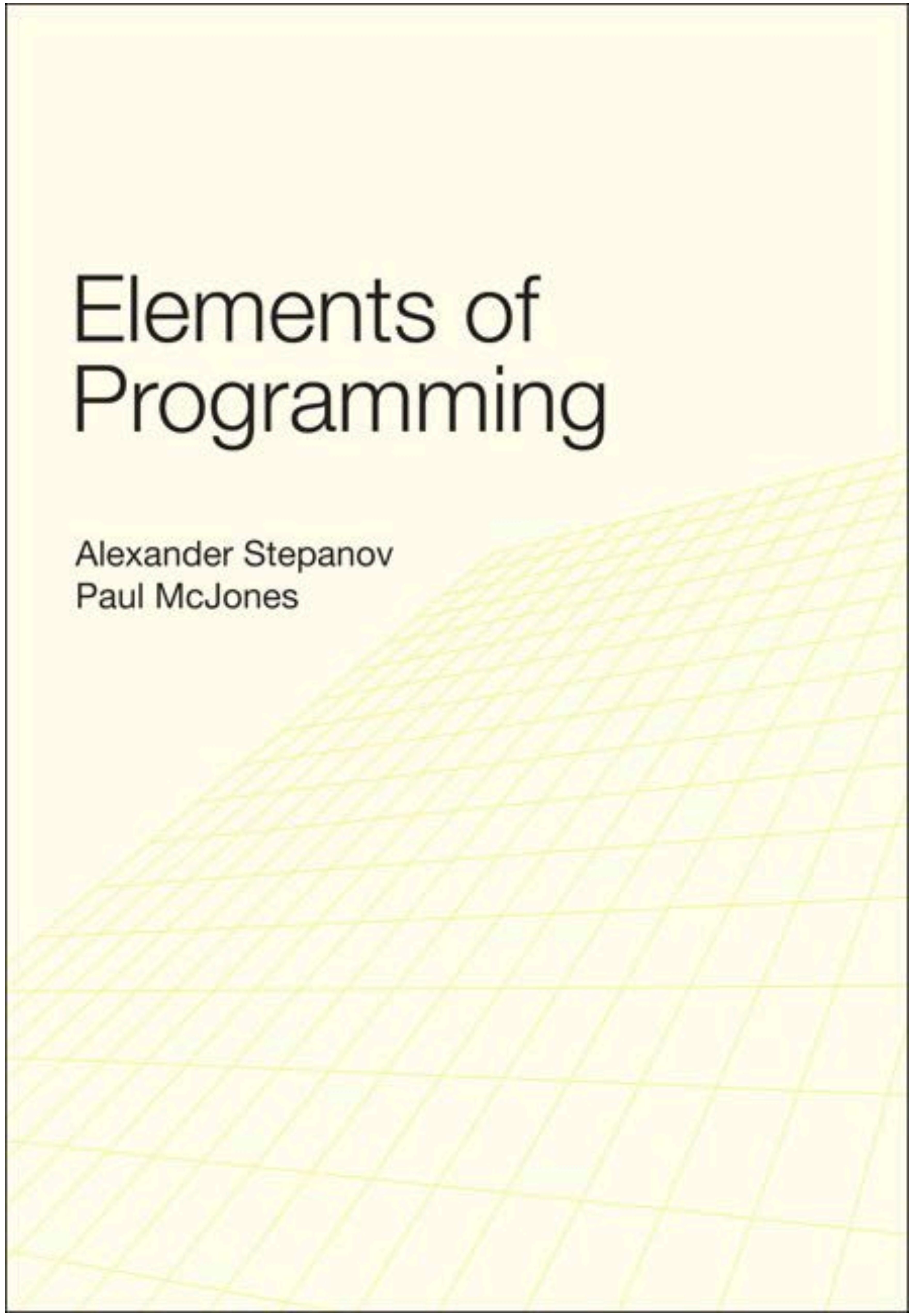
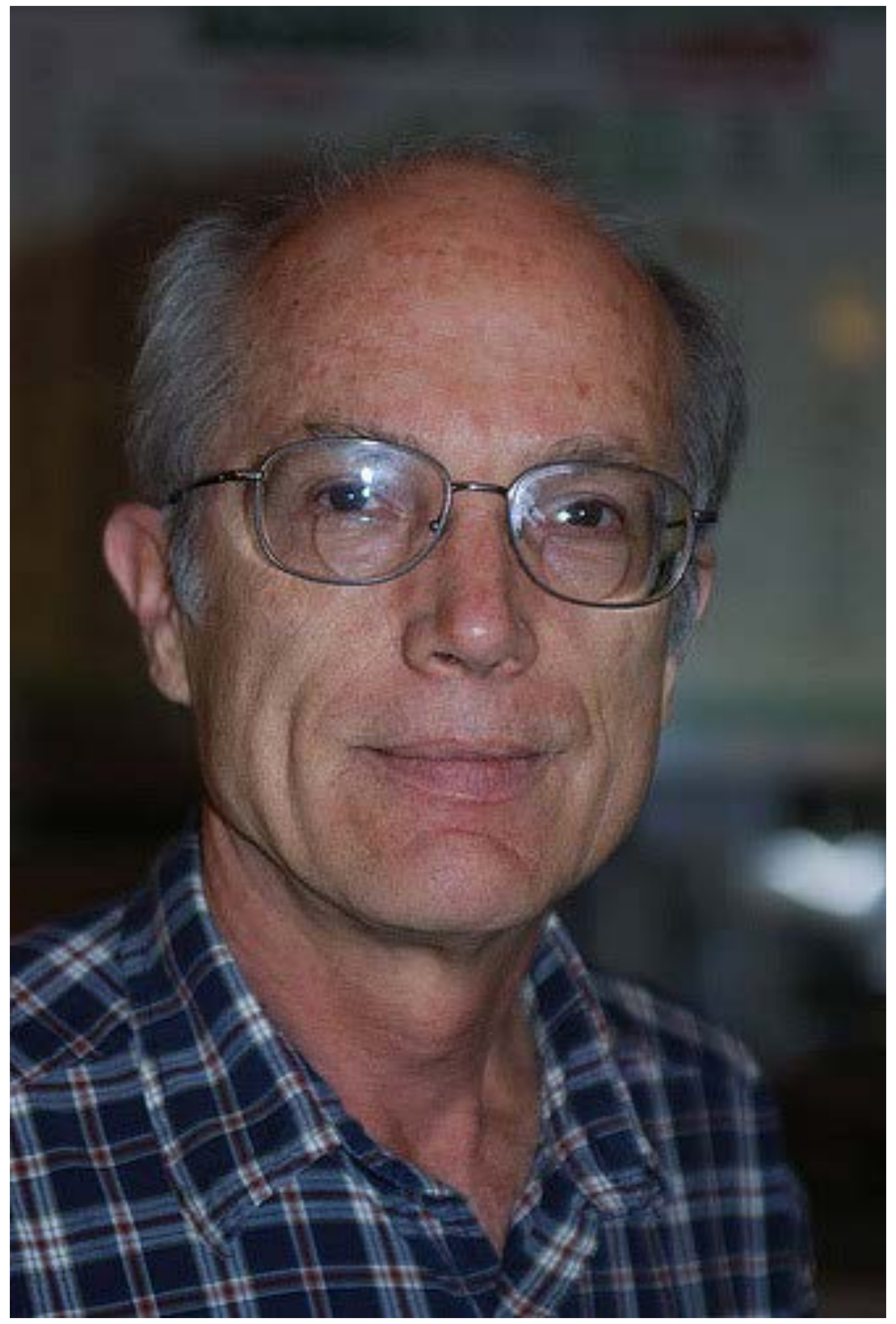
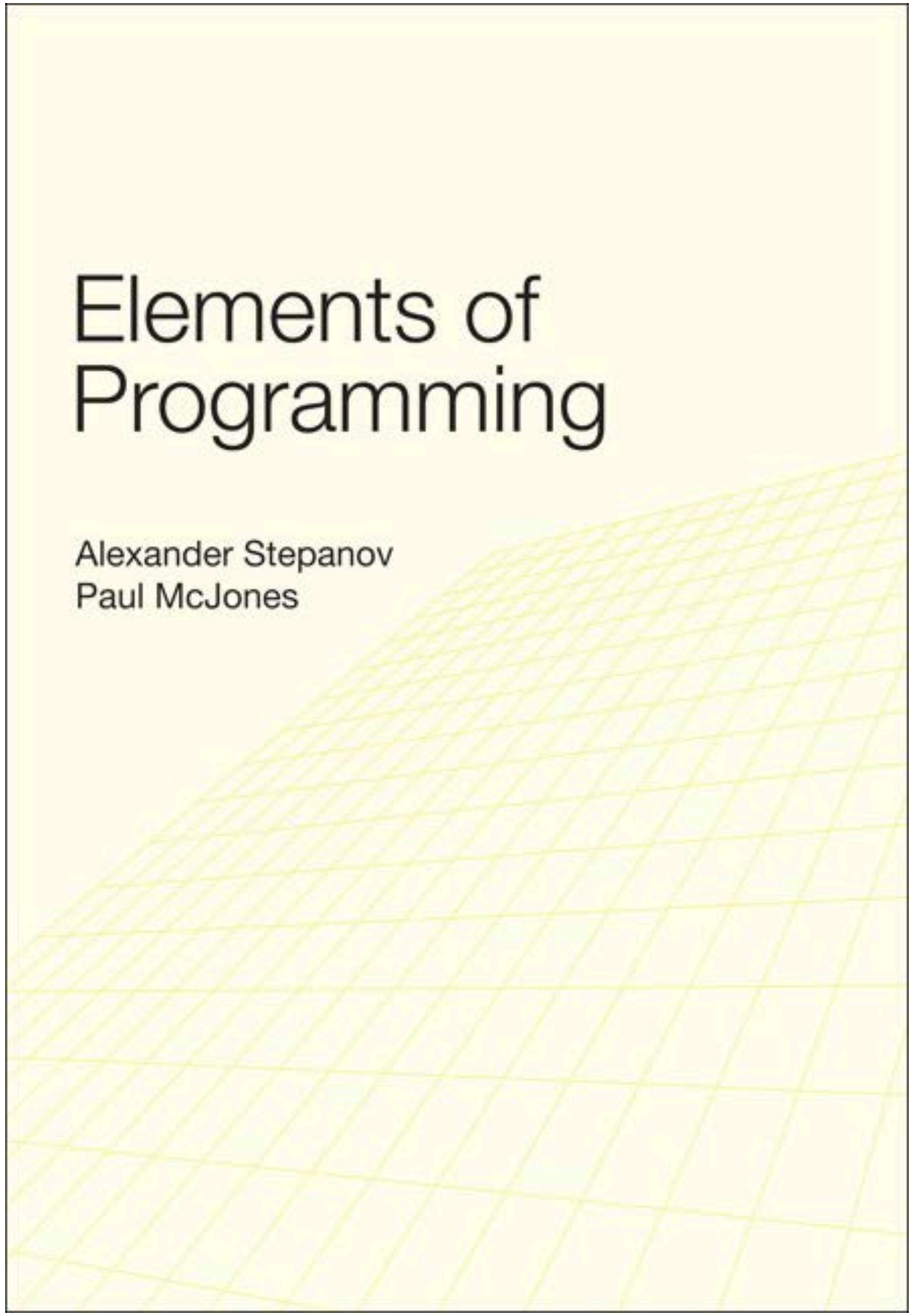


Elements of Programming

Alexander Stepanov
Paul McJones





Elements of Programming

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Paul McJones

```
template <typename I, typename P>
    requires(Mutable(I) && ForwardIterator(I) &&
        UnaryPredicate(P) && ValueType(I) == Domain(P))
I partition_semistable(I f, I l, P p) {
    // Precondition: mutable_bounded_range(f, l)
    I i = find_if(f, l, p);
    if (i == l) return i;
    I j = successor(i);
    while (true) {
        j = find_if_not(j, l, p);
        if (j == l) return i;
        swap_step(i, j);
    }
}
```

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template <typename I, typename P>
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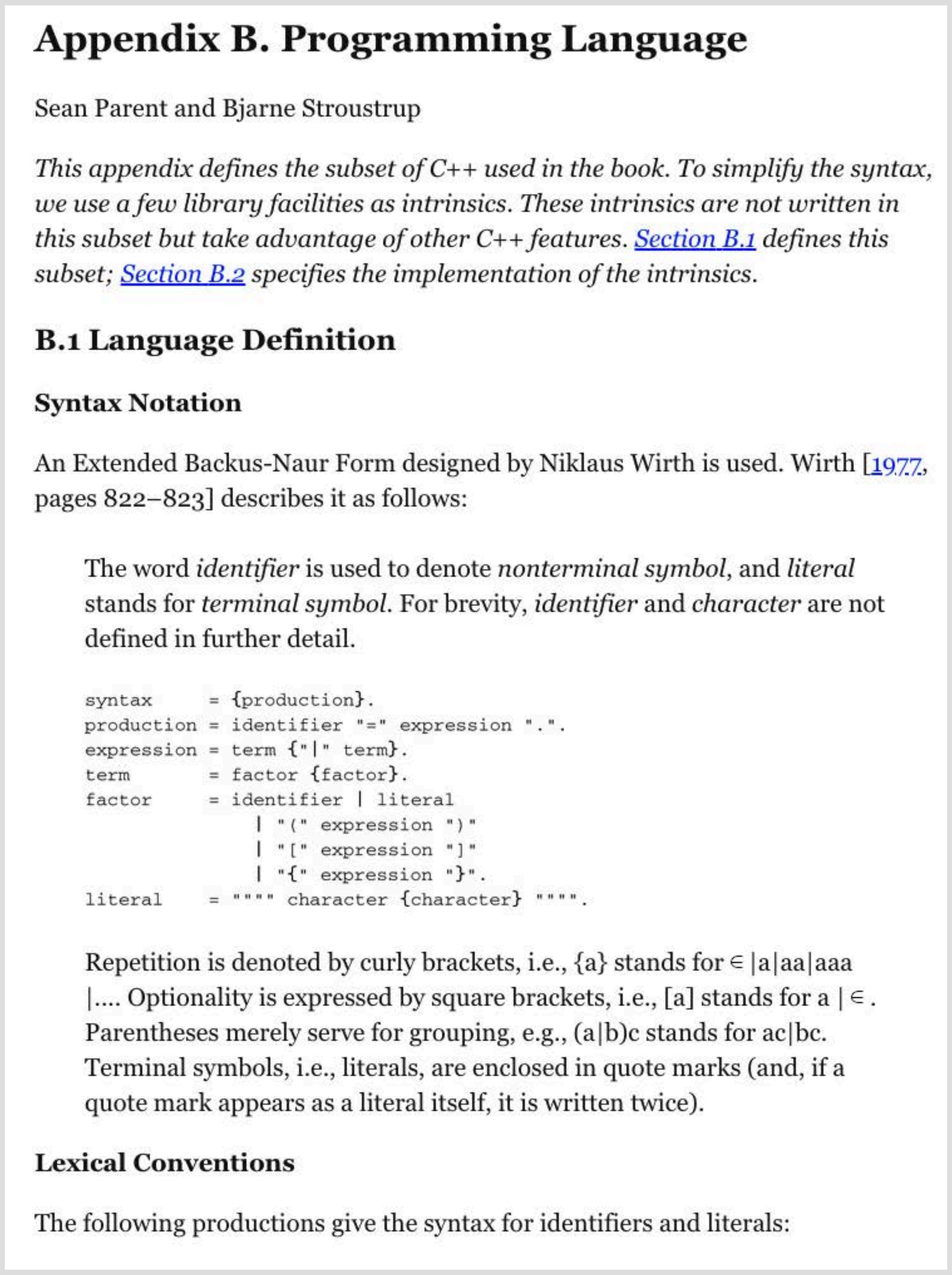


Appendix B. Programming Language

Sean Parent and Bjarne Stroustrup

This appendix defines the subset of C++ used in the book. To simplify the syntax, we use a few library facilities as intrinsics. These intrinsics are not written in this subset but take advantage of other C++ features. Section B.1 defines this subset; Section B.2 specifies the implementation of the intrinsics.

B.1 Language Definition

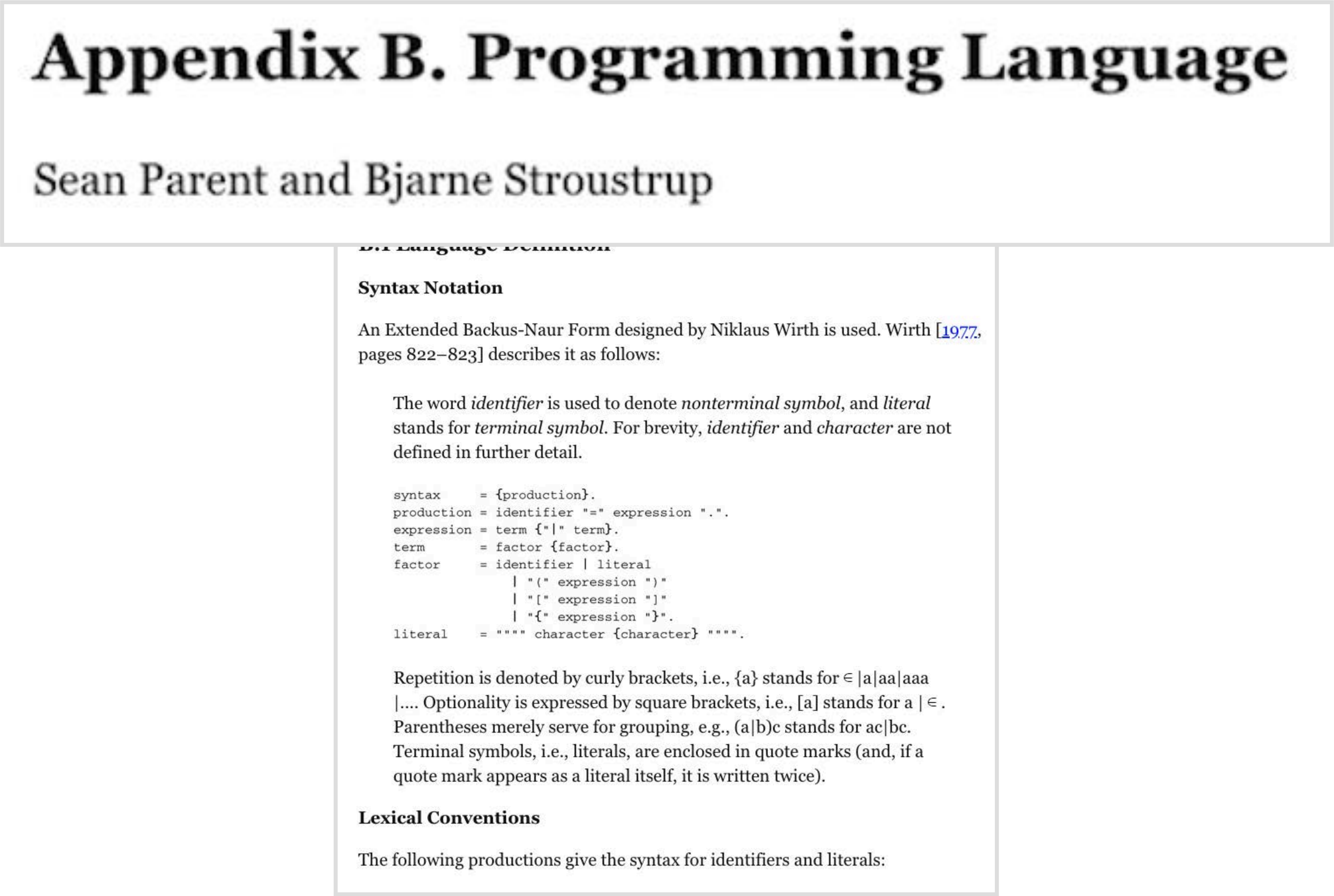
Syntax Notation

An Extended Backus-Naur Form designed by Niklaus Wirth is used. Wirth [1977, pages 822–823] describes it as follows:

> The word *identifier* is used to denote *nonterminal symbol*, and *literal* stands for *terminal symbol*. For brevity, *identifier* and *character* are not defined in further detail.
>
> ```
> syntax     = {production}.
> production = identifier "=" expression ".".
> expression = term {"|" term}.
> term       = factor {factor}.
> factor     = identifier | literal
>                | "(" expression ")"
>                | "[" expression "]"
>                | "{" expression "}".
> literal    = """" character {character} """".
> ```
>
> Repetition is denoted by curly brackets, i.e., {a} stands for ∈ |a|aa|aaa |.... Optionality is expressed by square brackets, i.e., [a] stands for a | ∈ . Parentheses merely serve for grouping, e.g., (a|b)c stands for ac|bc. Terminal symbols, i.e., literals, are enclosed in quote marks (and, if a quote mark appears as a literal itself, it is written twice).

Lexical Conventions

The following productions give the syntax for identifiers and literals:



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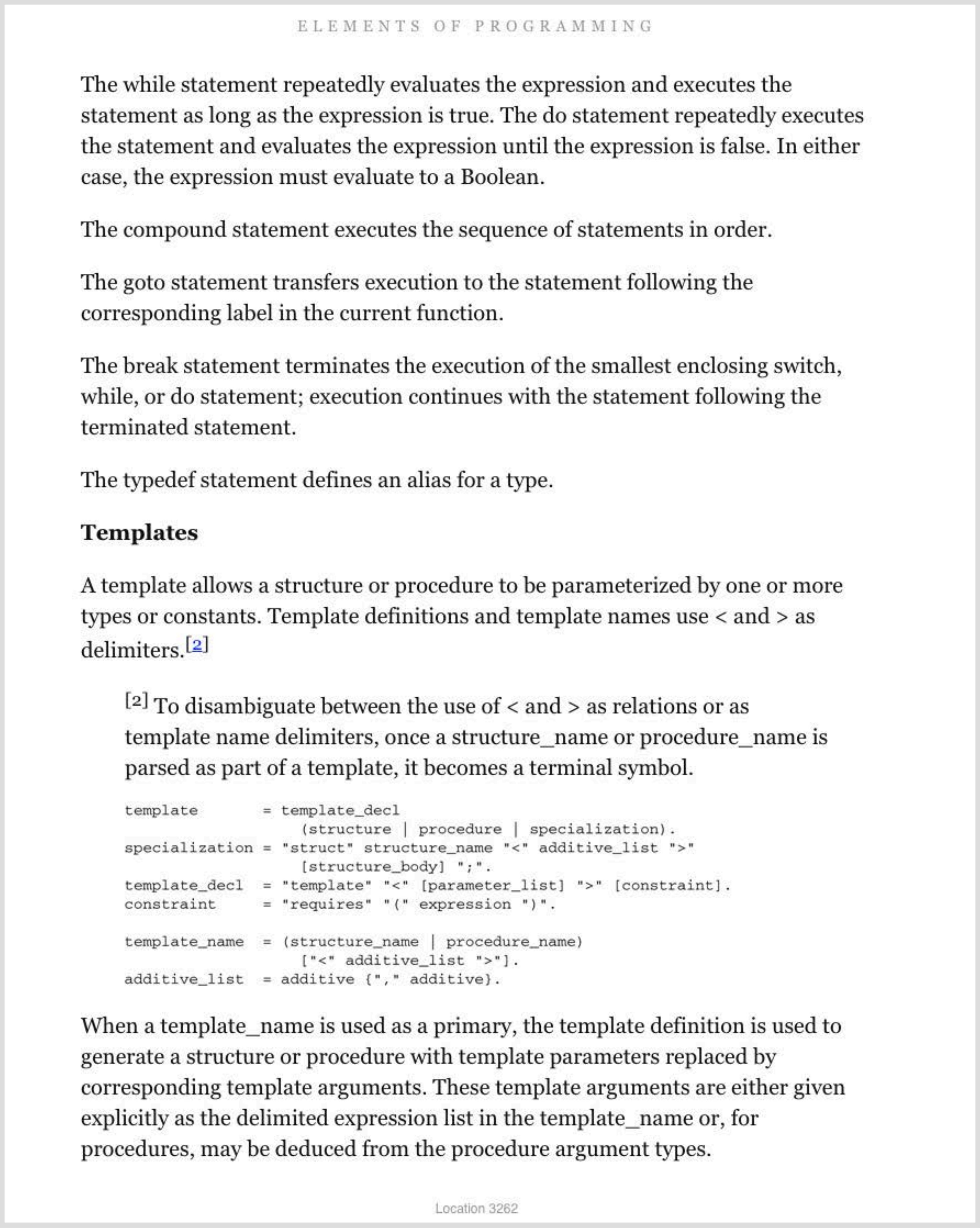


The while statement repeatedly evaluates the expression and executes the statement as long as the expression is true. The do statement repeatedly executes the statement and evaluates the expression until the expression is false. In either case, the expression must evaluate to a Boolean.

The compound statement executes the sequence of statements in order.

The goto statement transfers execution to the statement following the corresponding label in the current function.

The break statement terminates the execution of the smallest enclosing switch, while, or do statement; execution continues with the statement following the terminated statement.

The typedef statement defines an alias for a type.

Templates

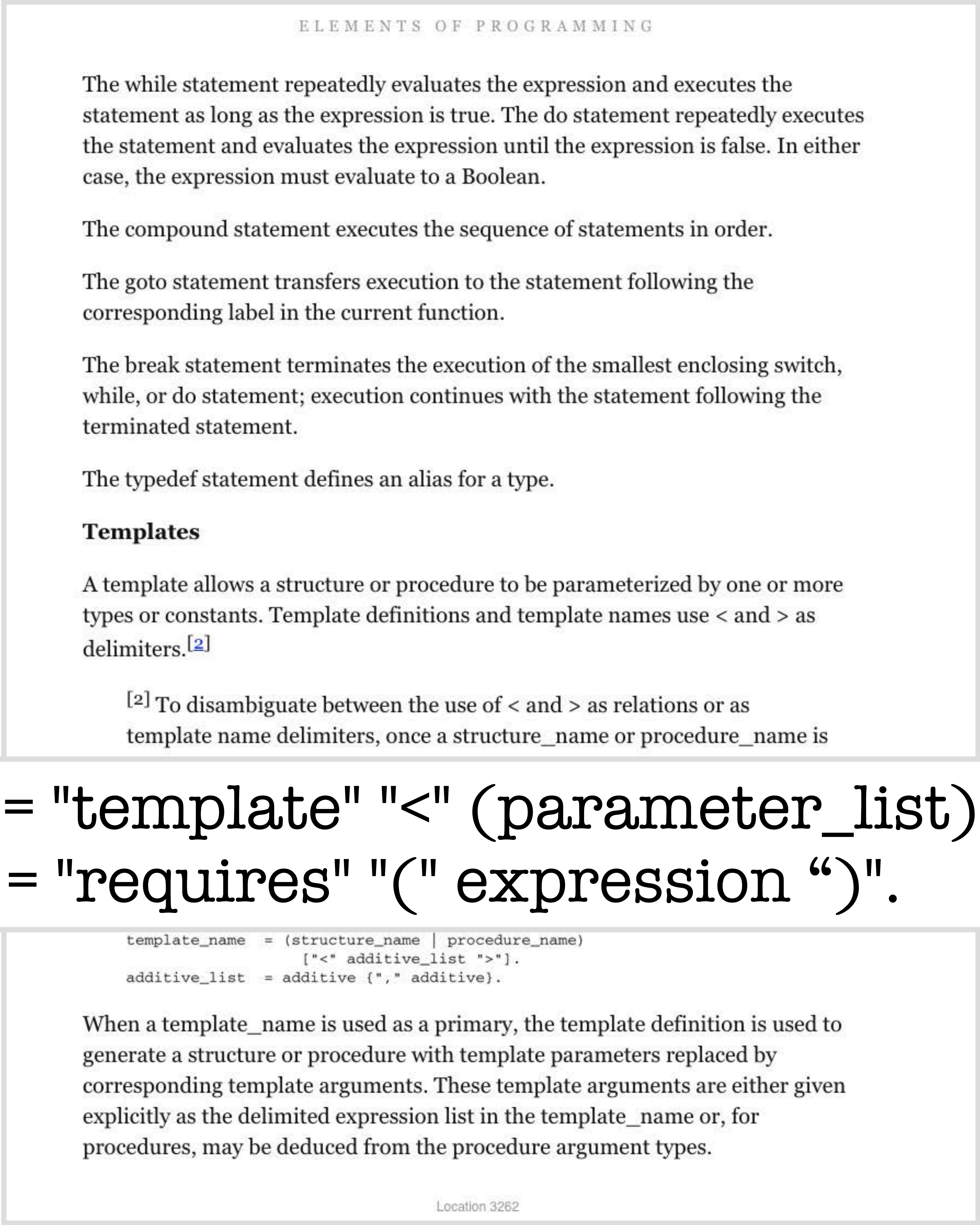
A template allows a structure or procedure to be parameterized by one or more types or constants. Template definitions and template names use < and > as delimiters.[2]

> [2] To disambiguate between the use of < and > as relations or as template name delimiters, once a structure_name or procedure_name is parsed as part of a template, it becomes a terminal symbol.

```
template       = template_decl
                   (structure | procedure | specialization).
specialization = "struct" structure_name "<" additive_list ">"
                   [structure_body] ";".
template_decl  = "template" "<" [parameter_list] ">" [constraint].
constraint     = "requires" "(" expression ")".

template_name  = (structure_name | procedure_name)
                   ["<" additive_list ">"].
additive_list  = additive {"," additive}.
```

When a template_name is used as a primary, the template definition is used to generate a structure or procedure with template parameters replaced by corresponding template arguments. These template arguments are either given explicitly as the delimited expression list in the template_name or, for procedures, may be deduced from the procedure argument types.



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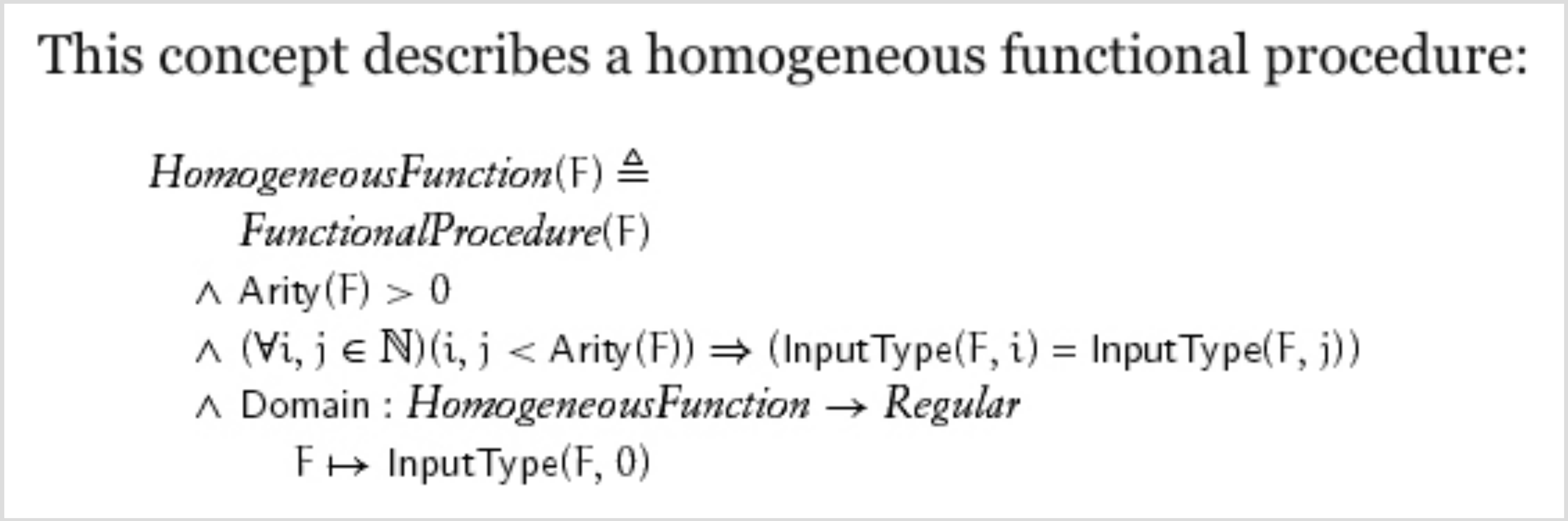
```
template_decl   = "template" "<" (parameter_list) ">" [constraint].
constraint      = "requires" "(" expression “)".
```



This concept describes a homogeneous functional procedure:

$$
\begin{aligned}
&\mathit{HomogeneousFunction}(\mathsf{F}) \triangleq \\
&\qquad \mathit{FunctionalProcedure}(\mathsf{F}) \\
&\quad \wedge\ \mathsf{Arity}(\mathsf{F}) > 0 \\
&\quad \wedge\ (\forall \mathsf{i}, \mathsf{j} \in \mathbb{N})(\mathsf{i}, \mathsf{j} < \mathsf{Arity}(\mathsf{F})) \Rightarrow (\mathsf{InputType}(\mathsf{F}, \mathsf{i}) = \mathsf{InputType}(\mathsf{F}, \mathsf{j})) \\
&\quad \wedge\ \mathsf{Domain} : \mathit{HomogeneousFunction} \to \mathit{Regular} \\
&\qquad\quad \mathsf{F} \mapsto \mathsf{InputType}(\mathsf{F}, 0)
\end{aligned}
$$

2006

Concepts: Linguistic Support for Generic Programming in C++

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Abstract

Generic programming has emerged as an important technique for the development of highly reusable and efficient software libraries. In C++, generic programming is enabled by the flexibility of templates, the C++ type parametrization mechanism. However, the power of templates comes with a price: generic (template) libraries can be more difficult to use and develop than non-template libraries and their misuse results in notoriously confusing error messages. As currently defined in C++98, templates are unconstrained, and type-checking of templates is performed late in the compilation process, i.e., after the use of a template has been combined with its definition. To improve the support for generic programming in C++, we introduce *concepts* to express the syntactic and semantic behavior of types and to constrain the type parameters in a C++ template. Using concepts, type-checking of template definitions is separated from their uses, thereby making templates easier to use and easier to compile. These improvements are achieved without limiting the flexibility of templates or decreasing their performance—in fact their expressive power is increased. This paper describes the language extensions supporting concepts, their use in the expression of the C++ Standard Template Library, and their implementation in the ConceptGCC compiler. Concepts are candidates for inclusion in the upcoming revision of the ISO C++ standard, C++0x.

Categories and Subject Descriptors D.3.3 [*Programming Languages*]: Language Constructs and Features—Abstract data types; D.3.3 [*Programming Languages*]: Language Constructs and Features—Polymorphism; D.2.13 [*Software Engineering*]: Reusable Software—Reusable libraries

General Terms Design, Languages

Keywords Generic programming, constrained generics, parametric polymorphism, C++ templates, C++0x, concepts

1. Introduction

The C++ language [25, 62] supports parametrized types and functions in the form of *templates*. Templates provide a unique combination of features that have allowed them to be used for many different programming paradigms, including Generic Programming [3,44], Generative Programming [11], and Template Metaprogramming [1, 66]. Much of the flexibility of C++ templates comes from their unconstrained nature: a template can perform any operation on its template parameters, including compile-time type computations, allowing the emulation of the basic features required for diverse programming paradigms. Another essential part of templates is their ability to provide abstraction without performance degradation: templates provide sufficient information to a compiler's optimizers (especially the inliner) to generate code that is optimal in both time and space.

Consequently, templates have become the preferred implementation style for a vast array of reusable, efficient C++ libraries [2, 6, 14, 20, 32, 54, 55, 65], many of which are built upon the Generic Programming methodology exemplified by the C++ Standard Template Library (STL) [42, 60]. Aided by the discovery of numerous *ad hoc* template techniques [28, 46, 56, 66, 67], C++ libraries are becoming more powerful, more flexible, and more expressive.

However, these improvements come at the cost of implementation complexity [61, 63]: authors of C++ libraries typically rely on a grab-bag of template tricks, many of which are complex and poorly documented. Where library interfaces are rigorously separated from library implementation, the complexity of implementation of a library is not a problem for its users. However, templates rely on the absence of modular (separate) type-checking for flexibility and performance. Therefore, the complexities of library implementation leak through to library users. This problem manifests itself most visibly in spectacularly poor error messages for simple mistakes. Consider:

```
list<int> lst;
sort(lst.begin(), lst.end());
```

Attempting to compile this code with a recent version of the GNU C++ compiler [17] produces more than two kilobytes of output, containing six different error messages. Worse, the errors reported provide line numbers and file names that point to the implementation of the STL sort() function and its helper functions. The only clue provided to users that this error was triggered by their own code (rather than by a bug in the STL implementation) is the following innocuous line of output:

```
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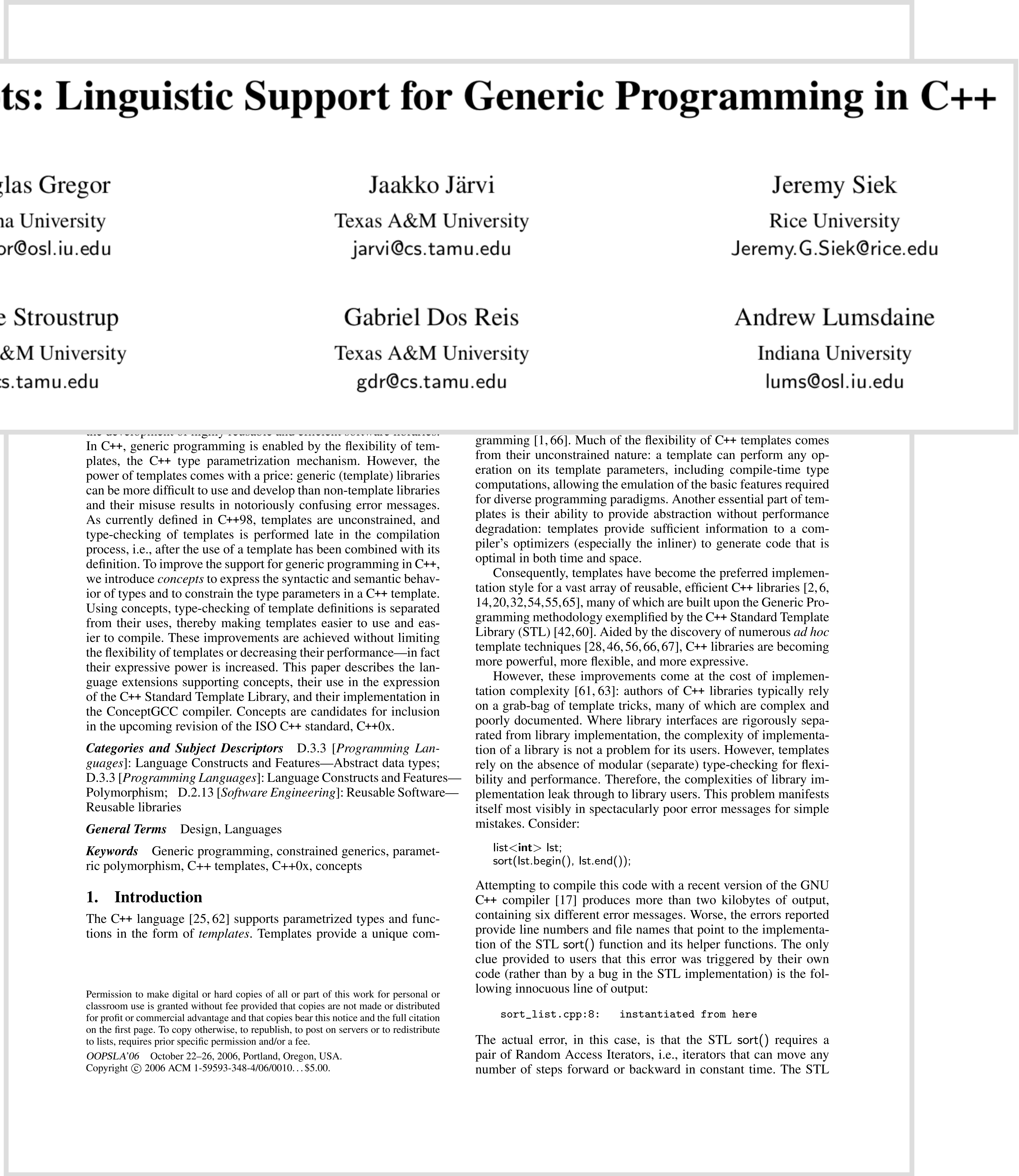

Concepts: Linguistic Support for Generic Programming in C++

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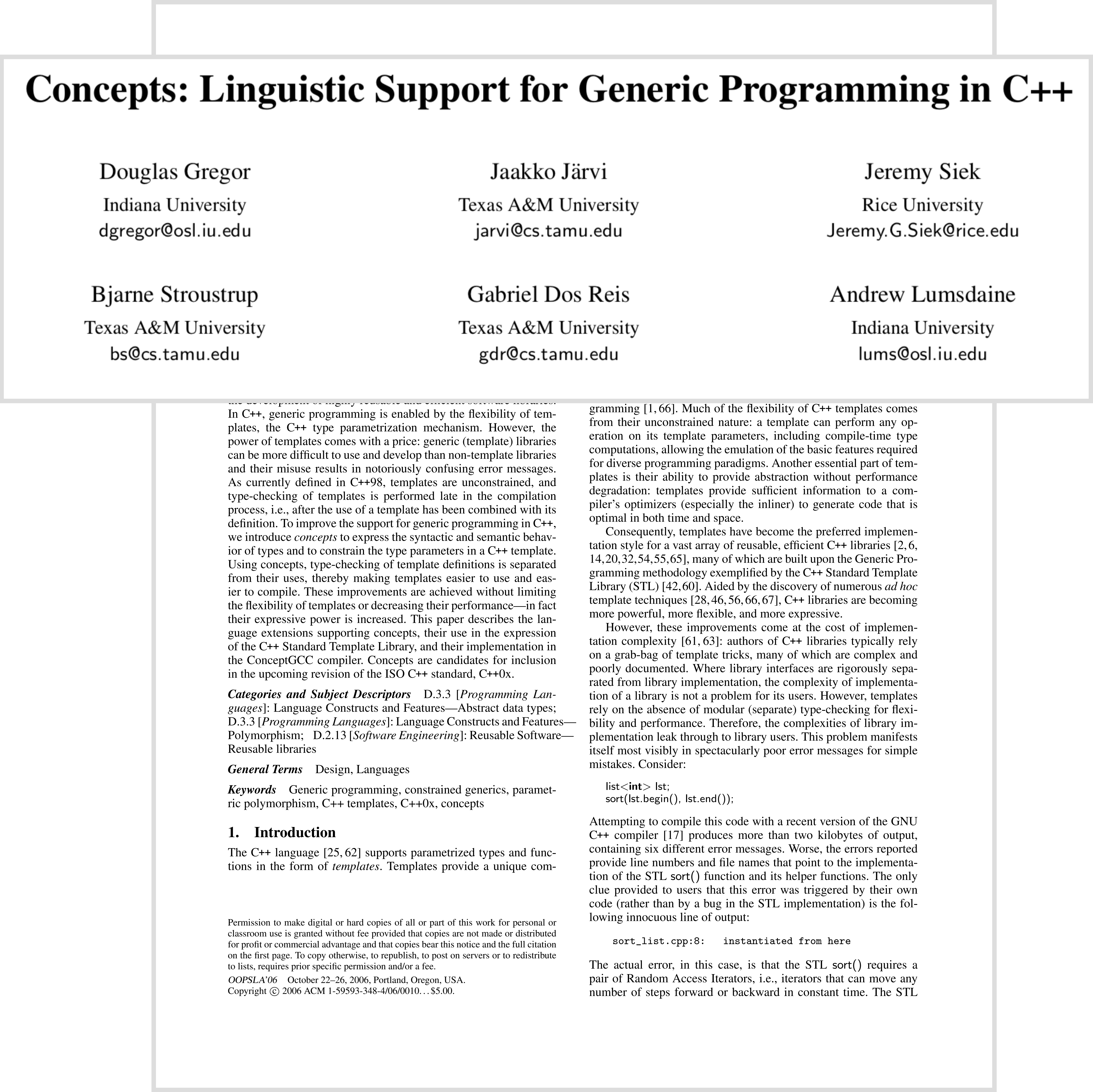

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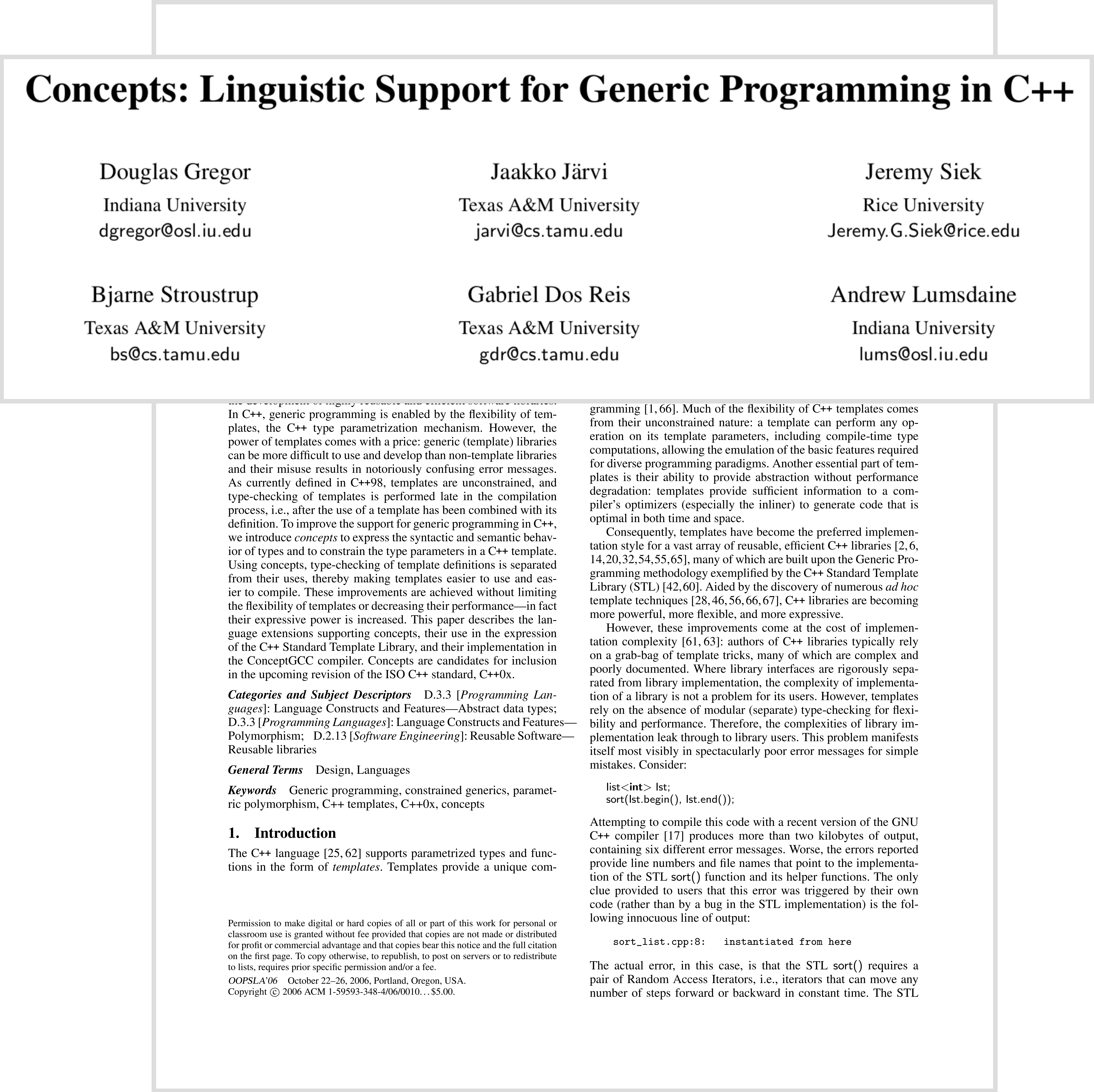

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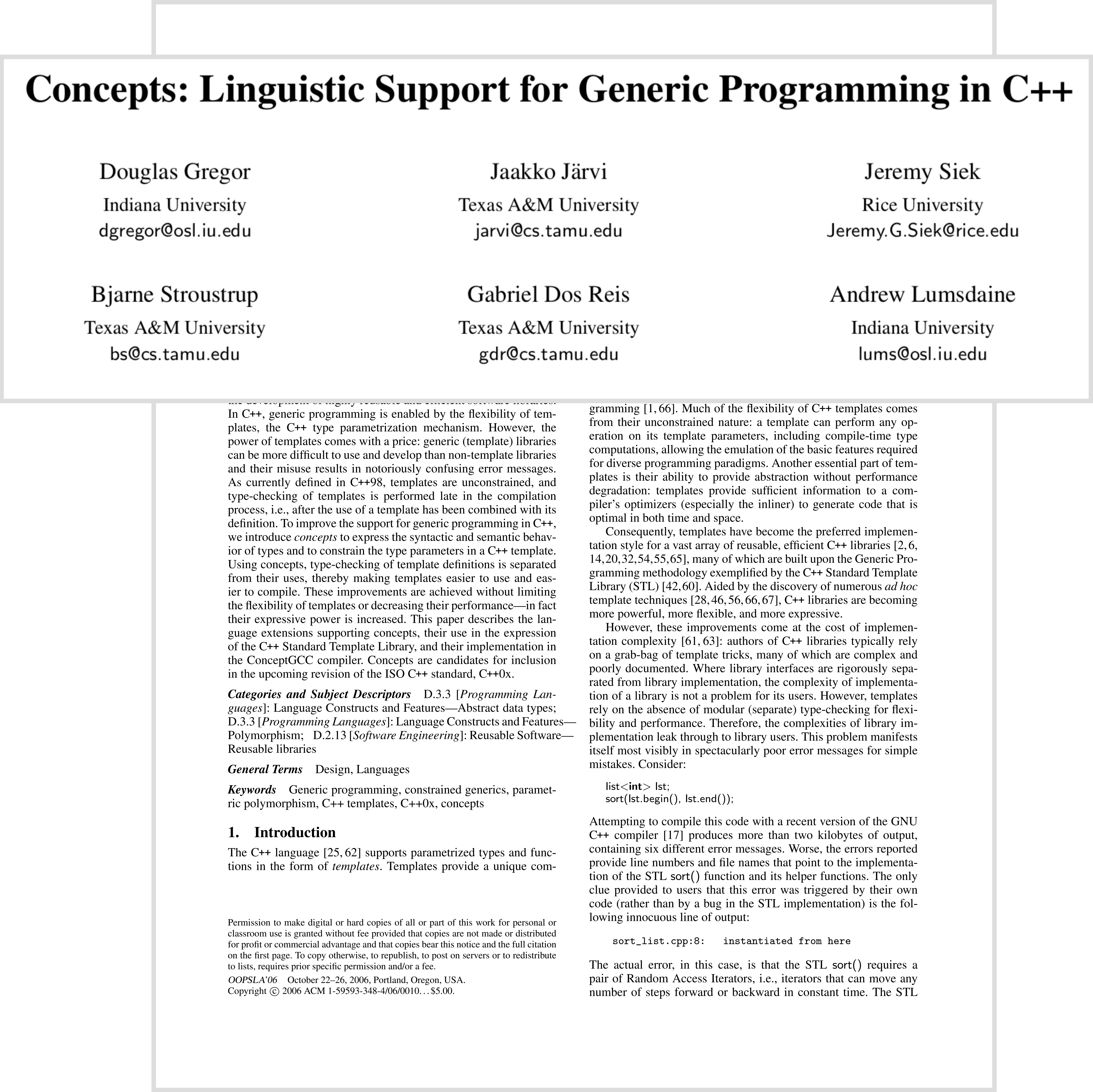

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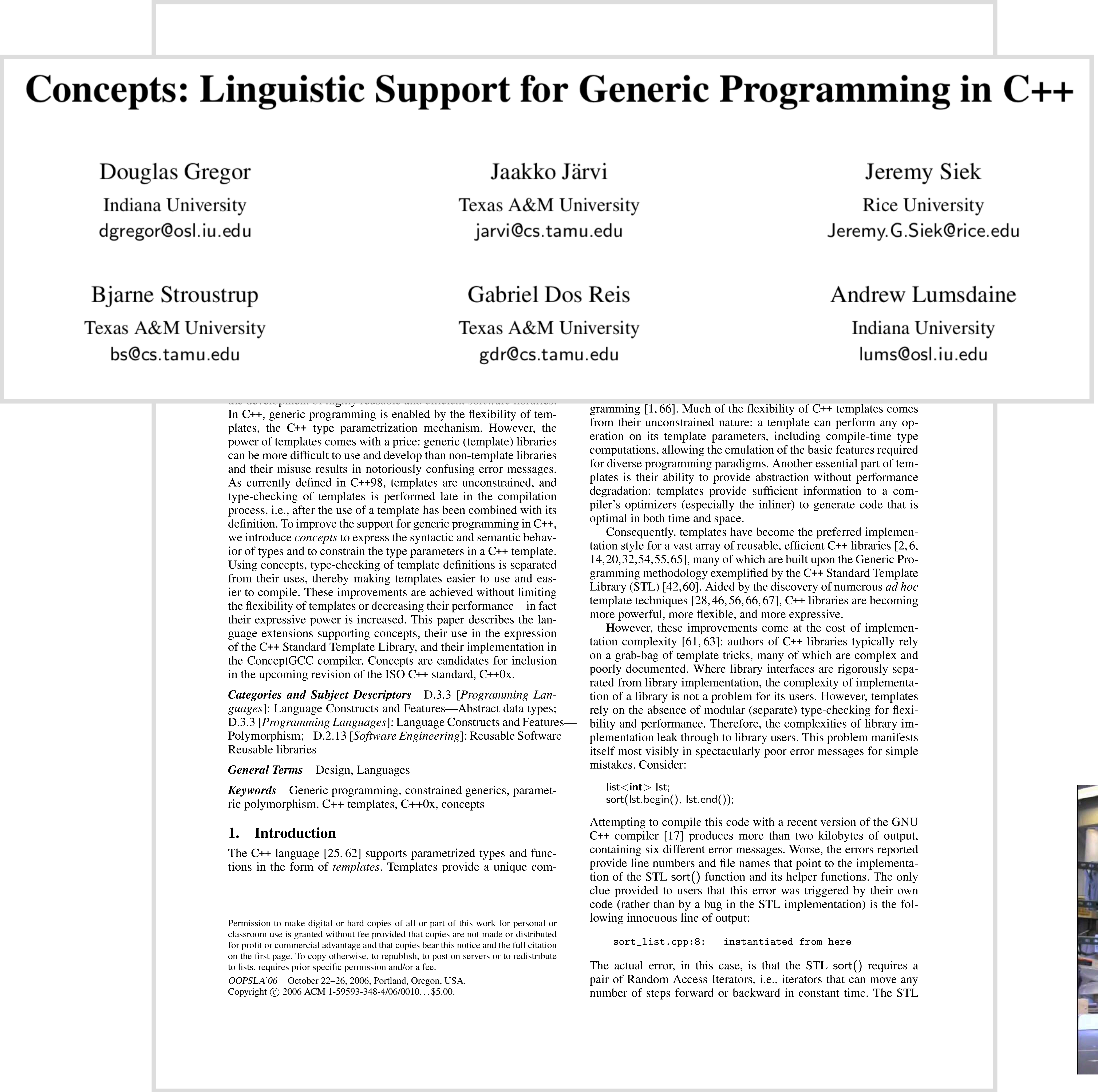

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2011

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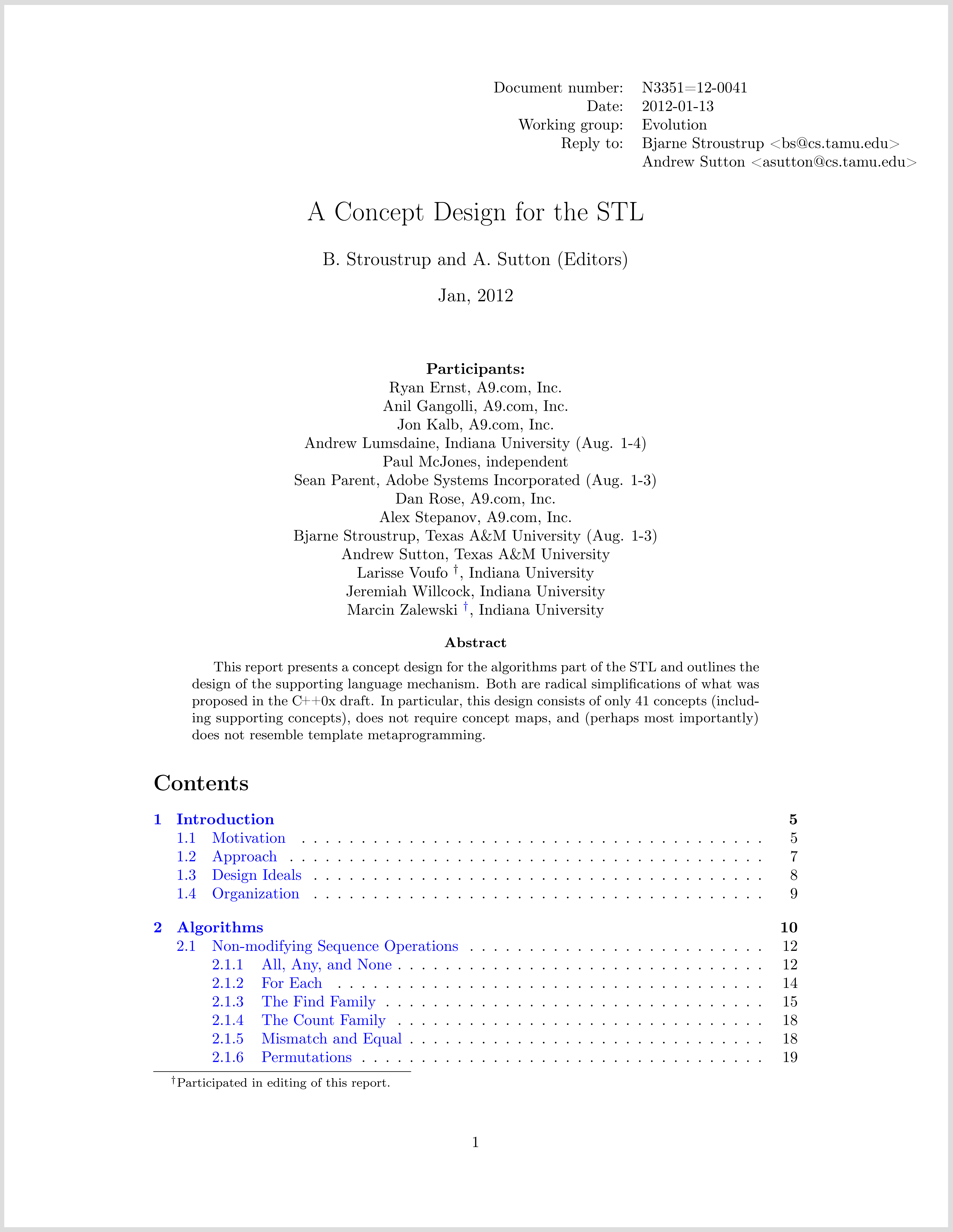
2012



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Date: 2012-01-13
Working group: Evolution
Reply to: Bjarne Stroustrup <bs@cs.tamu.edu>
Andrew Sutton <asutton@cs.tamu.edu>

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B. Stroustrup and A. Sutton (Editors)

Jan, 2012

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Jon Kalb, A9.com, Inc.
Andrew Lumsdaine, Indiana University (Aug. 1-4)
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Dan Rose, A9.com, Inc.
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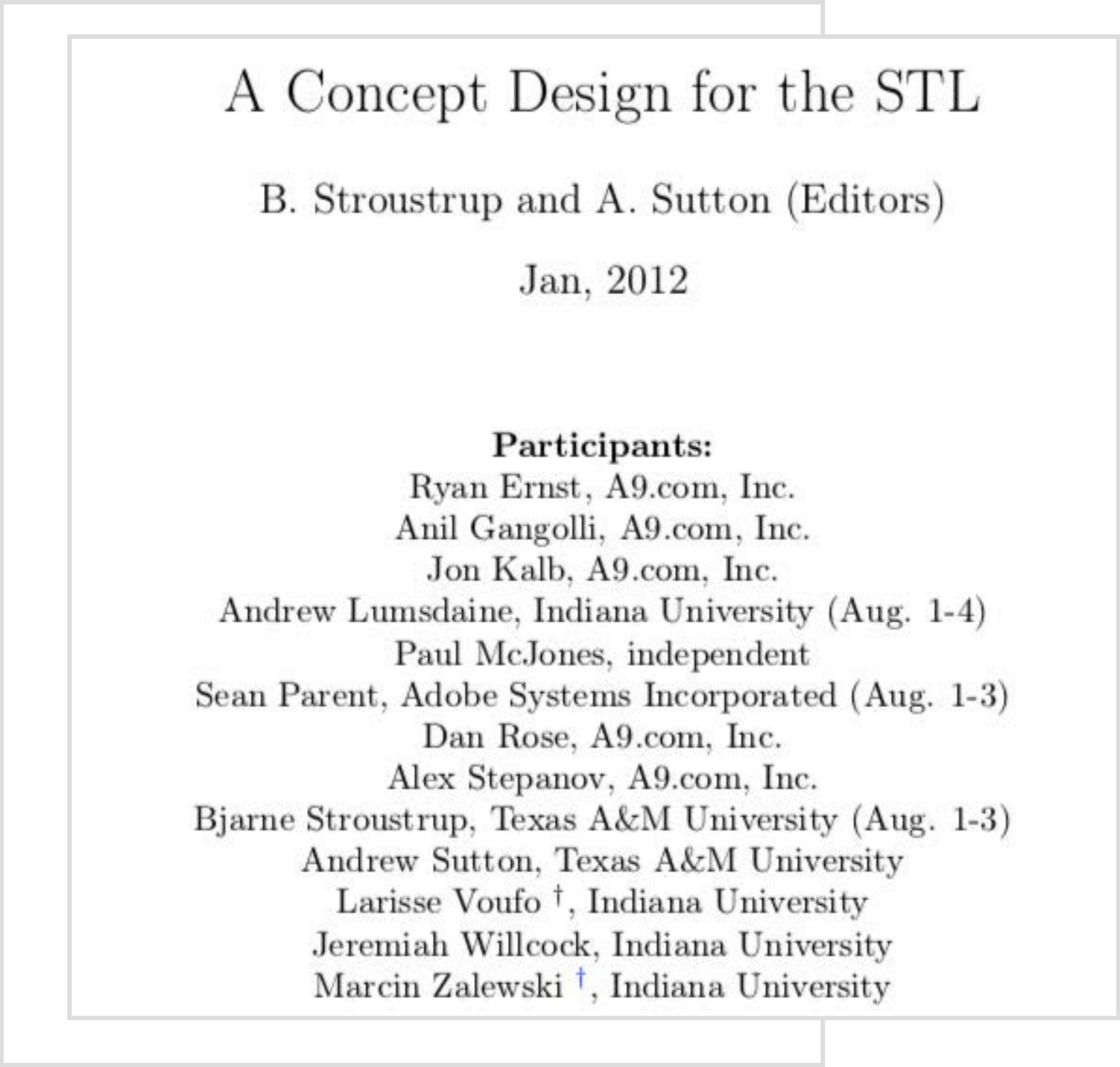
Bjarne Stroustrup, Texas A&M University (Aug. 1-3)
Andrew Sutton, Texas A&M University
Larisse Voufo [†], Indiana University
Jeremiah Willcock, Indiana University
Marcin Zalewski [†], Indiana University

Abstract

This report presents a concept design for the algorithms part of the STL and outlines the design of the supporting language mechanism. Both are radical simplifications of what was proposed in the C++0x draft. In particular, this design consists of only 41 concepts (including supporting concepts), does not require concept maps, and (perhaps most importantly) does not resemble template metaprogramming.

Contents

[†] Participated in editing of this report.



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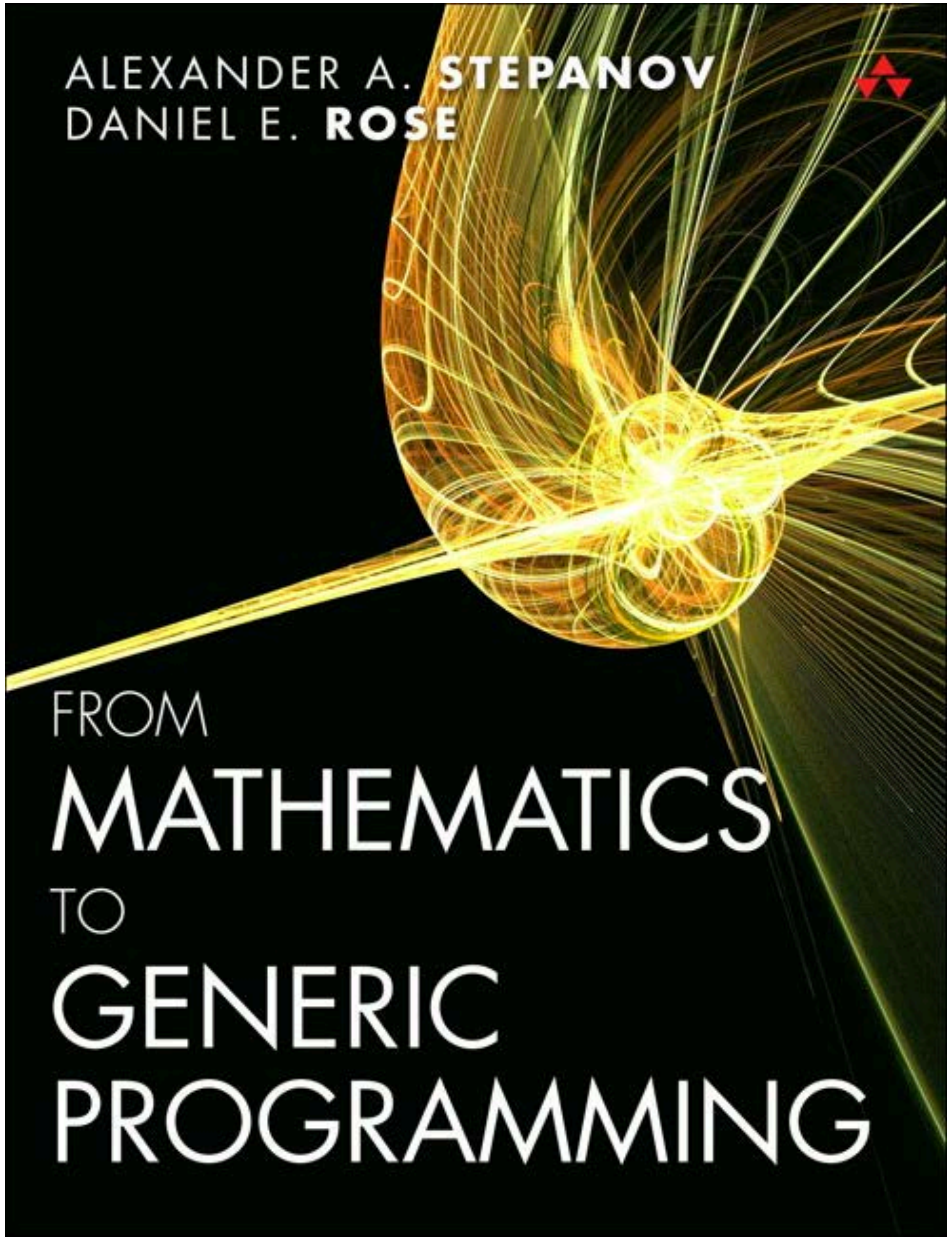

ALEXANDER A. **STEPANOV**
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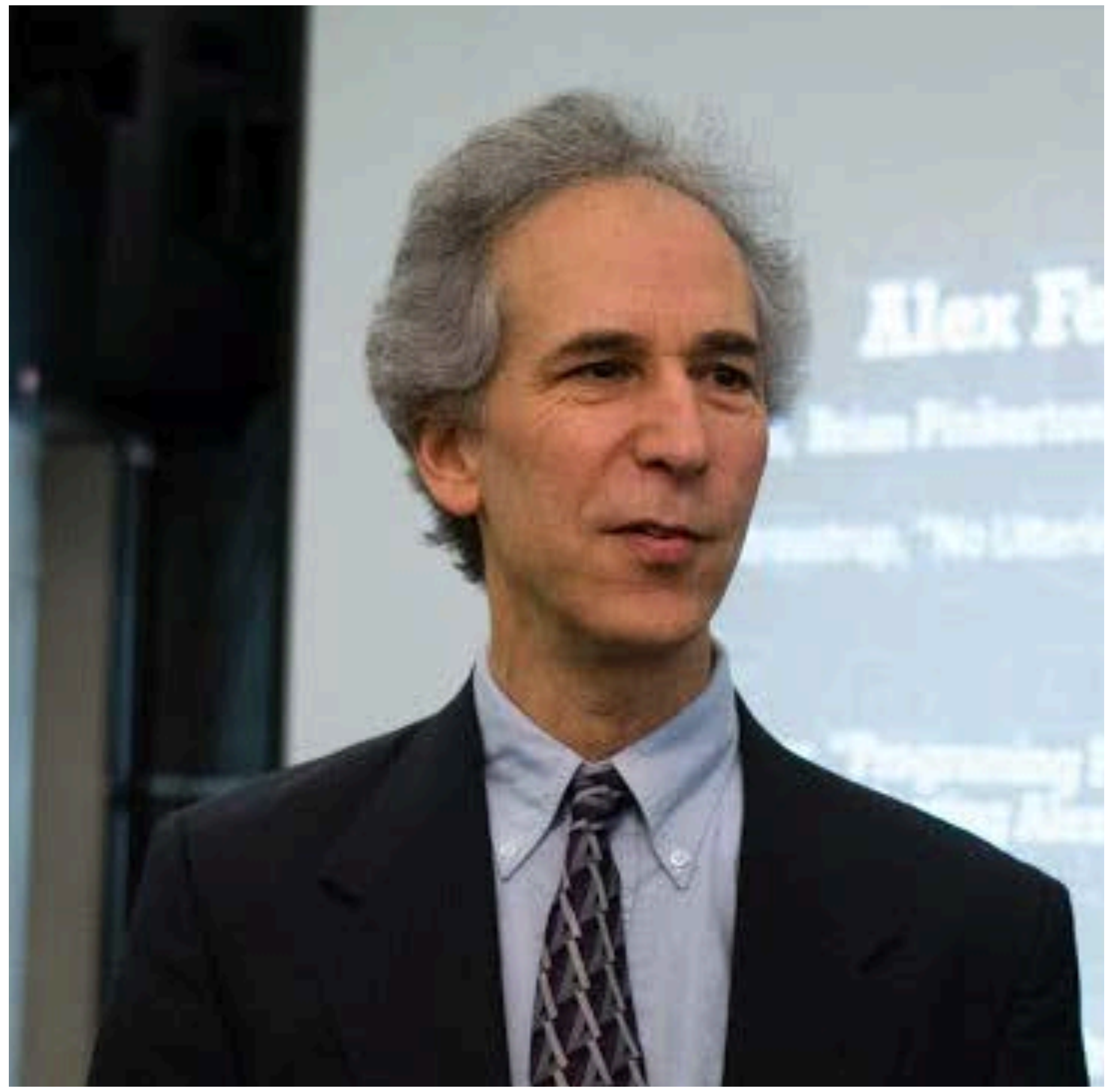

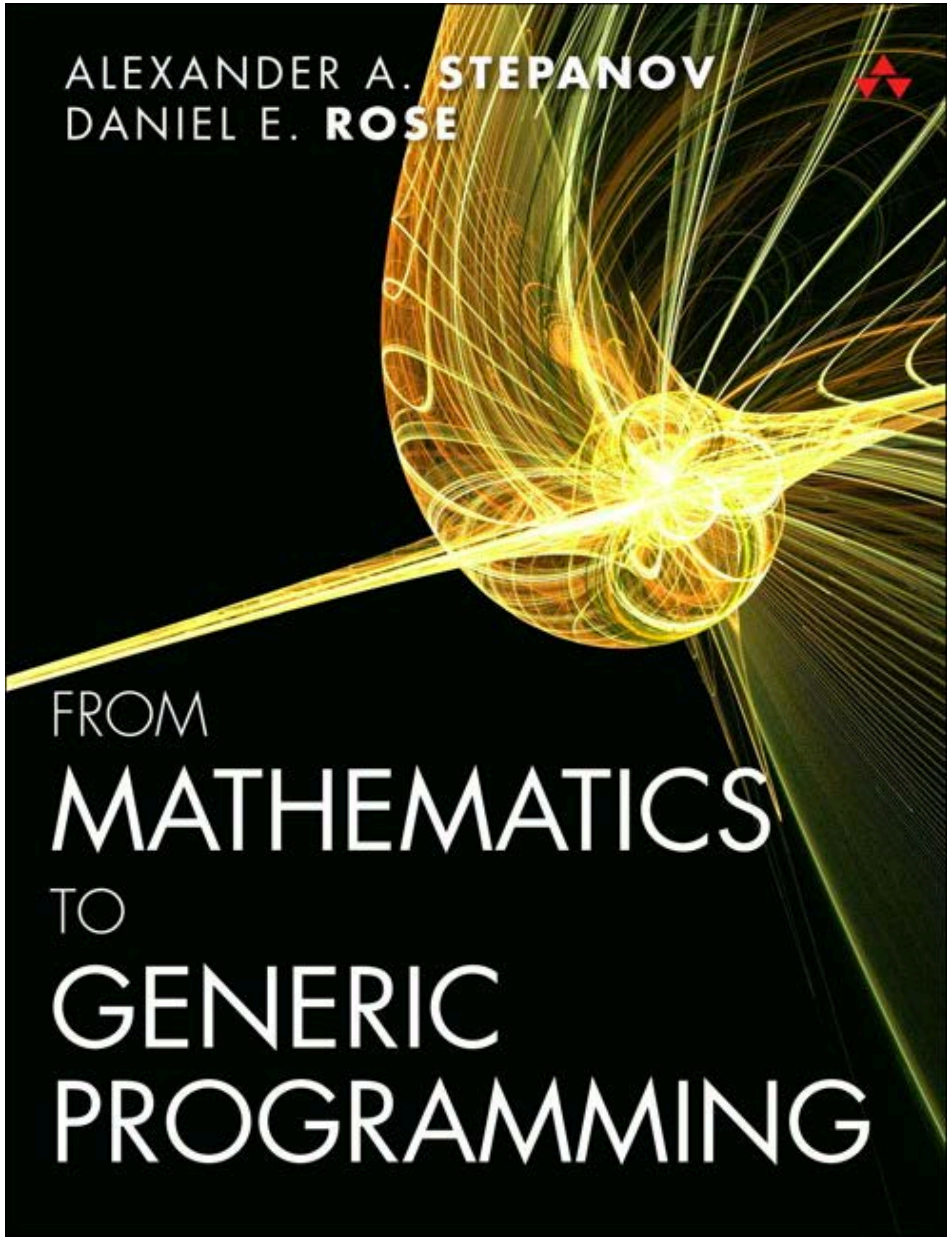

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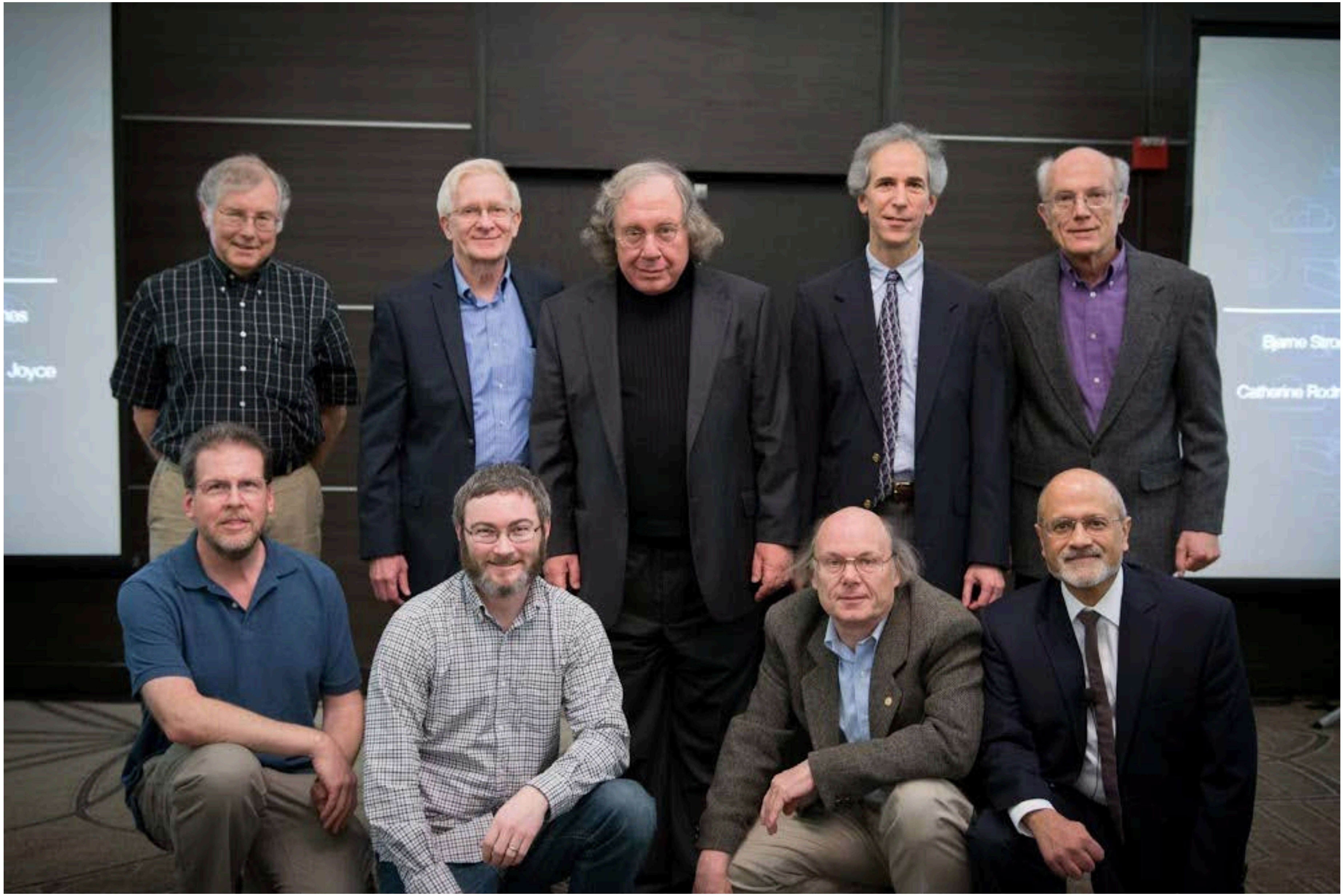

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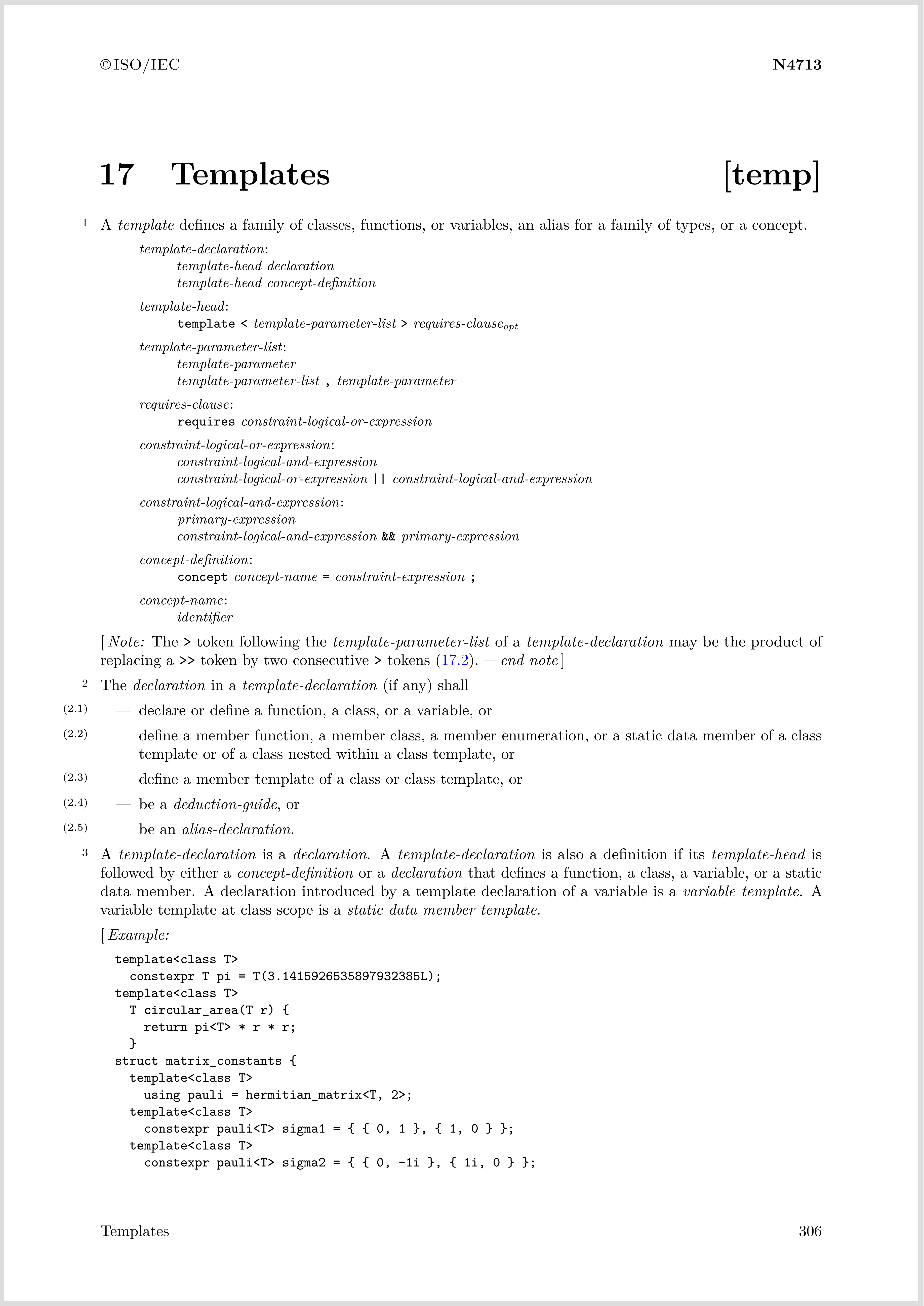


17 Templates [temp]

1 A *template* defines a family of classes, functions, or variables, an alias for a family of types, or a concept.

template-declaration:
  template-head declaration
  template-head concept-definition

template-head:
  `template <` *template-parameter-list* `>` *requires-clause*$_{opt}$

template-parameter-list:
  template-parameter
  template-parameter-list `,` *template-parameter*

requires-clause:
  `requires` *constraint-logical-or-expression*

constraint-logical-or-expression:
  constraint-logical-and-expression
  constraint-logical-or-expression `||` *constraint-logical-and-expression*

constraint-logical-and-expression:
  primary-expression
  constraint-logical-and-expression `&&` *primary-expression*

concept-definition:
  `concept` *concept-name* `=` *constraint-expression* `;`

concept-name:
  identifier

[*Note:* The `>` token following the *template-parameter-list* of a *template-declaration* may be the product of replacing a `>>` token by two consecutive `>` tokens (17.2). — *end note*]

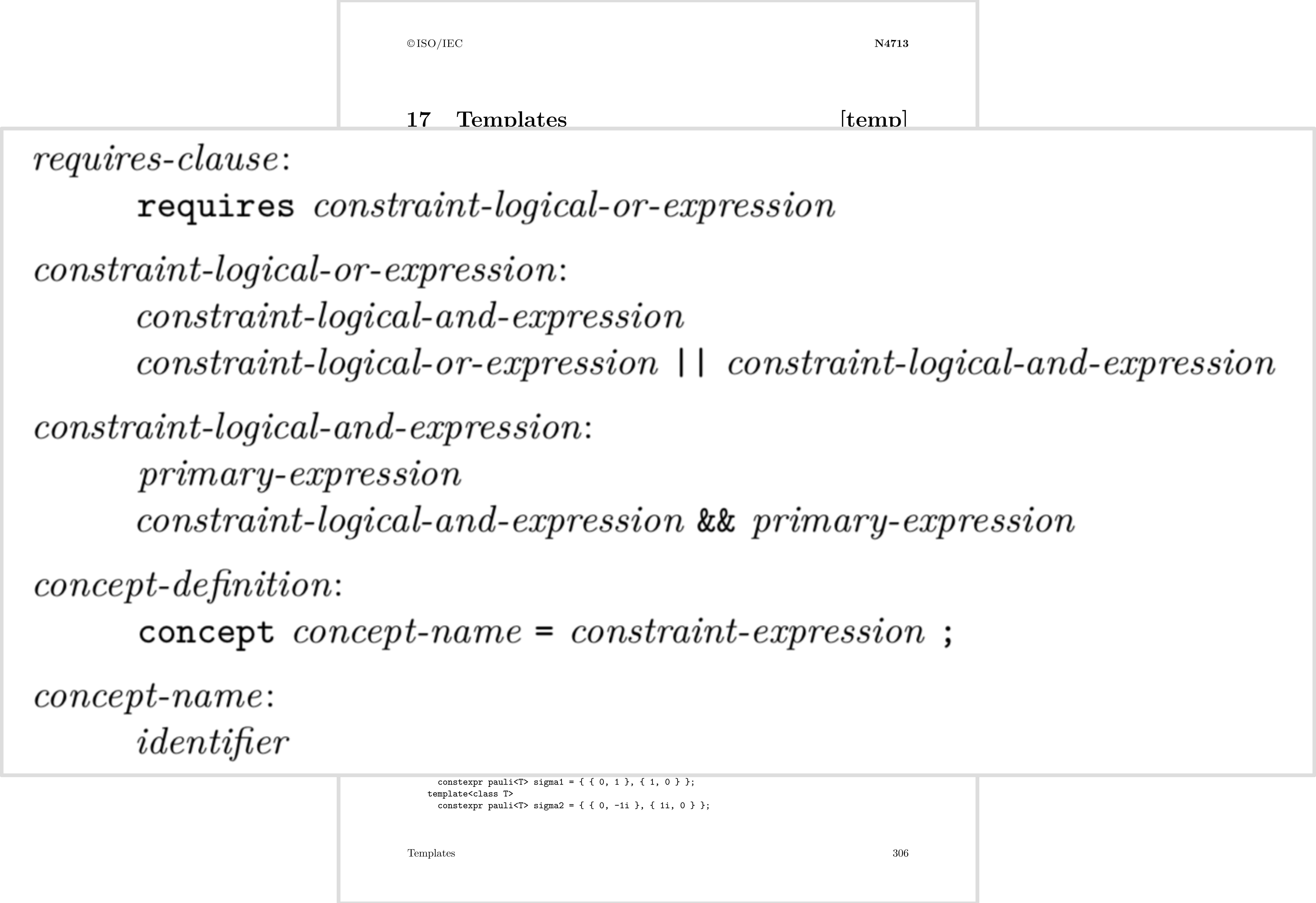
2 The *declaration* in a *template-declaration* (if any) shall

- (2.1) — declare or define a function, a class, or a variable, or
- (2.2) — define a member function, a member class, a member enumeration, or a static data member of a class template or of a class nested within a class template, or
- (2.3) — define a member template of a class or class template, or
- (2.4) — be a *deduction-guide*, or
- (2.5) — be an *alias-declaration*.

3 A *template-declaration* is a *declaration*. A *template-declaration* is also a definition if its *template-head* is followed by either a *concept-definition* or a *declaration* that defines a function, a class, a variable, or a static data member. A declaration introduced by a template declaration of a variable is a *variable template*. A variable template at class scope is a *static data member template*.

[*Example:*

```
template<class T>
  constexpr T pi = T(3.1415926535897932385L);
template<class T>
  T circular_area(T r) {
    return pi<T> * r * r;
  }
struct matrix_constants {
  template<class T>
    using pauli = hermitian_matrix<T, 2>;
  template<class T>
    constexpr pauli<T> sigma1 = { { 0, 1 }, { 1, 0 } };
  template<class T>
    constexpr pauli<T> sigma2 = { { 0, -1i }, { 1i, 0 } };
```



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17 Templates [temp]

requires-clause:
    `requires` *constraint-logical-or-expression*

constraint-logical-or-expression:
    constraint-logical-and-expression
    constraint-logical-or-expression `||` *constraint-logical-and-expression*

constraint-logical-and-expression:
    primary-expression
    constraint-logical-and-expression `&&` *primary-expression*

concept-definition:
    `concept` *concept-name* `=` *constraint-expression* `;`

concept-name:
    identifier

```
  constexpr pauli<T> sigma1 = { { 0, 1 }, { 1, 0 } };
template<class T>
  constexpr pauli<T> sigma2 = { { 0, -1i }, { 1i, 0 } };
```

Templates 306

"Generic programming is about abstracting and classifying algorithms and data structures.

It gets its inspiration from Knuth and not from type theory.

Its goal is the incremental construction of systematic catalogs of useful, efficient and abstract algorithms and data structures.

Such an undertaking is still a dream."

– Alex Stepanov

References

Much of the material in this talk can be found at http://stepanovpapers.com/

A special thanks to Paul McJones for organizing this site

Can Programming Be Liberated from the von Neumann Style?

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Exception-Safety in Generic Components

https://dl.acm.org/citation.cfm?id=724067

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Greatest Common Measure: The Last 2500 Years

https://youtu.be/fanm5y00joc

Sincere apologies to anyone I left out, your contribution was important.

Adobe